

# Computational Geometry – WSPD

## WSPD

LEHRSTUHL FÜR ALGORITHMIK I · INSTITUT FÜR THEORETISCHE INFORMATIK · FAKULTÄT FÜR INFORMATIK

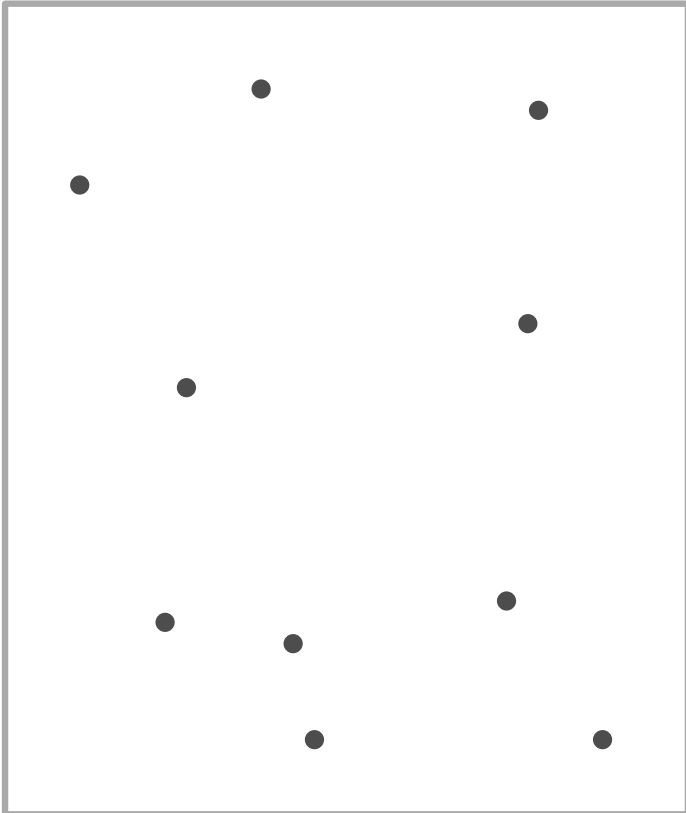
Guido Brückner  
06.07.2018



# Motivation: Spanners

## Task:

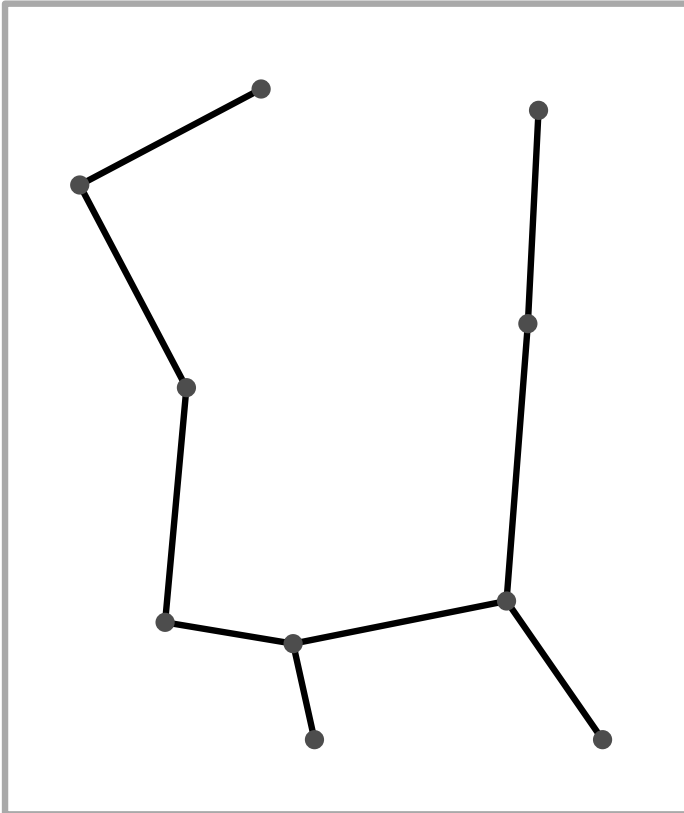
A set of cities shall be connected by a new road network.



# Motivation: Spanners

## Task:

A set of cities shall be connected by a new road network.



