Cops and Robber: When Capturing is not Surrounding

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2-Players: k Cops 1 Robber



Rules:

- Cops go first.
- Robber is second.
- Moves are between adjacent vertices.



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¹ Cops and Robber – When Capturing is not Surrounding Paul Jungeblut, Samuel Schneider, Torsten Ueckerdt

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Cop number c(G): How many cops are necessary to capture the robber?



vertex surround



vertex surround







Restrictive variants:

 $c_{V,r}(G)$: Robber must not end his move on a cop.

 $c_{E,r}(G)$: Robber must not move through a cop.



c(G):

Quilliot '78 Nowakowski, Winkler '83 Aigner, Fromme '84

0.0 O c(G): Quilliot '78

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 $c_{V,r}(G)$: Burgess et al. '20

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 $c_V(G)$: new

3 $c_{E,r}(G)$: Crytser et al. '20 $\mathbf{c}_{\mathsf{E}}(\mathsf{G})$: new

Our Results

Unification

of the four surrounding variants

Theorem: (informal) Strategy for k cops in some variant gives a strategy for $\leq k \cdot 2\Delta$ cops in any other variant.

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max. degree

Separation

from classical variant

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Theorem:
There are infinitely many G with:
c(G) = 2 and Δ(G) = 3
c_V(G), c_{V,r}(G), c_E(G) and c_{E,r}(G) are unbounded

Upper bounds: Simulation

	not restricted	restricted
vertices	$c_V(G)$	$c_{V,r}(G)$
edges	$c_{E}(G)$	$c_{E,r}(G)$

Upper bounds: Simulation



A restricted robber is a weaker robber.

Upper bounds: Simulation



Replace each cop by a group of $\Delta(G)$ cops.



Upper bounds: Simulation



Replace edge cop by two vertex cops.



Upper bounds: Simulation



Replace vertex cop by a group of $\Delta(G)$ edge cops.



follow original cop

Upper bounds: Simulation



Lower bounds: Constructions

Tight examples for all claimed inequalities:

complete (bipartite) graphs

regular graphs (with "leaves")

 based on "MOLS" (mutually orthogonal Latin squares)

line graphs of complete graphs

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Question: other direction?



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