

# Algorithmen zur Visualisierung von Graphen

## Kompaktierung von orthogonalen Layouts

Vorlesung im Sommersemester 2009

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# Flussnetzwerk Längenzuweisung

Definition Flussnetzwerk  $N_{\text{hor}} = ((W_{\text{hor}}, A_{\text{hor}}); l; u; b; \text{cost})$

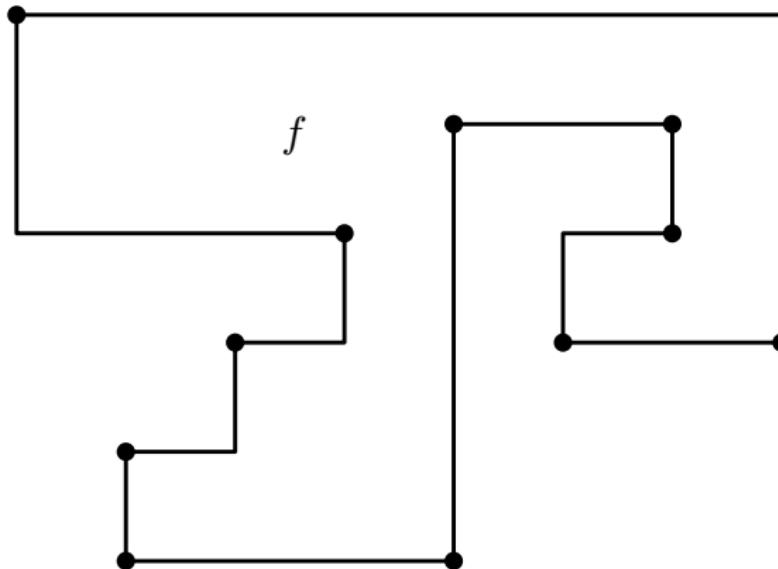
- »  $W_{\text{hor}} = \mathcal{F}$
- »  $A_{\text{hor}} = \{(f, g) \mid f, g \text{ besitzen gemeinsames horizontales Kantensegment und } f \text{ liegt unterhalb von } g\}$
- »  $l(a) = 1 \quad \forall a \in A_{\text{hor}}$
- »  $u(a) = \infty \quad \forall a \in A_{\text{hor}}$
- »  $\text{cost}(a) = 1 \quad \forall a \in A_{\text{hor}}$
- »  $b(f) = 0 \quad \forall f \in W_{\text{hor}}$

# Flussnetzwerk Längenzuweisung

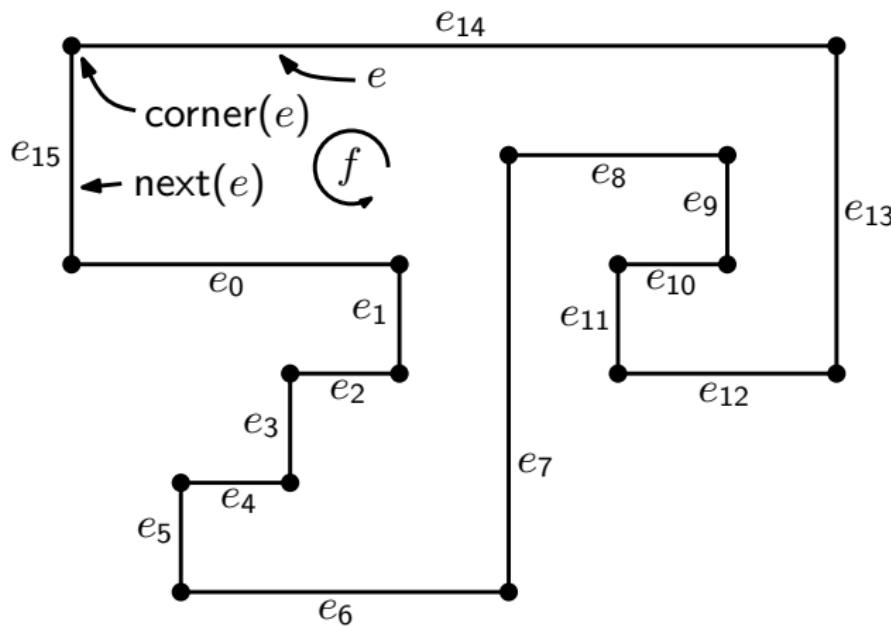
Definition Flussnetzwerk  $N_{\text{vert}} = ((W_{\text{vert}}, A_{\text{vert}}); l; u; b; \text{cost})$

- »  $W_{\text{vert}} = \mathcal{F}$
- »  $A_{\text{vert}} = \{(f, g) \mid f, g \text{ besitzen gemeinsames vertikales Kantensegment und } f \text{ liegt links von } g\}$
- »  $l(a) = 1 \quad \forall a \in A_{\text{vert}}$
- »  $u(a) = \infty \quad \forall a \in A_{\text{vert}}$
- »  $\text{cost}(a) = 1 \quad \forall a \in A_{\text{vert}}$
- »  $b(f) = 0 \quad \forall f \in W_{\text{vert}}$

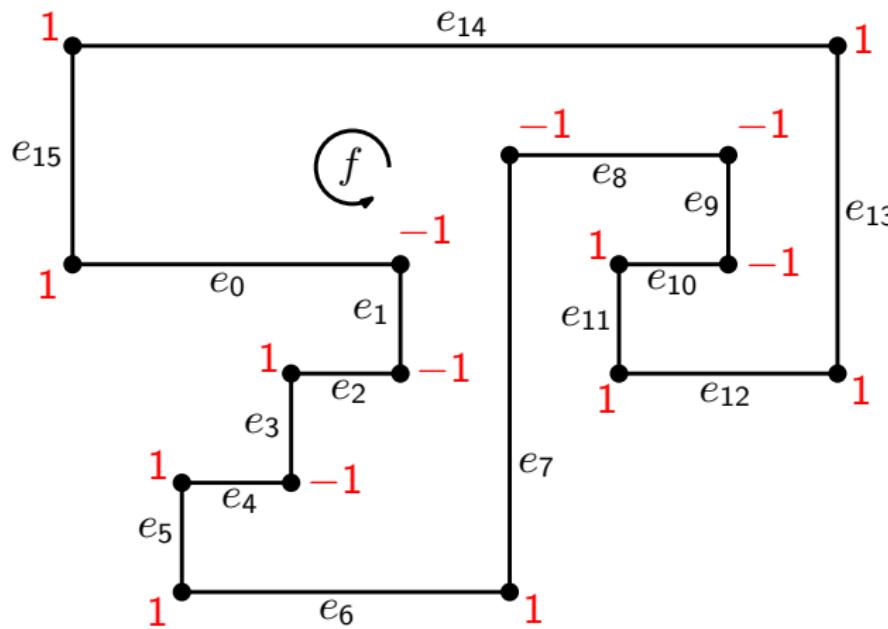
# Verfeinerung von $(G, H)$ – innere Facette