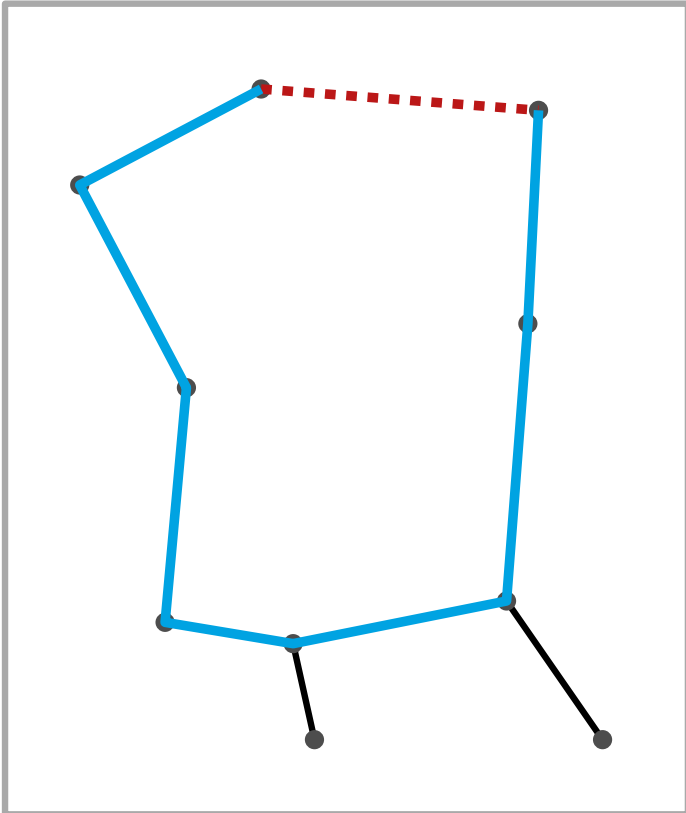
**Idea 1:** Euclidean minimum spanning tree

# Motivation: Spanners

## Task:

A set of cities shall be connected by a new road network.

But for no pair  $(x, y)$  the path length in the road network should be much larger than the distance  $\|xy\|$ .



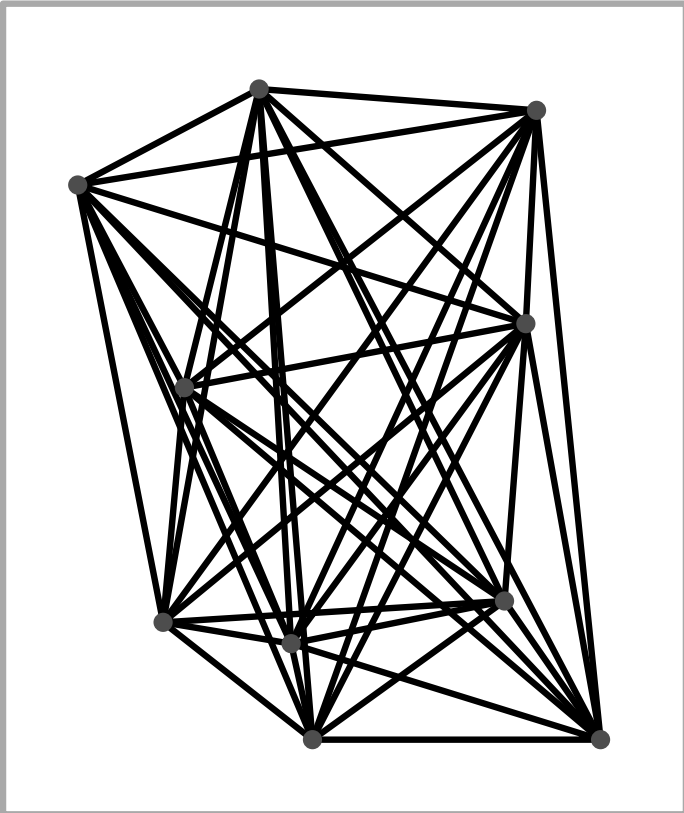
**Idea 1:** Euclidean minimum spanning tree

# Motivation: Spanners

## Task:

A set of cities shall be connected by a new road network.

But for no pair  $(x, y)$  the path length in the road network should be much larger than the distance  $\|xy\|$ .



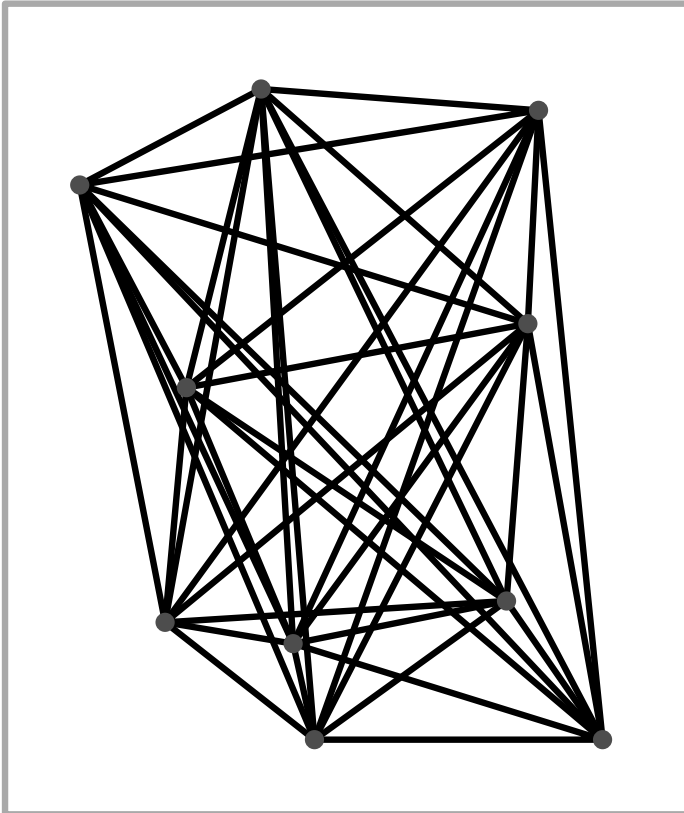
**Idea 1:** Euclidean minimum spanning tree

**Idea 2:** complete graph

# Motivation: Spanners

## Task:

A set of cities shall be connected by a new road network.



