Algorithmen zur Visualisierung von Graphen Wintersemester 2019/2020

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## Exercise Sheet 5

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Exercise 1: Properties of st-Graphs *
Let $D=(V, A)$ be a planar st-graph with a given embedding. For a face $f$ of $D$ denote by $V_{f}$ and $E_{f}$ the vertices and edges on $f$. Let $\operatorname{start}(\mathrm{f})$ and $\operatorname{target}(f)$ be the source and sink of the graph $\left(V_{f}, E_{f}\right)$, respectively. Prove or disprove:
(a) $D$ is bimodal.
(b) The boundary of each face $f$ consists of two directed paths from start $(f)$ to target $(f)$.
(c) For every vertex $v \in V$ there is a simple directed st-path that contains $v$.

## Exercise 2: Duals of $s t$-Graphs

Let $D$ be a planar embedded st-graph. For a directed edge $e=(u, v)$, let $\ell(e)$ denote the face left of $e$, and let $r(e)$ denote the face right of $e$. Without loss of generality assume that $D$ is embedded such that $r(s, t)$ is the external face. The directed dual graph $D^{\star}=\left(V^{\star}, A^{\star}\right)$ of $D$ is defined as follows:

- $V^{\star}$ is the set of faces of $D$, where $s^{\star}=r(s, t)$ and $t^{\star}=\ell(s, t)$.
- $A^{\star}=\{(\ell(e), r(e)) \mid e \in A \backslash\{(s, t)\}\} \cup\left\{\left(s^{\star}, t^{\star}\right)\right\}$
(a) Prove that $D^{\star}$ is a planar $s t$-graph.
(b) Prove that for any two faces $f$ and $g$ of $D$ exactly one of the following properties holds:
i) $D$ contains a directed path from $\operatorname{target}(f)$ to $\operatorname{start}(g)$
ii) $D$ contains a directed path from $\operatorname{target}(g)$ to $\operatorname{start}(f)$
iii) $D^{\star}$ contains a directed path from $f$ to $g$
iv) $D^{\star}$ contains a directed path from $g$ to $f$

Hint: Consider a topological numbering $\sigma: V \rightarrow \mathbb{N}$ of the nodes of $D$, such that for every $(u, v) \in A$ it holds that $\sigma(u)<\sigma(v)$.

## Exercise 3: Extended Canonical Ordering for 4-Connected Graphs

A planar graph $G=(V, E)$ is called proper triangular planar (PTP, for short) if every interior face of $G$ is a triangle and the exterior face of $G$ is a quadrangle, and $G$ has no separating triangles.

Let $G=(V, E)$ be a PTP graph with vertices $a, b, c, d$ on the outer face. A labeling $v_{1}=a, v_{2}=$ $b, v_{3}, \ldots, v_{n}=d$ of the vertices of $G$ is called an extended canonical ordering of $G$ if for every $4 \leq k \leq n$ :
(i) The subgraph $G_{k-1}$ induced by $v_{1}, \ldots, v_{k-1}$ is biconnected and the boundary $C_{k-1}$ of $G_{k-1}$ contains the edge $(a, b)$, and
(ii) the vertex $v_{k}$ is on the boundary of the exterior face of $G_{k}$, and its neighbors in $G_{k-1}$ form a subinterval of the path $C_{k-1} \backslash(a, b)$ with at least two elements. If $k \leq n-2$, then $v_{k}$ has at least two neighbors in $G \backslash G_{k}$.
Let $G=(V, E)$ be a PTP graph with vertices $a, b, c, d$ on the outer face. Prove the following statements. We denote by $G_{C}$ the graph that is induced by the vertices in the interior and on the boundary of a simple cycle $C$.
(a) The graph obtained from $G$ by the removal of the vertices $c, d$ and all edges incident to them is biconnected.
(b) Let $C=\left\{a=u_{1}, \ldots, u_{m}=b, a\right\}$ be a simple cycle of $G$ such that $c, d \notin C$. Let $u_{i} \in C$, $2 \leq i \leq m-1$ such that no internal chord of $C$ is incident to $u_{i}$. Then the graph $G_{C} \backslash\left\{u_{i}\right\}$ is biconnected.
(c) Let $C$ be as above and let $\left(v_{i}, v_{j}\right), 1 \leq i<j \leq m$, be an internal chord of $C$. Then there exists a vertex $u_{l}, i<l<j$ that is adjacent to at least two vertices of $G \backslash G_{C}$.

Use the previous statements to prove the following lemma.

Lemma 1 Every PTP graph $G$ with four vertices $a, b, c, d$ on the outer face has an extended canonical ordering such that $v_{1}=a, v_{2}=b, v_{n-1}=c, v_{n}=d$.