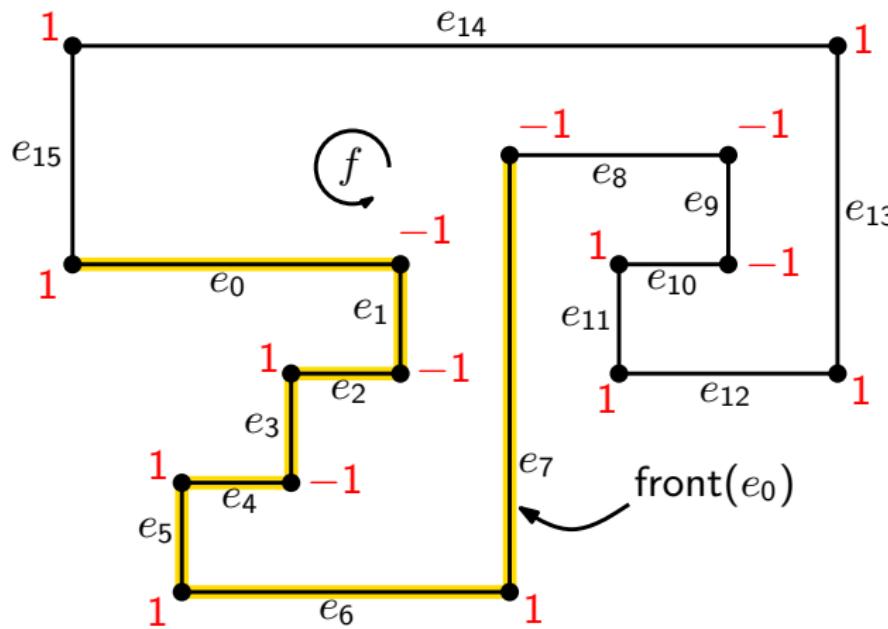
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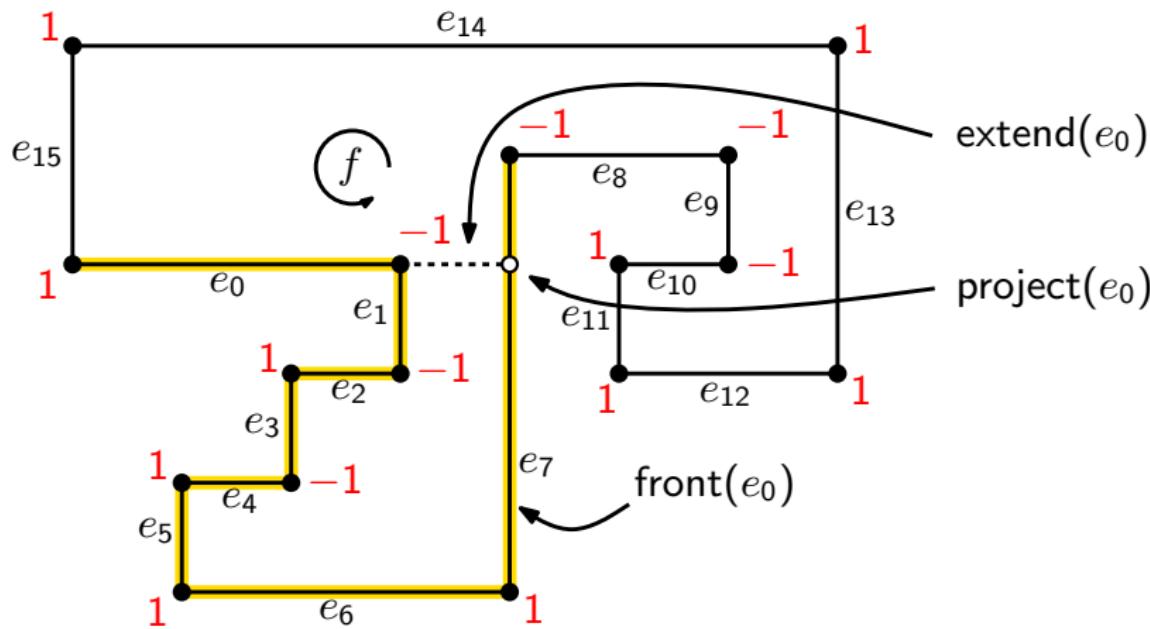
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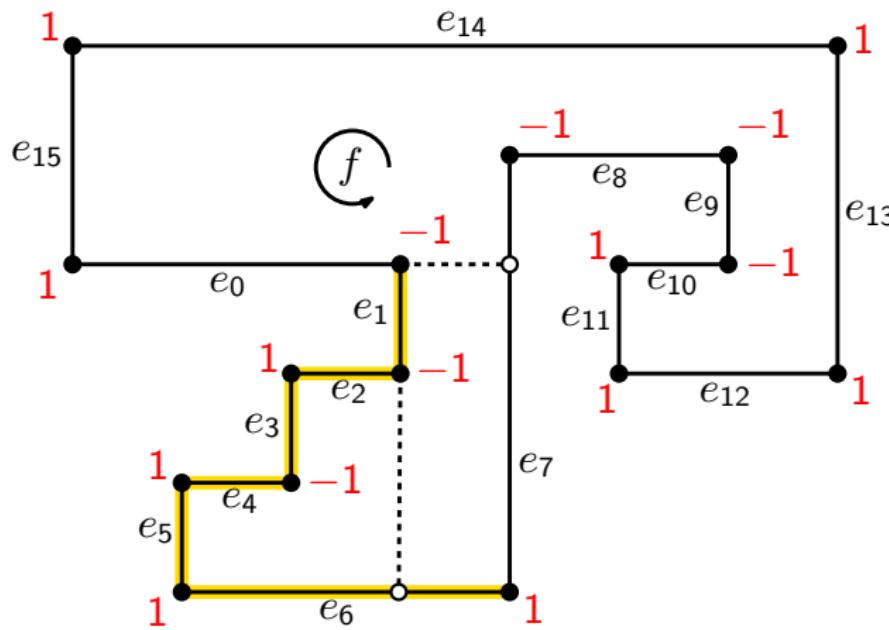
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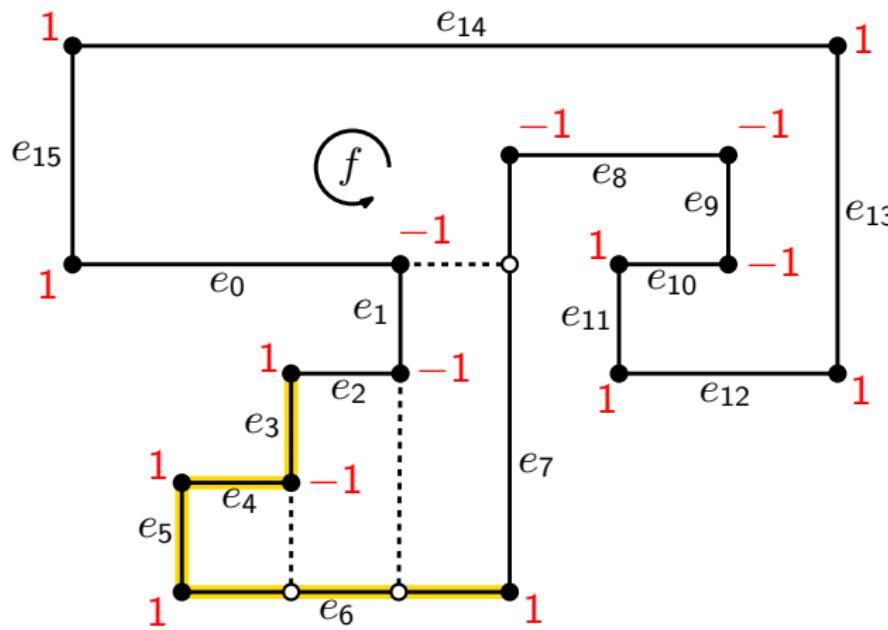
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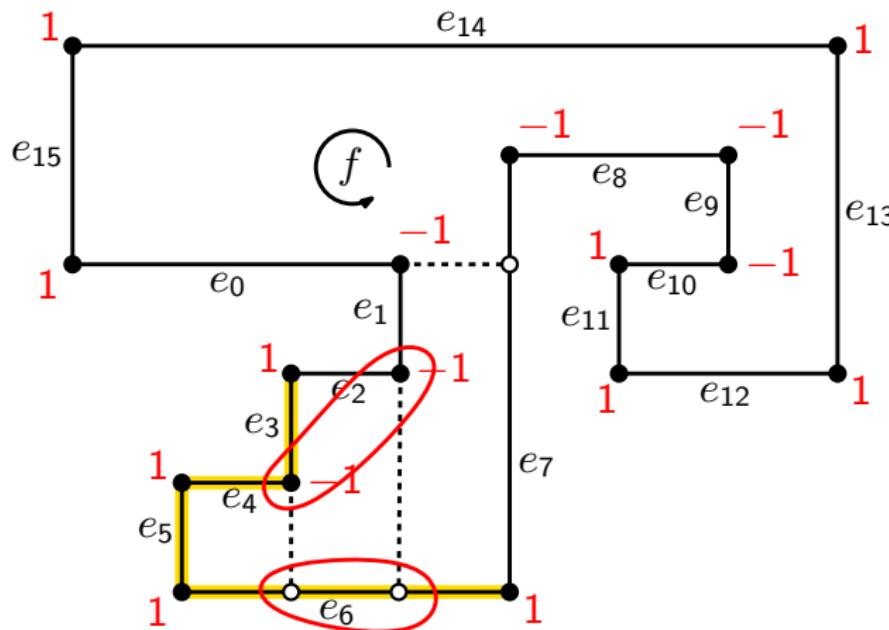
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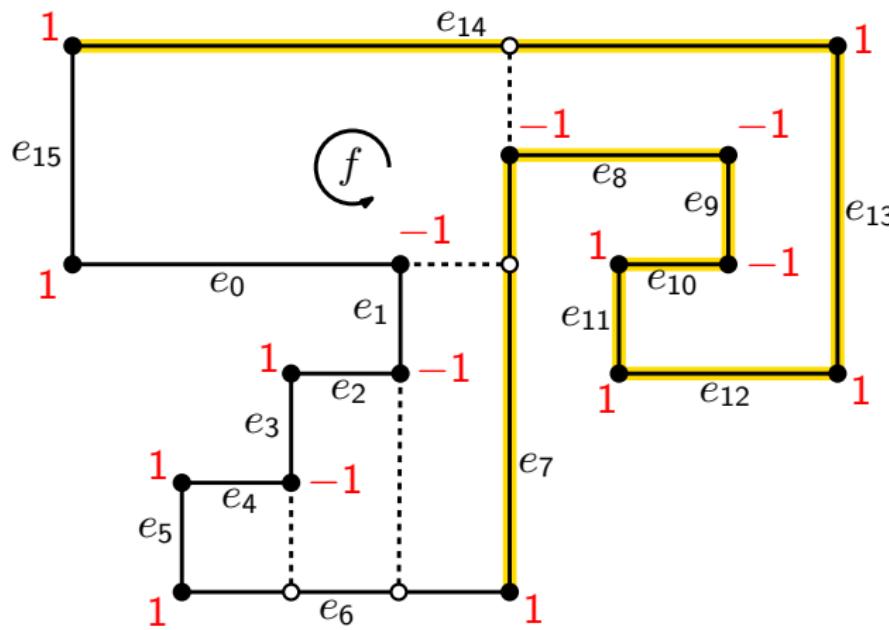
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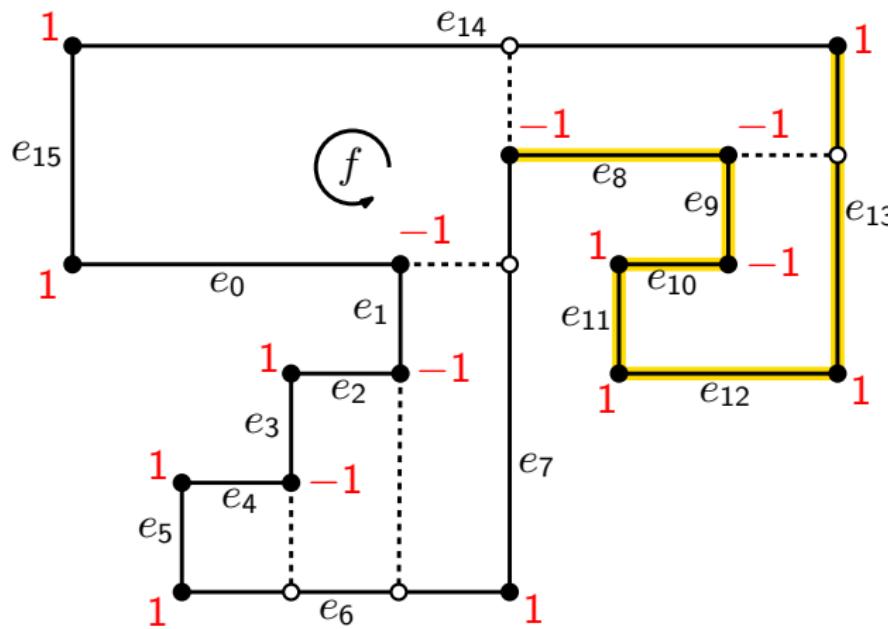