But for no pair  $(x, y)$  the path length in the road network should be much larger than the distance  $\|xy\|$ .

Construction costs must remain reasonable, e.g., only  $O(n)$  edges.

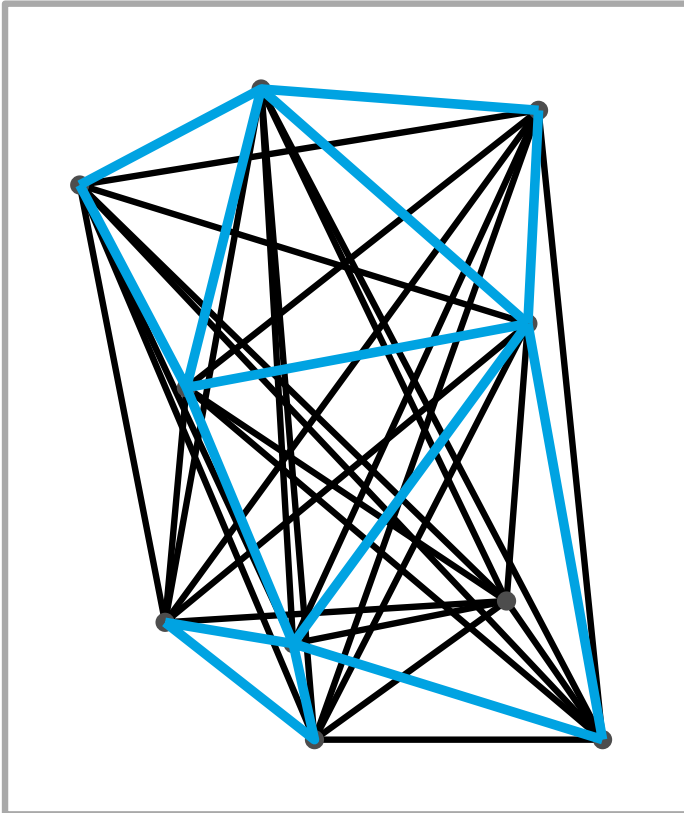
**Idea 1:** Euclidean minimum spanning tree

**Idea 2:** complete graph

# Motivation: Spanners

## Task:

A set of cities shall be connected by a new road network.



But for no pair  $(x, y)$  the path length in the road network should be much larger than the distance  $\|xy\|$ .

Construction costs must remain reasonable, e.g., only  $O(n)$  edges.

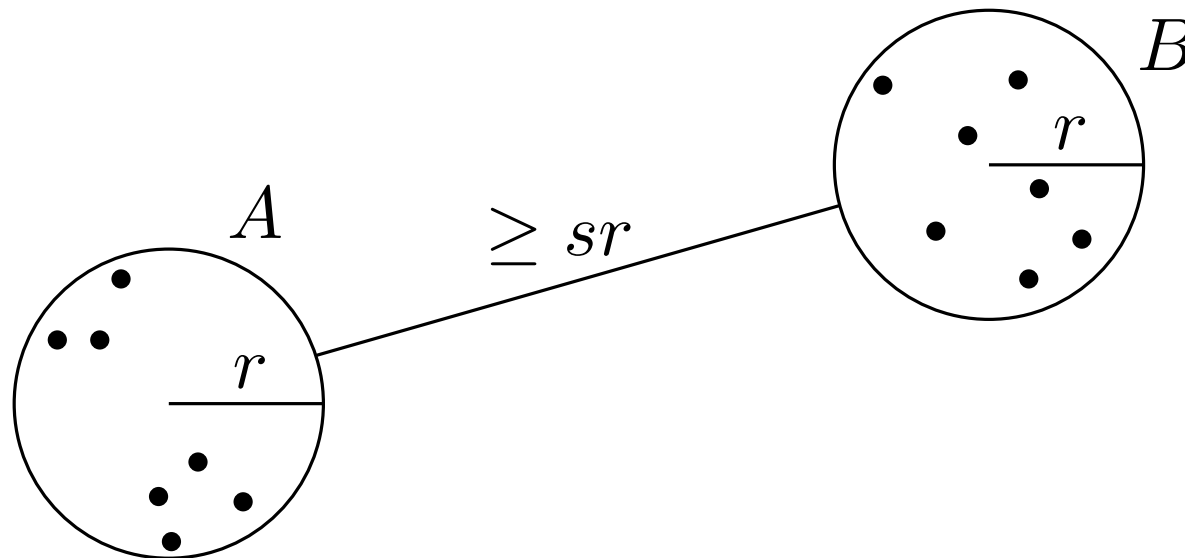
**Idea 1:** Euclidean minimum spanning tree