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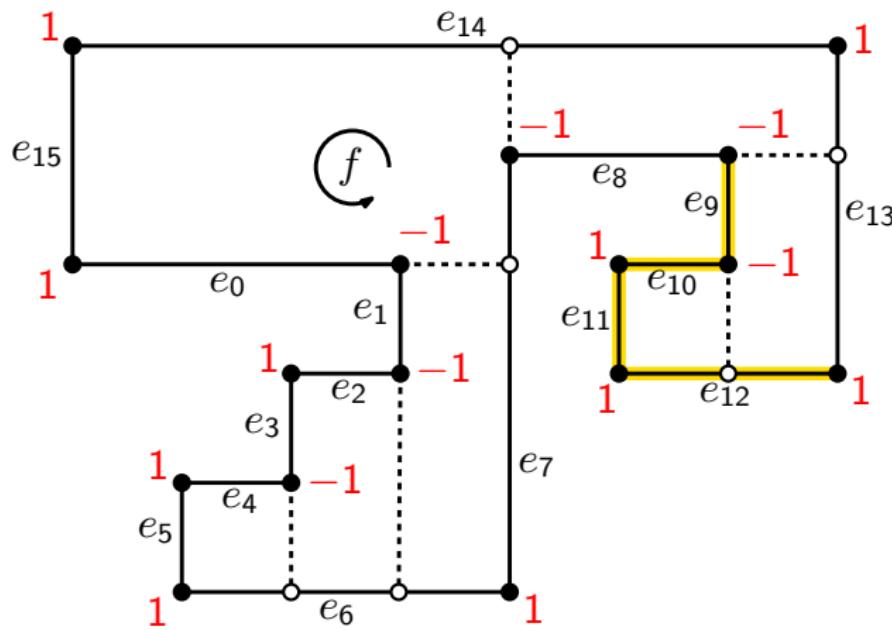
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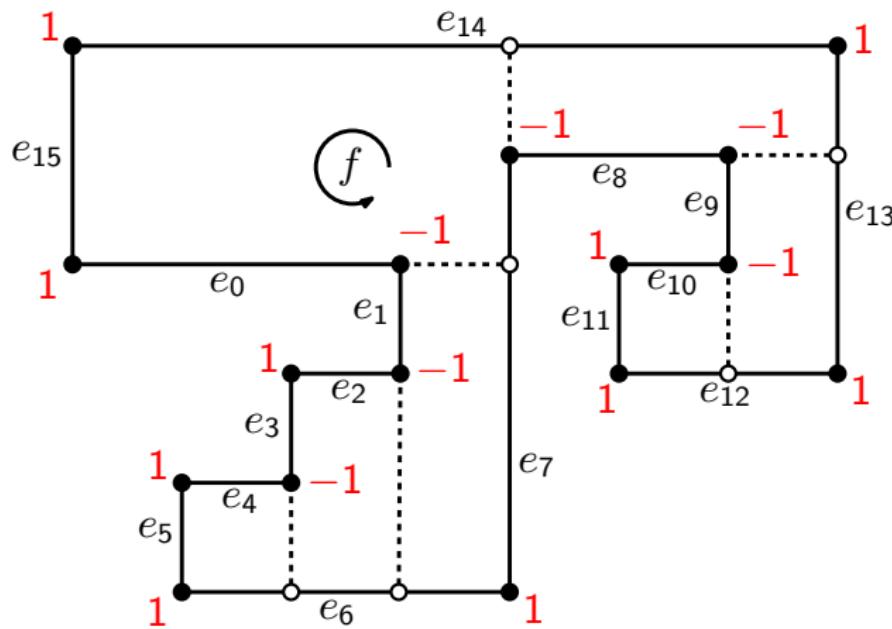
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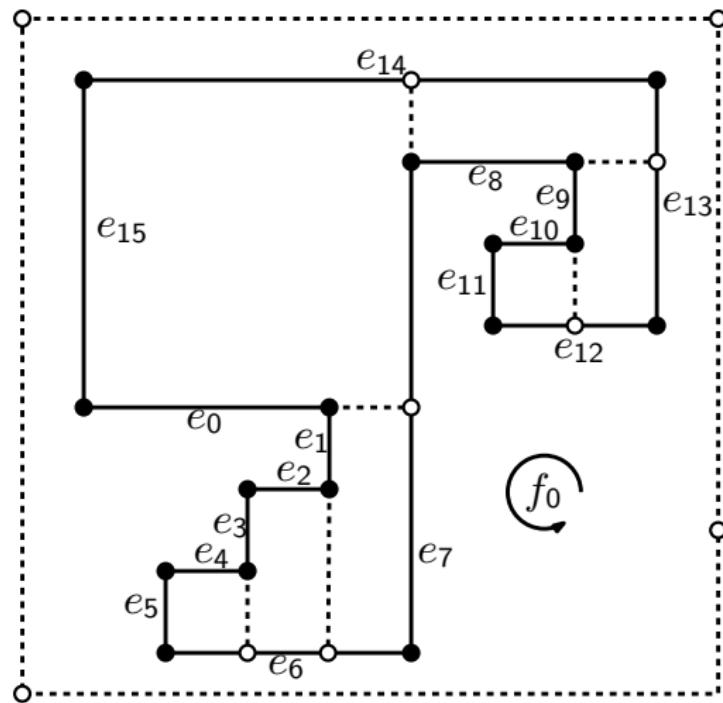
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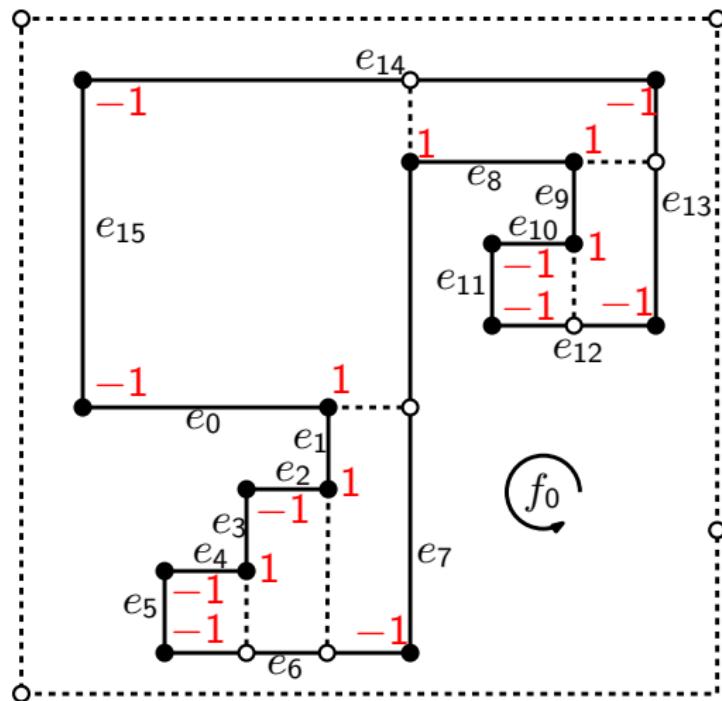
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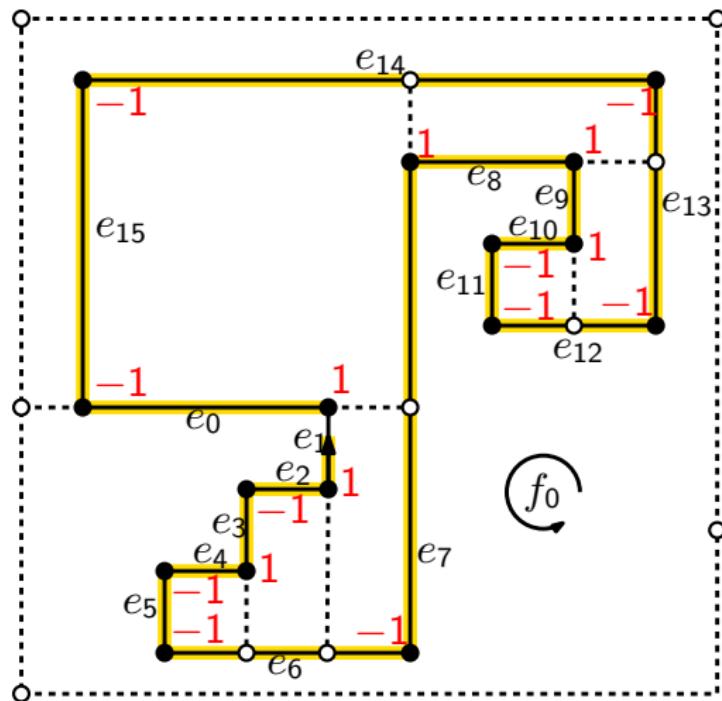
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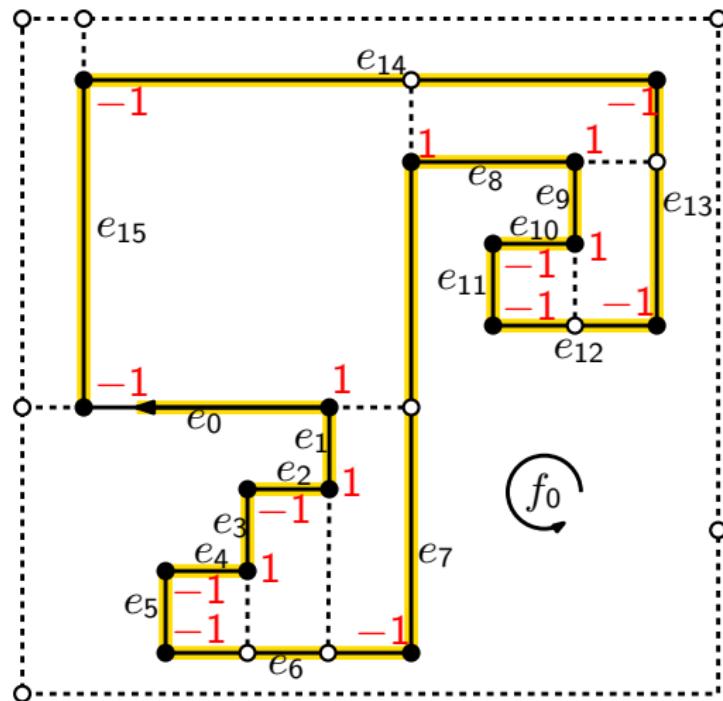