**Idea 2:** complete graph

**Idea 3:** sparse  $t$ -spanner

# Well-Separated Pairs

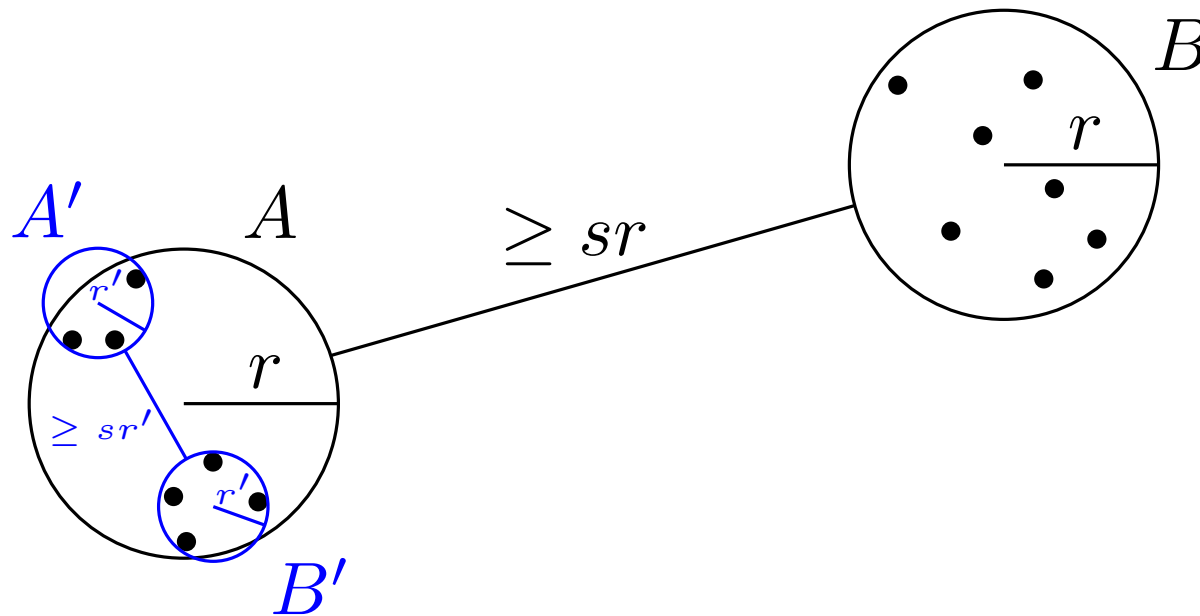
**Def:** A pair of disjoint point sets  $A$  and  $B$  in  $\mathbb{R}^d$  is called  **$s$ -well separated** for some  $s > 0$ , if  $A$  and  $B$  can each be covered by a ball of radius  $r$  whose distance is at least  $sr$ .





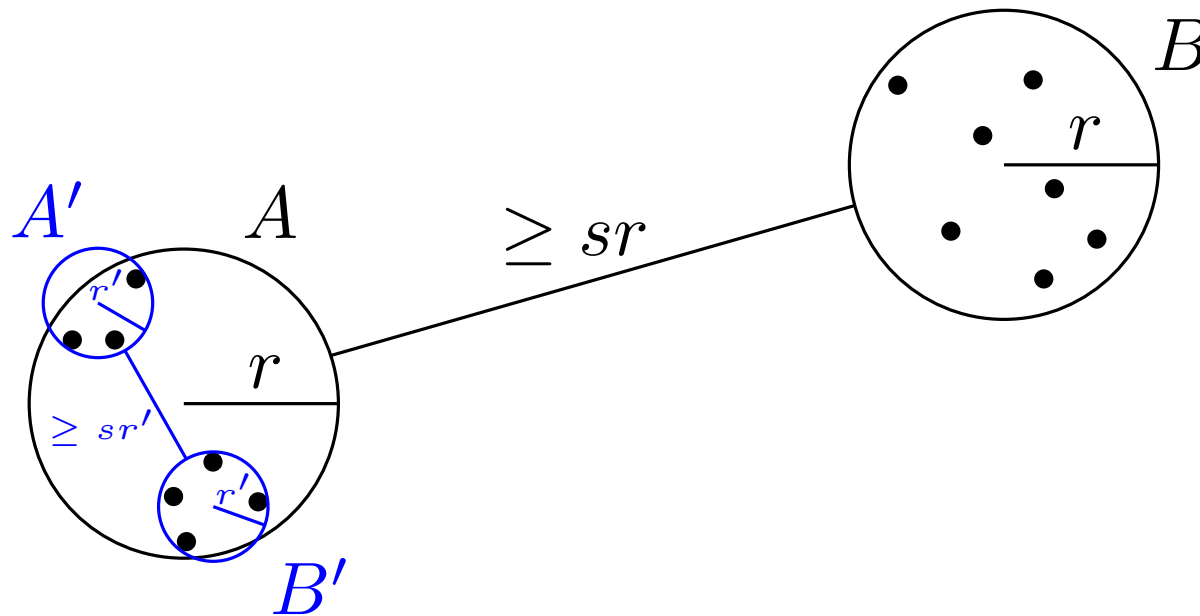
# Well-Separated Pairs

**Def:** A pair of disjoint point sets  $A$  and  $B$  in  $\mathbb{R}^d$  is called  **$s$ -well separated** for some  $s > 0$ , if  $A$  and  $B$  can each be covered by a ball of radius  $r$  whose distance is at least  $sr$ .