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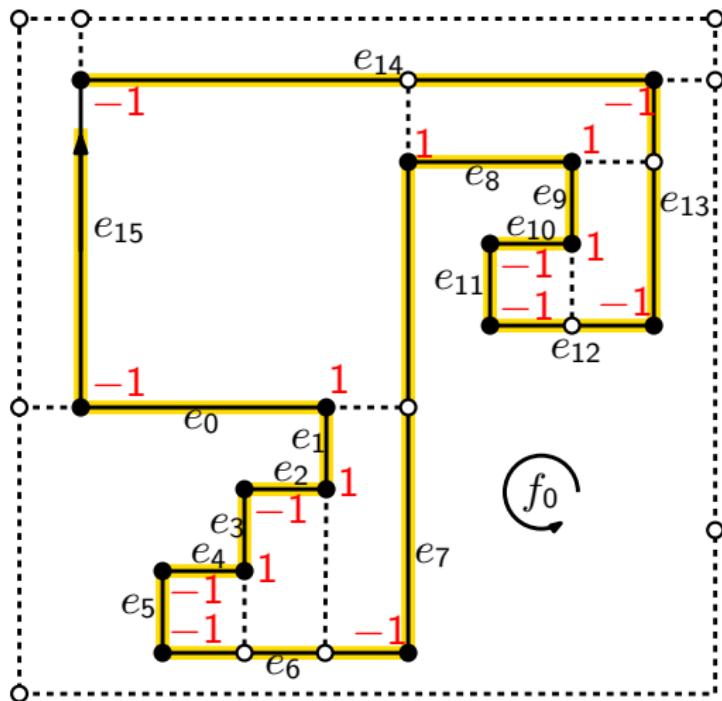
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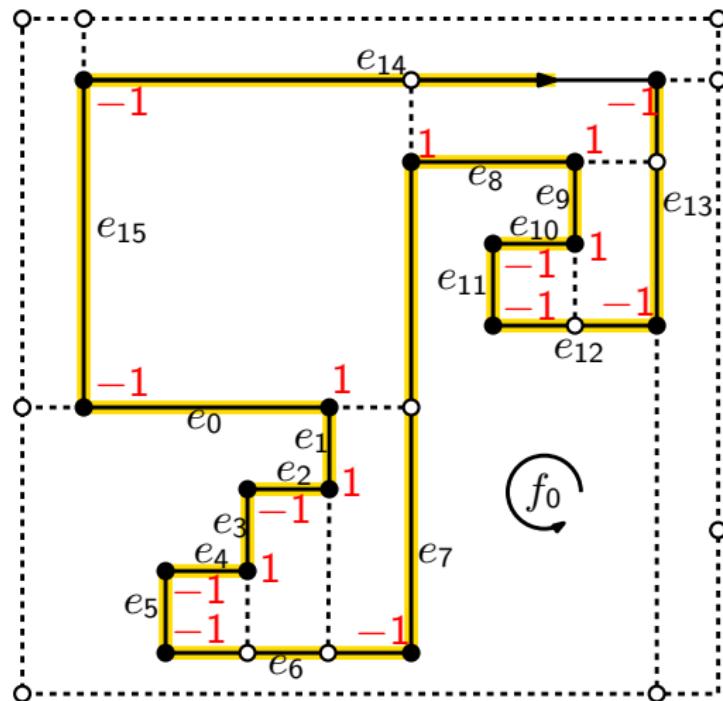
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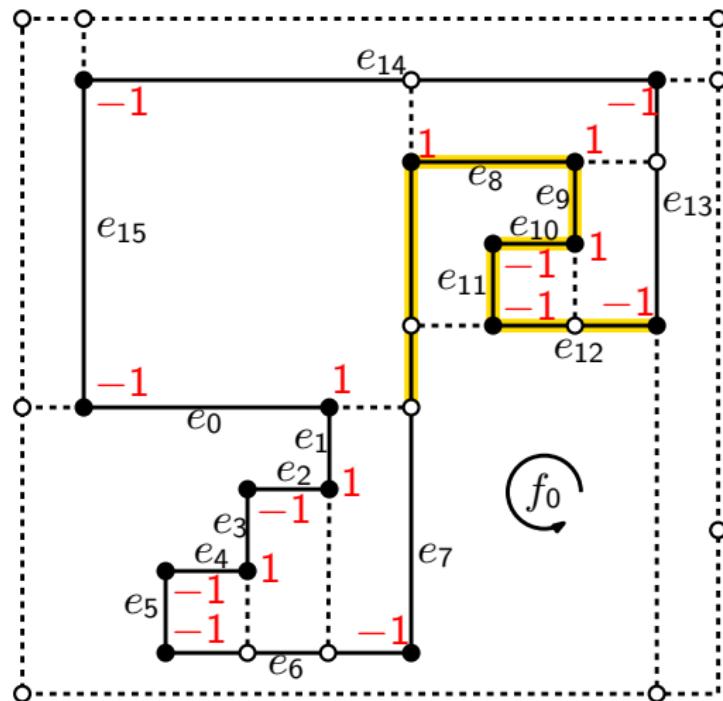
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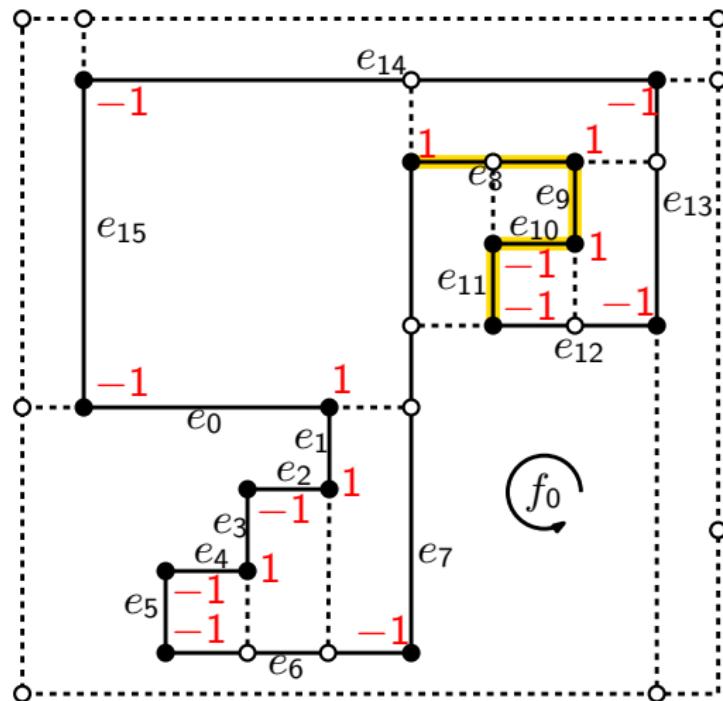
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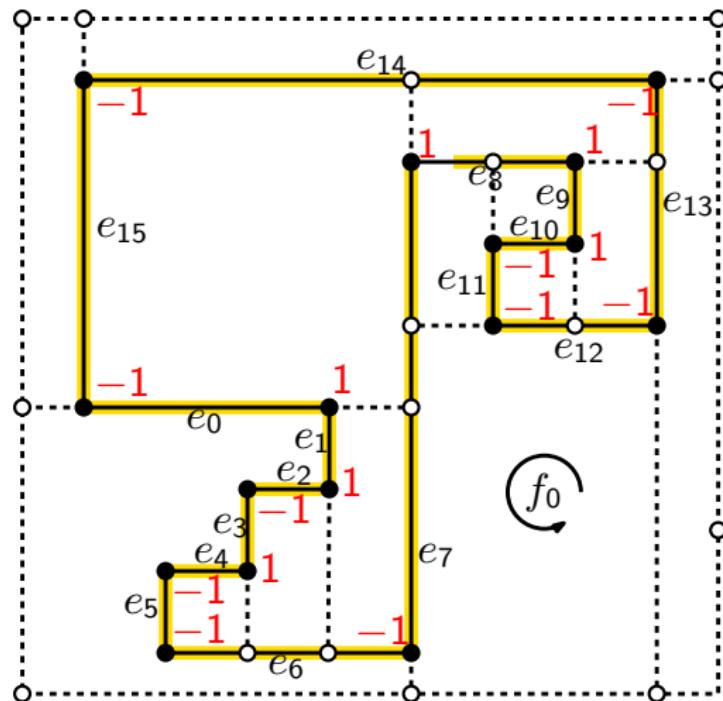
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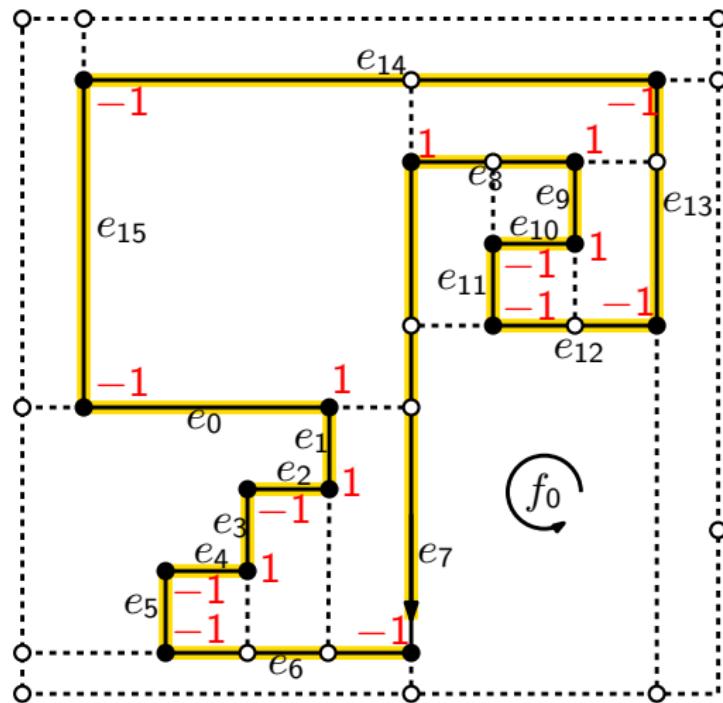