# Well-Separated Pairs

**Def:** A pair of disjoint point sets  $A$  and  $B$  in  $\mathbb{R}^d$  is called  **$s$ -well separated** for some  $s > 0$ , if  $A$  and  $B$  can each be covered by a ball of radius  $r$  whose distance is at least  $sr$ .



- Obs:**
- $s$ -well separated  $\Rightarrow$   $s'$ -well separated for all  $s' \leq s$
  - singletons  $\{a\}$  and  $\{b\}$  are  $s$ -well separated for all  $s > 0$

# Well-Separated Pair Decomposition (WSPD)

For well-separated pair  $\{A, B\}$  we know that the distance for all point pairs in  $A \otimes B = \{\{a, b\} \mid a \in A, b \in B, a \neq b\}$  is similar.

**Goal:**  $o(n^2)$ -sized data structure that approximates the distances of all  $\binom{n}{2}$  pairs of points in a set  $P = \{p_1, \dots, p_n\}$ .

# Well-Separated Pair Decomposition (WSPD)

For well-separated pair  $\{A, B\}$  we know that the distance for all point pairs in  $A \otimes B = \{\{a, b\} \mid a \in A, b \in B, a \neq b\}$  is similar.

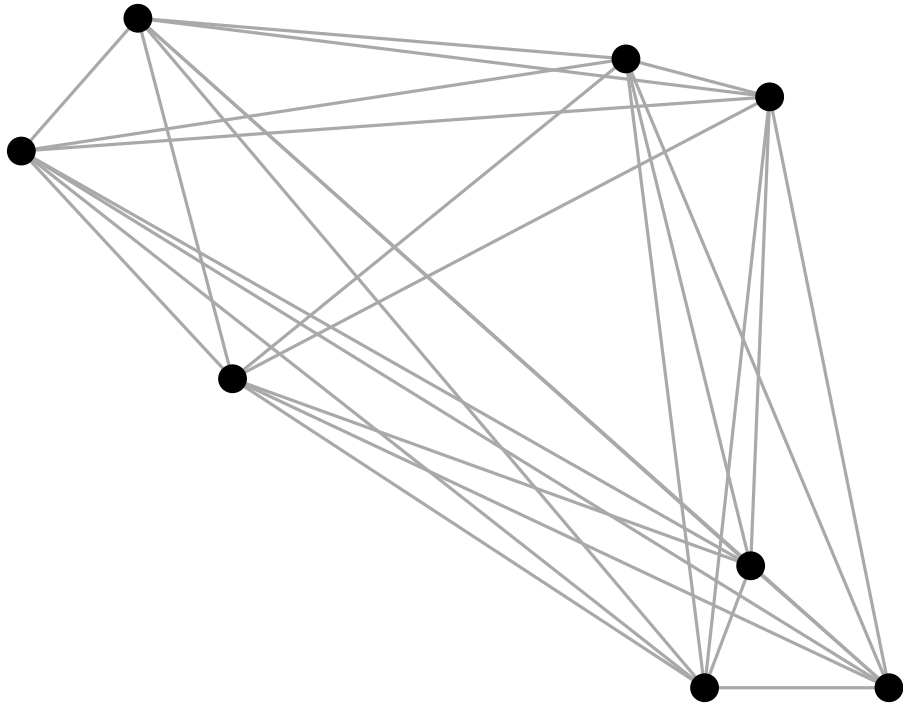
**Goal:**  $o(n^2)$ -sized data structure that approximates the distances of all  $\binom{n}{2}$  pairs of points in a set  $P = \{p_1, \dots, p_n\}$ .

**Def:** For a point set  $P$  and some  $s > 0$  an  $s$ -**well separated pair decomposition** ( $s$ -WSPD) is a set of pairs