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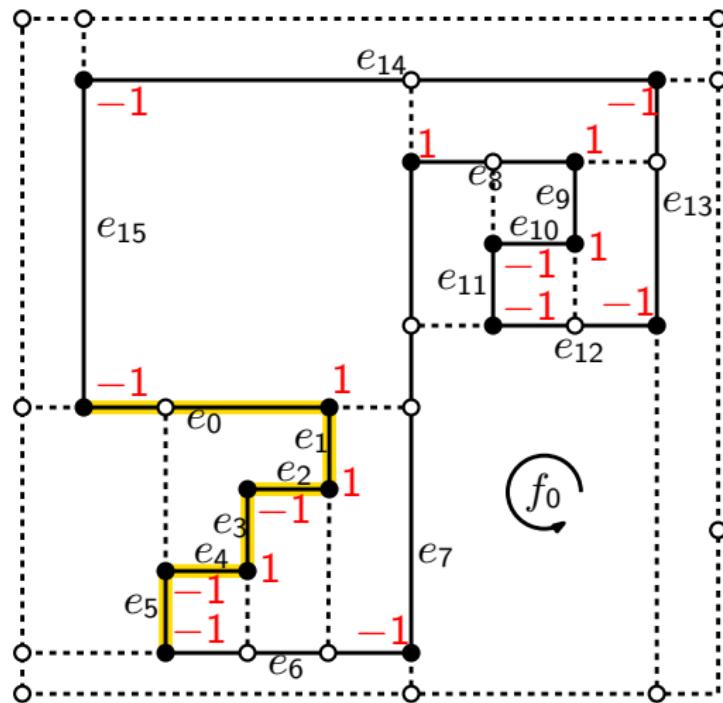
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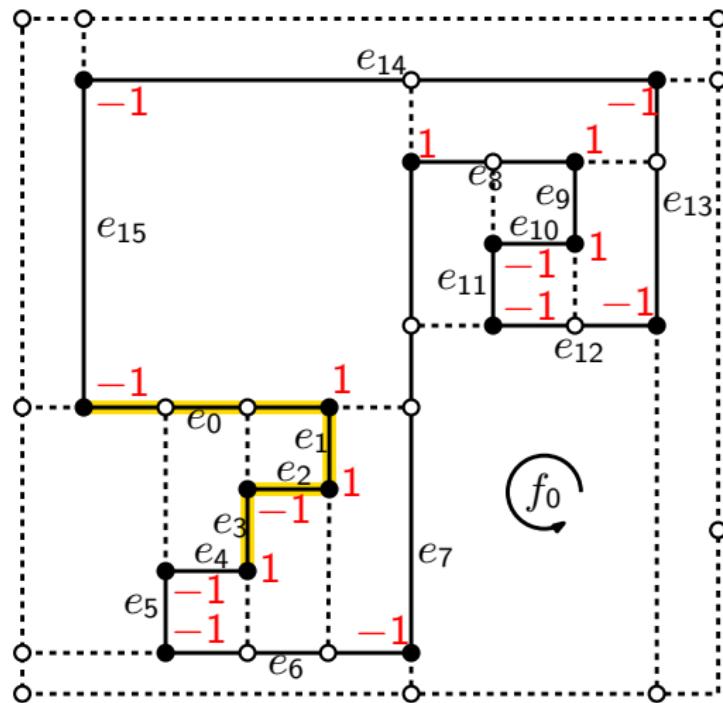
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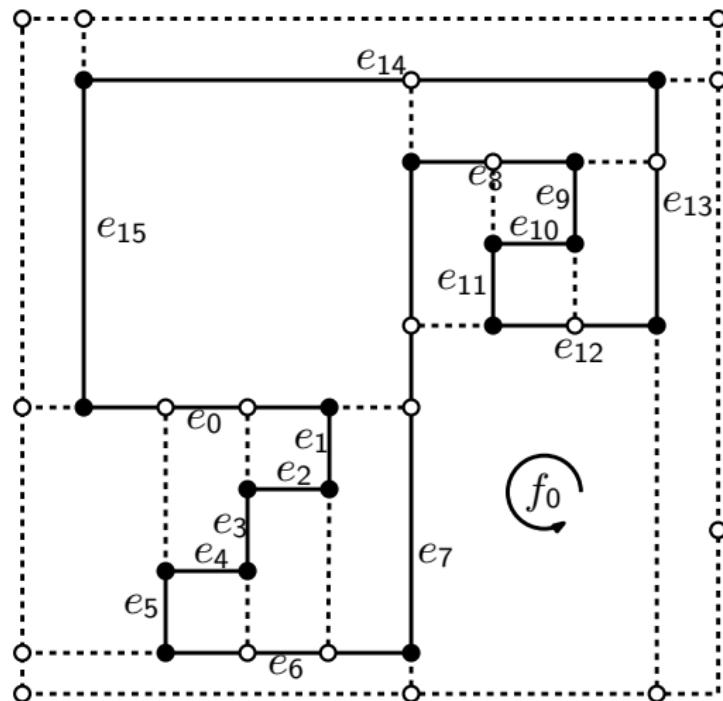
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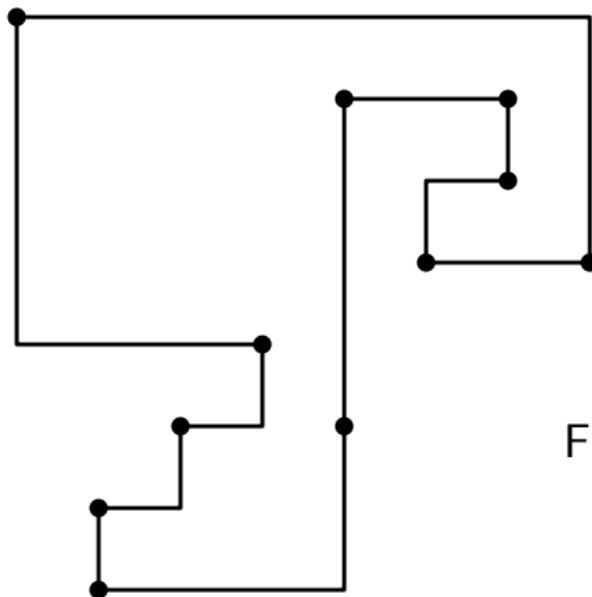
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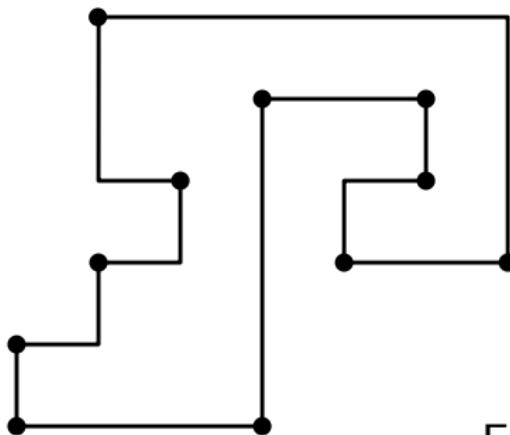
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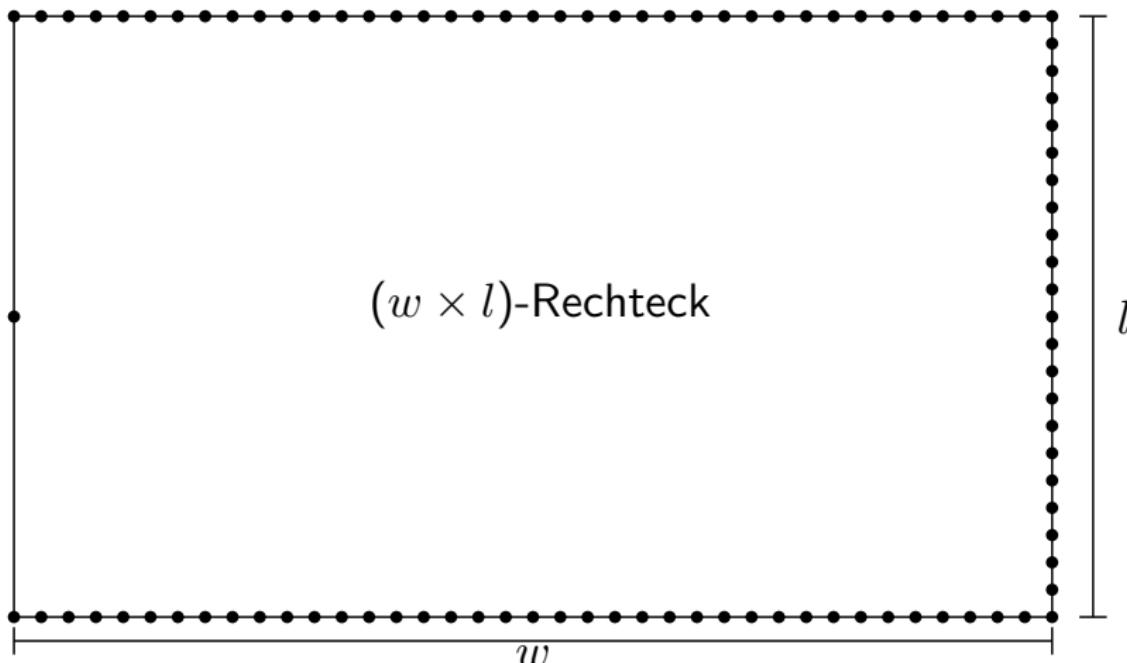
Flächenminimal?

Nein!

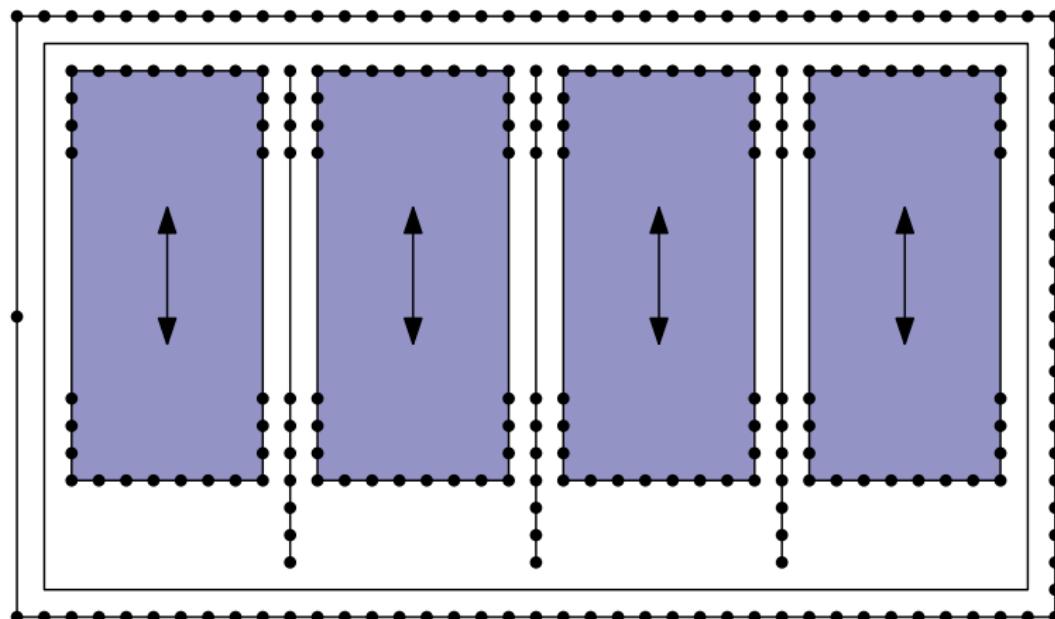
# Kompaktierung ist NP-schwer [Patrignani '01]