$\{\{A_1, B_1\}, \dots, \{A_m, B_m\}\}$  with

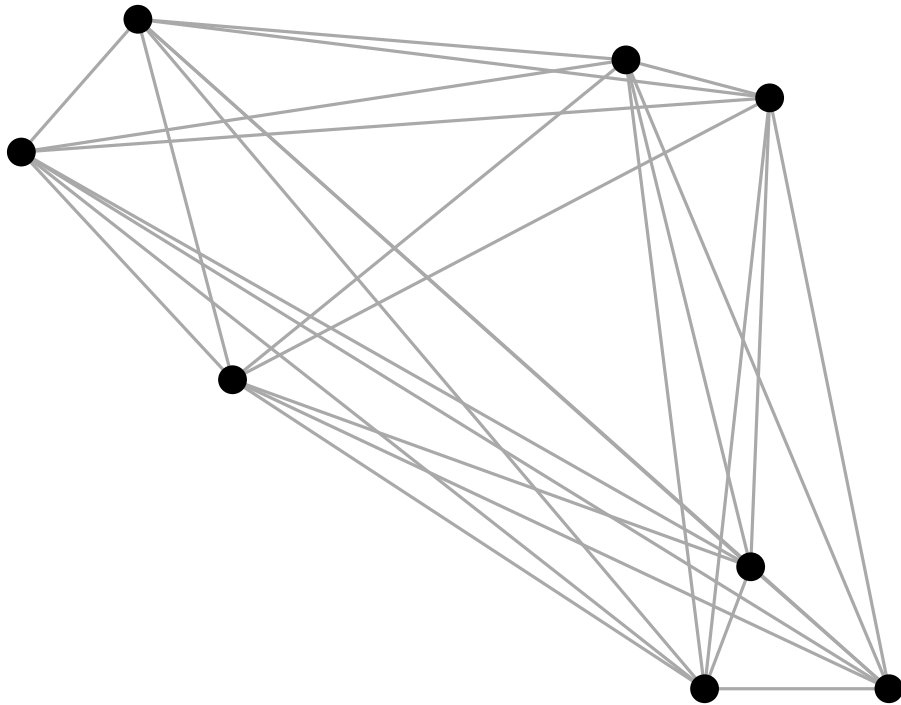
- $A_i, B_i \subset P$  for all  $i$
- $A_i \cap B_i = \emptyset$  for all  $i$
- $\bigcup_{i=1}^m A_i \otimes B_i = P \otimes P$
- $\{A_i, B_i\}$   $s$ -well separated for all  $i$