- » Grobstruktur von  $(G, H)$
- » Begrenzung
- » Gürtel
- » Klauselgadgets
- » Variablengadgets
- » bestimme geeigneten Wert  $K$
- »  $(G, H)$  lässt sich in Fläche  $K$  zeichnen gdw.  $\Phi$  erfüllbar

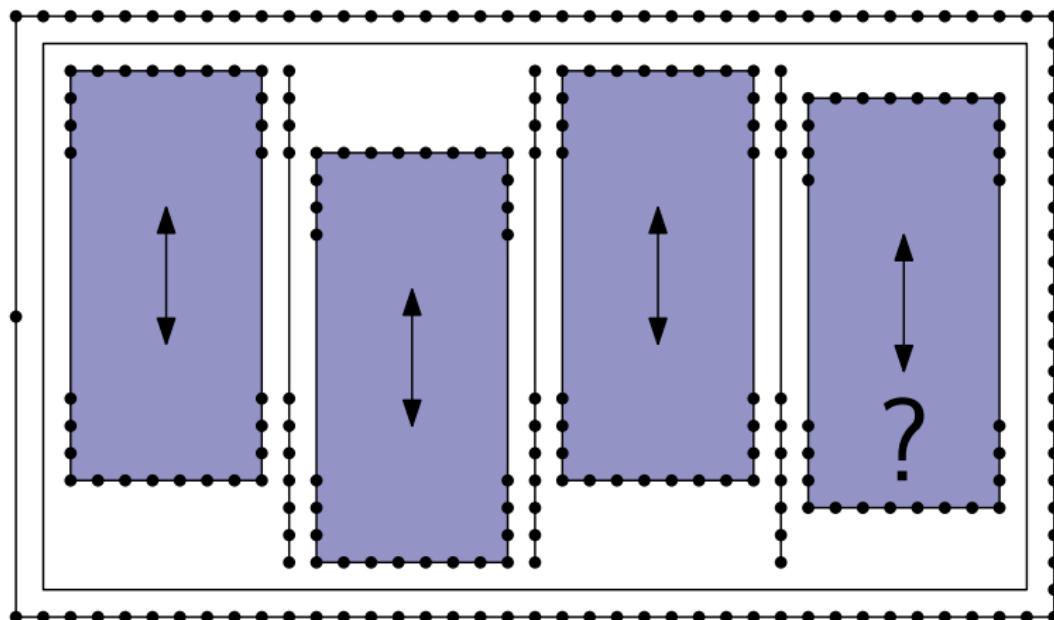
# Begrenzung, Gürtel, Kolbengadget



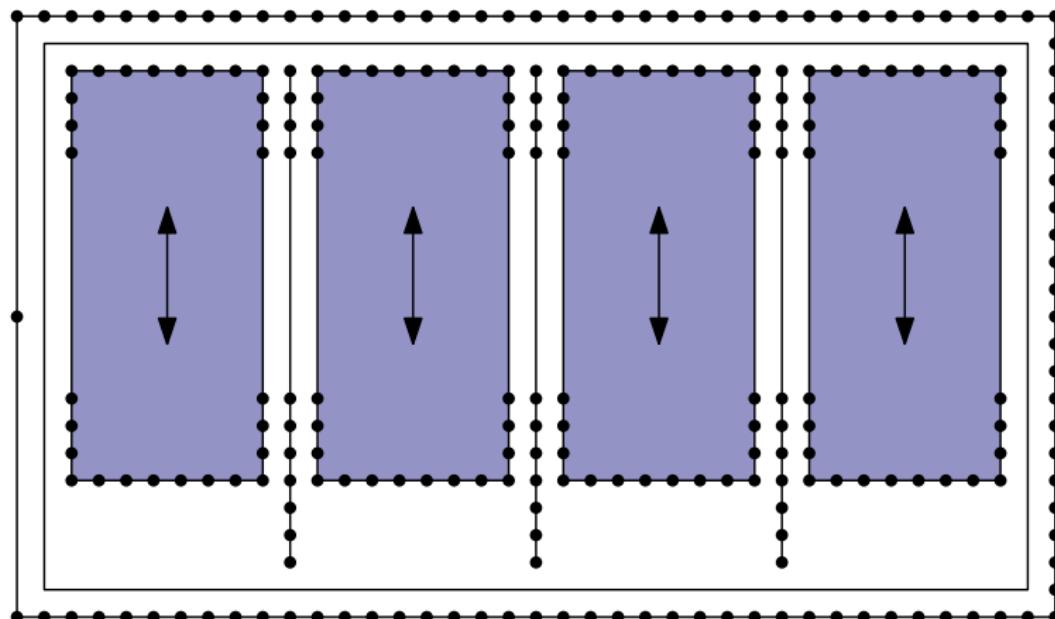
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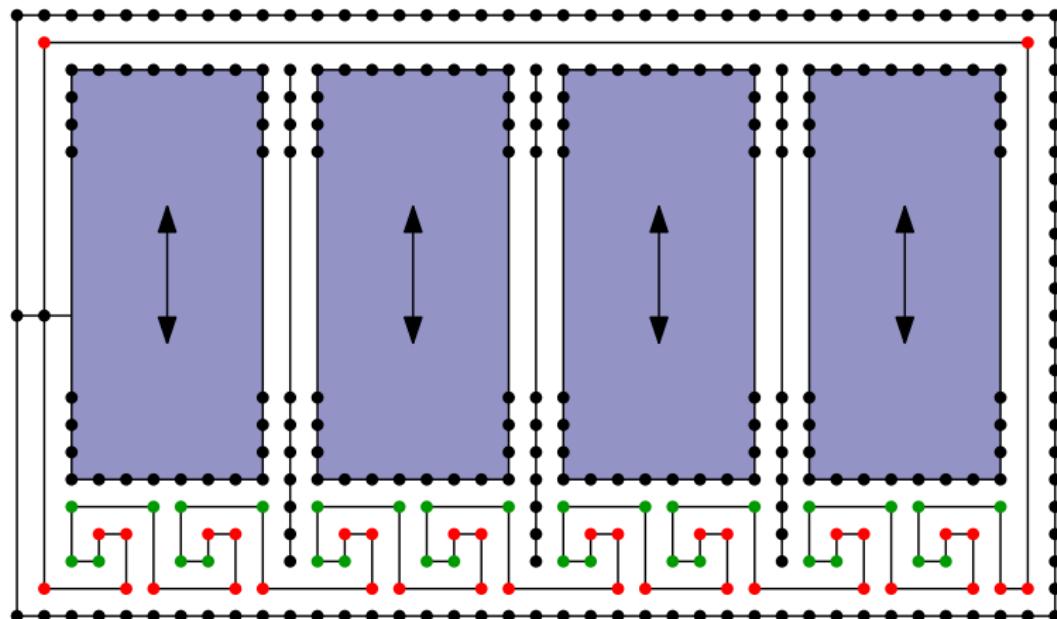
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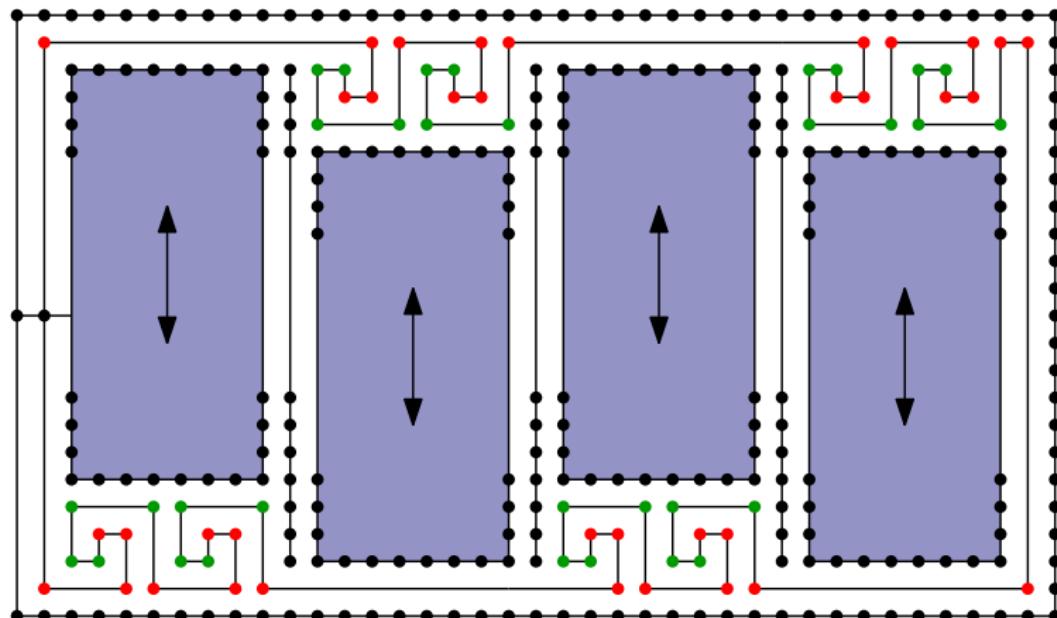
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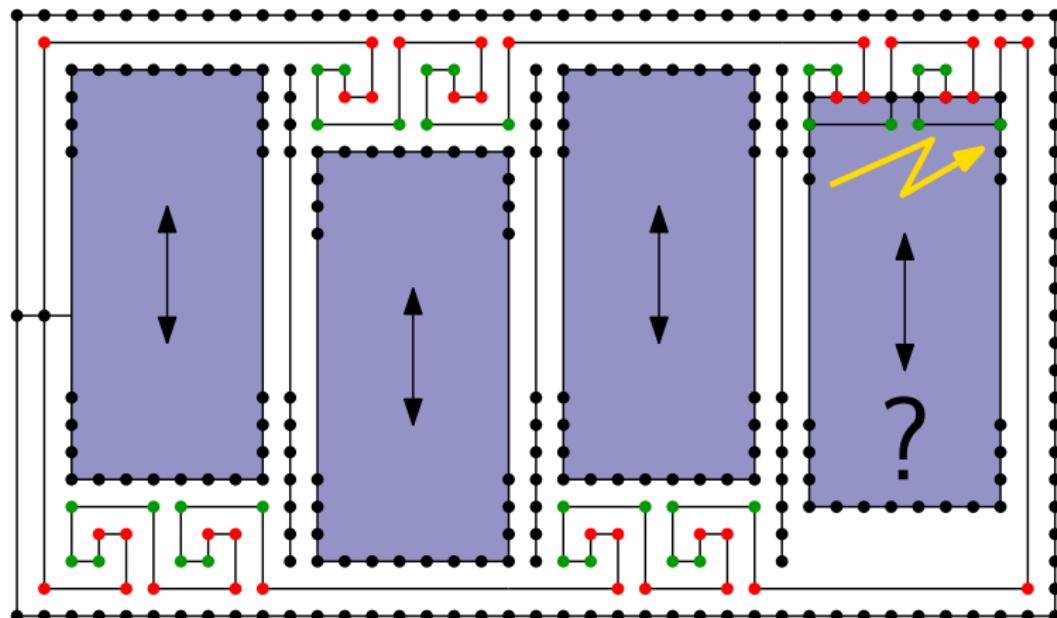
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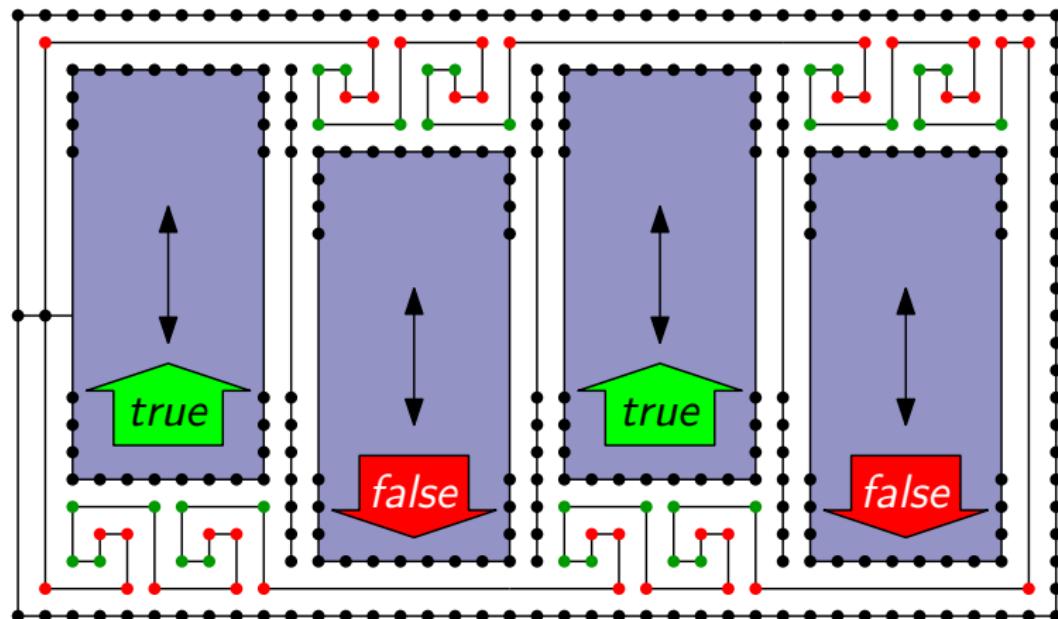
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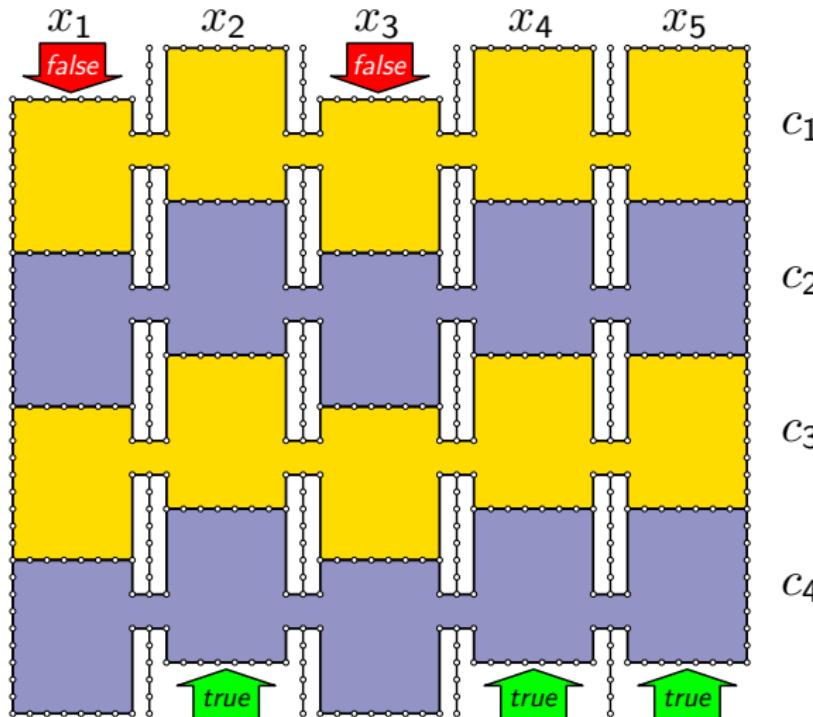
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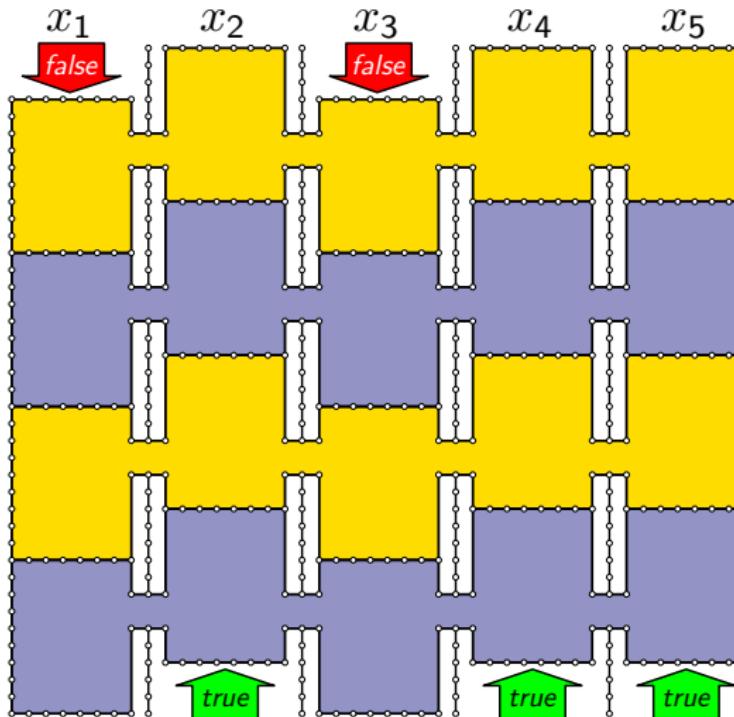
# Begrenzung, Gürtel, Kolbengadget



# Klauselgadgets



# Klauselgadgets



Beispiel:

$$c_1 = x_2 \vee \overline{x_4}$$

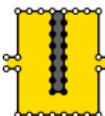
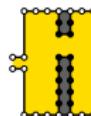
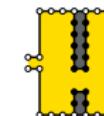
$$c_2 = x_1 \vee x_2 \vee \overline{x_3}$$

$$c_3 = x_5$$

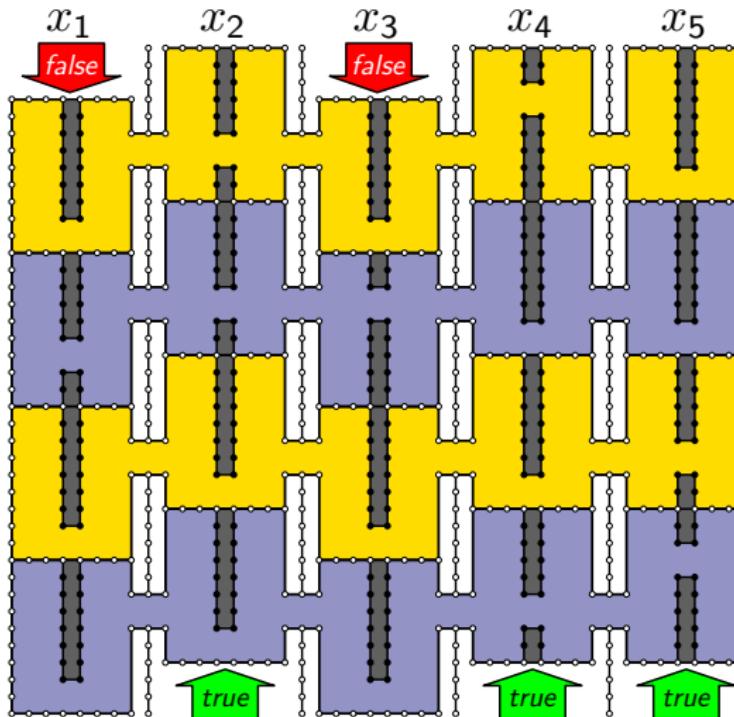
$$c_4 = x_4 \vee \overline{x_5}$$

$$c_3$$

$$c_4$$



# Klauselgadgets



Beispiel:

$$c_1 = x_2 \vee \overline{x_4}$$

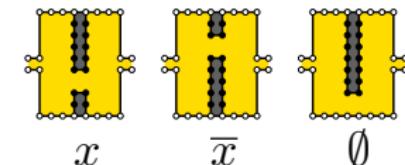
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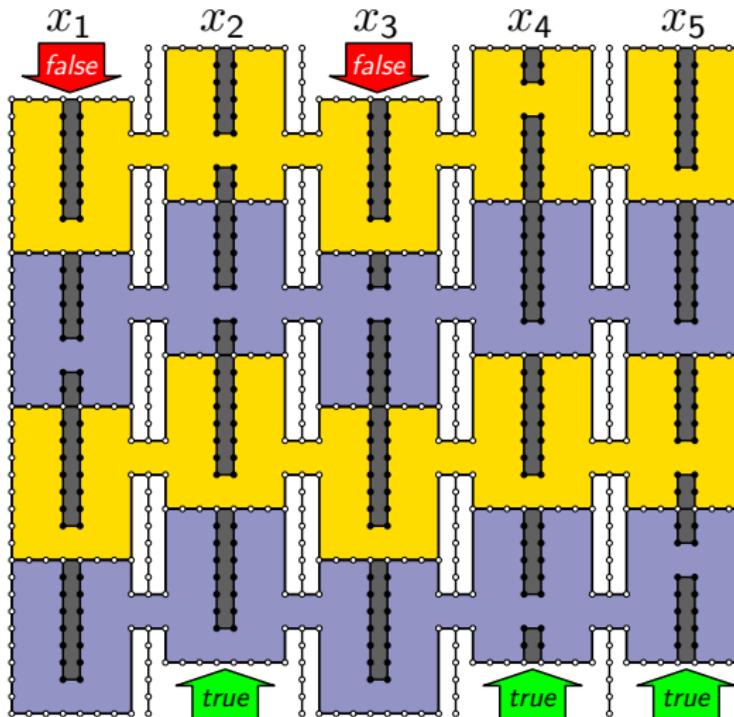
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$$c_4$$



# Klauselgadgets



Beispiel:

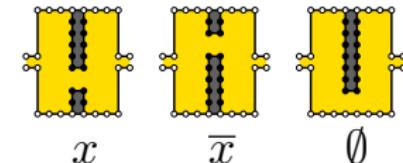
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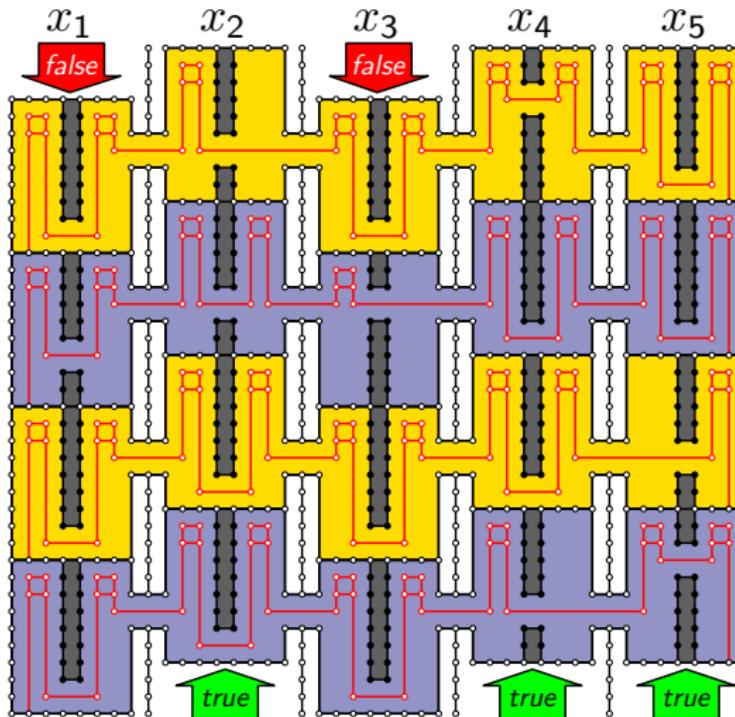
$$c_3$$



$$c_4$$

lege  $(2n - 1)$ -A-Kette  
durch jede Klausel

# Klauselgadgets



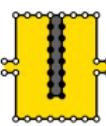
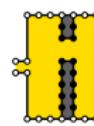
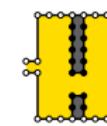
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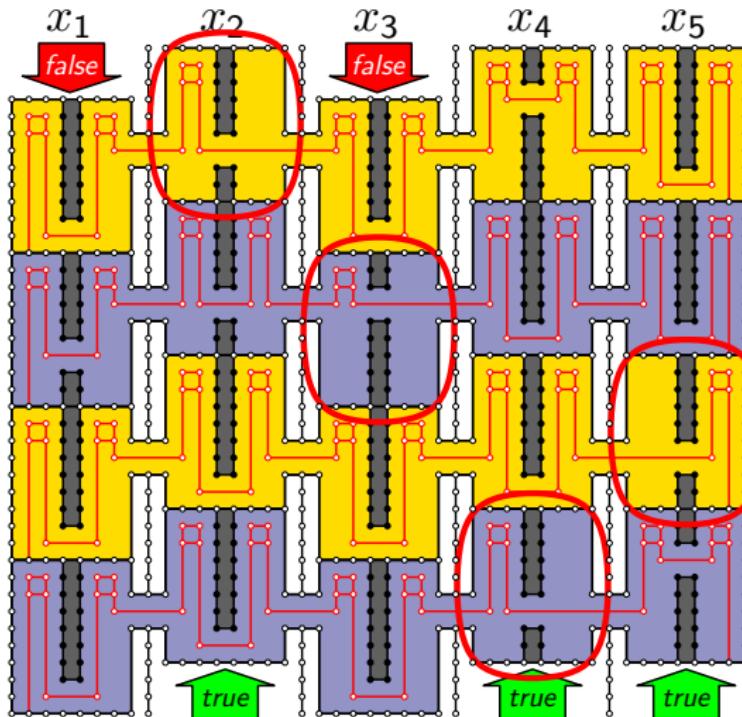


$$c_3$$

$$c_4$$

lege \$(2n - 1)\$-A-Kette  
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# Klauselgadgets



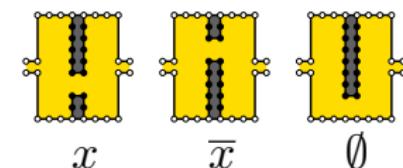
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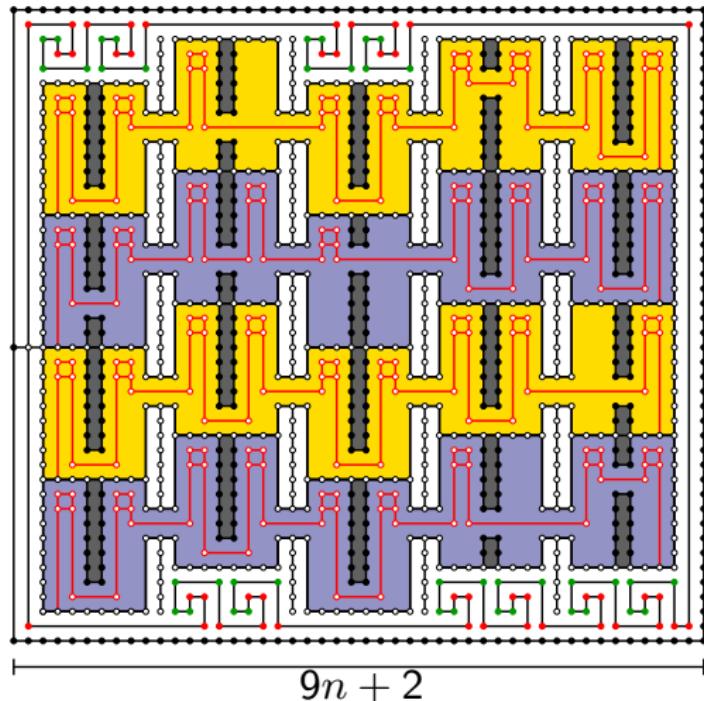
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$c_4$

lege  $(2n - 1)$ -A-Kette  
durch jede Klausel

# Komplette Reduktion



Setze  
 $K = (9n + 2) \cdot (9m + 7)$

$$9m + 7$$

Es gilt:  
 $(G, H)$  auf Fläche  $K$   
zeichnenbar  
 $\Leftrightarrow$   
 $\Phi$  erfüllbar