# Example

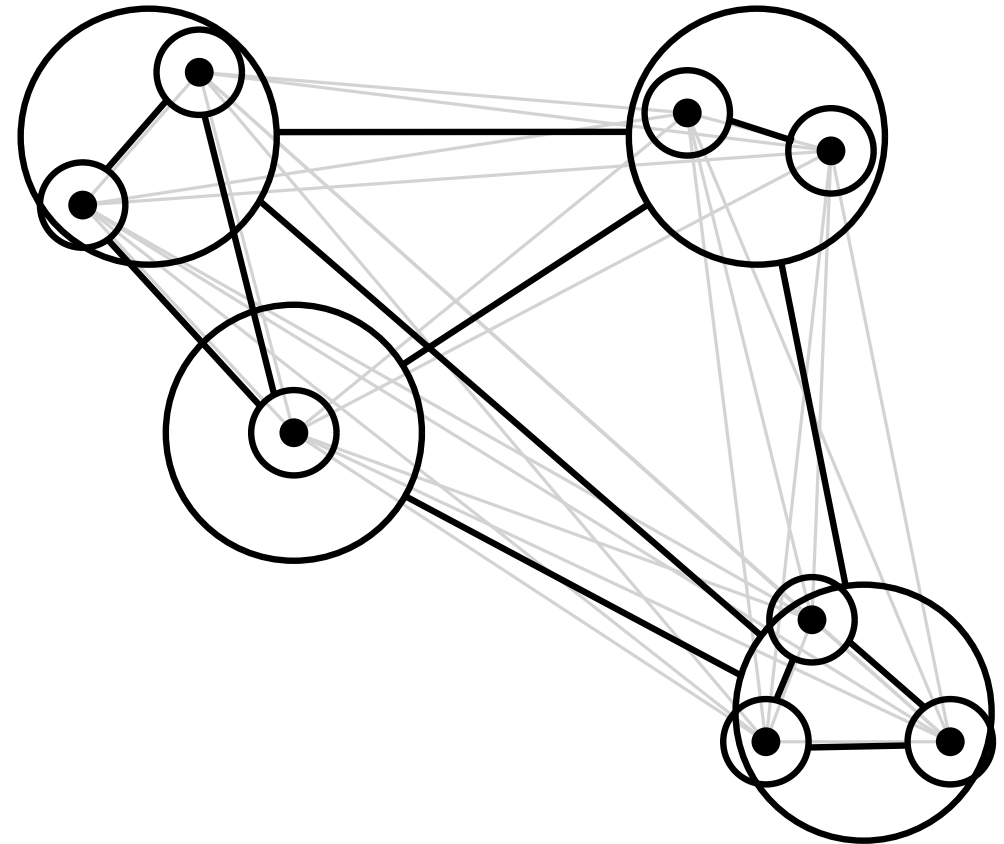


28 point pairs

# Example

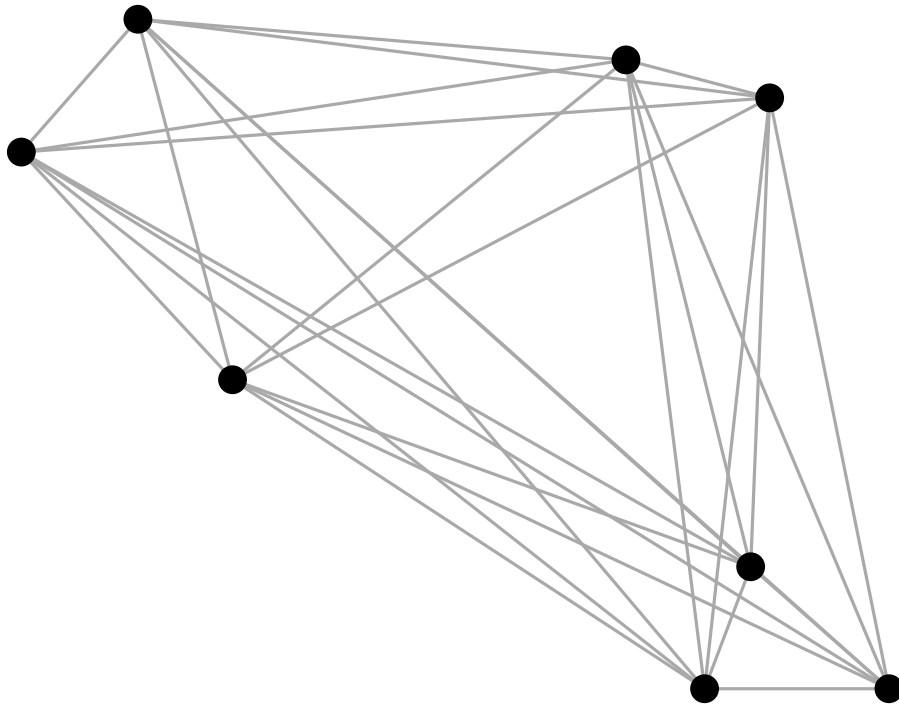


28 point pairs

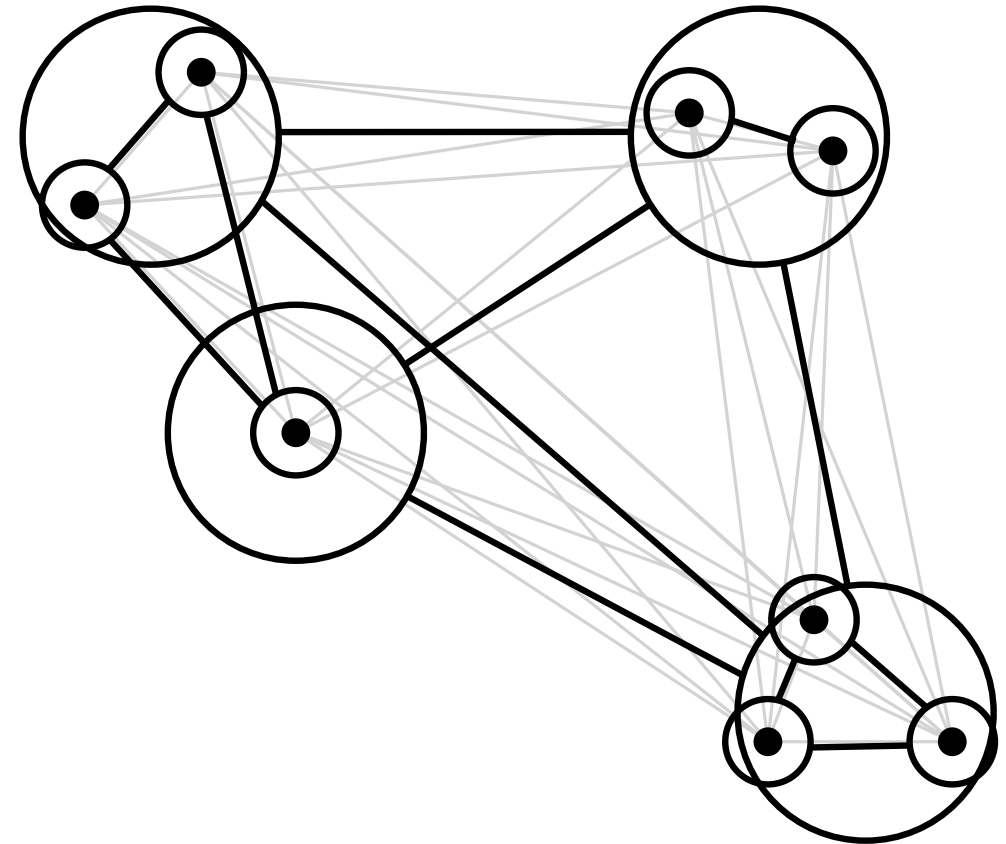


12  $s$ -well separated pairs

# Example



28 point pairs



12  $s$ -well separated pairs

**Thm 3:** Given a point set  $P$  in  $\mathbb{R}^d$  and  $s \geq 1$  we can construct an  $s$ -WSPD with  $O(s^d n)$  pairs in time  $O(n \log n + s^d n)$ .

# Exercise 5

- $x := 2/s + 1$
- $S := \{x^i \mid 0 \leq i \leq n - 1\}$

$\mathcal{W} = \{A_j, B_j\}$  arbitrary  $s$ -WSPD for  $S$  ( $s > 0$ )  
 $1 \leq j \leq m$

**Show:**

$$\sum_{j=1}^m (|A_j| + |B_j|) = \binom{n}{2} + m$$

*Hint:* Show that for each  $j$  at least one of both sets  $A_j$  and  $B_j$  is a singleton.



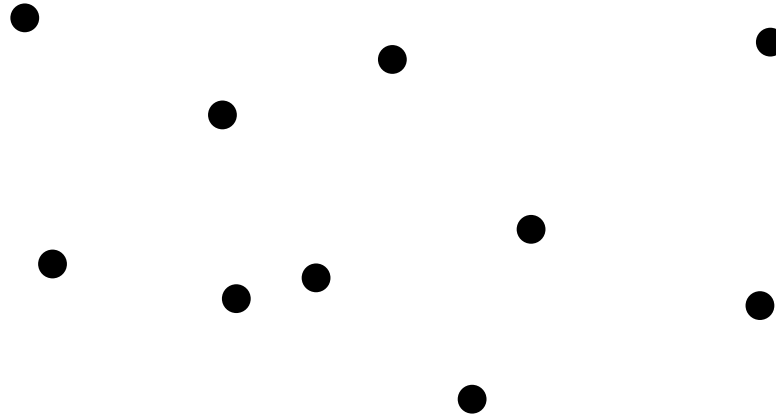
**Def.:** For a point set  $P$  and some  $s > 0$  an  **$s$ -well separated pair decomposition** ( $s$ -WSPD) is a set of pairs

$\{\{A_1, B_1\}, \dots, \{A_m, B_m\}\}$  with

- $A_i, B_i \subset P$  for all  $i$
- $\{A_i, B_i\}$   $s$ -well separated for all  $i$
- for two distinct points  $p, q \in P$  there is exactly one index  $i$  with  $1 \leq i \leq m$  such that
  - $p \in A_i$  and  $q \in B_i$ , or
  - $q \in A_i$  and  $p \in B_i$ .

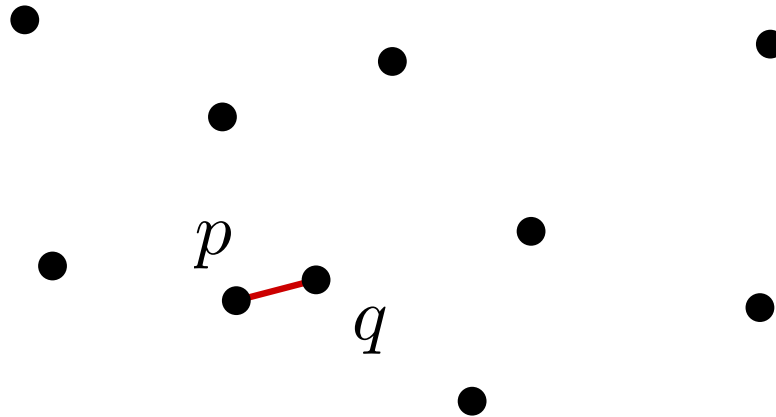
# Exercise 6/7

- $P$ :  $n$  Punkte aus dem  $\mathbb{R}^d$



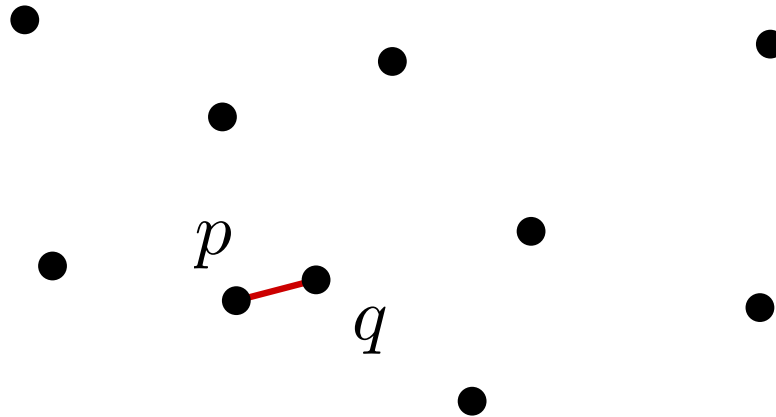
# Exercise 6/7

- $P$ :  $n$  Punkte aus dem  $\mathbb{R}^d$
- $p, q \in P$  and  $q$  is the next neighbor of  $p$



# Exercise 6/7

- $P$ :  $n$  Punkte aus dem  $\mathbb{R}^d$
- $p, q \in P$  and  $q$  is the next neighbor of  $p$



**Given:**  $s$ -WSPD  $\mathcal{W}$  for  $P$  with  $s > 2$

Let  $\{A, B\} \in \mathcal{W}$  with  $p \in A$  and  $q \in B$

**Show that:**

- $A$  is a singleton.
- size of  $\mathcal{W}$  is at least  $n/2$ .
- if  $p, q$  have minimal distance among all pairs, then  $\{\{p\}, \{q\}\}$  lies in WSPD.