## Cops and Robber:

## When Capturing is not Surrounding

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## Cops and Robber

2-Players: k Cops
1 Robber

## 승, 숭… 숭

Rules:

## Cops and Robber

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- Cops go first.


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Cop number $\mathbf{c}(\mathbf{G})$ :
How many cops are necessary to capture the robber?

## Surrounding Variants


vertex surround

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$\mathbf{c}_{V}(\mathbf{G})$
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## Restrictive variants:

$\mathbf{c}_{\mathrm{V}, \mathrm{r}}(\mathbf{G})$ :
Robber must not end his move on a cop.
$\mathbf{c}_{\mathrm{E}, \mathrm{r}}(\mathbf{G})$ :
Robber must not move through a cop.

## A Little Bit of History


$c(\mathbf{G}):$
Quilliot '78
Nowakowski, Winkler '83
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$c_{V}(\mathbf{G}):$ new

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## Our Results

## Unification

of the four surrounding variants

Theorem: (informal)
Strategy for $k$ cops in some variant gives a strategy for $\leqslant k \cdot 2 \Delta$ cops in any other variant.
max. degree

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from classical variant
Conjecture: (Crytser et al.)
$c_{\mathrm{E}, \mathrm{r}}(\mathrm{G}) \leqslant \mathrm{c}(\mathrm{G}) \cdot \Delta(\mathrm{G})$

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## Theorem:

There are infinitely many $G$ with:

- $c(G)=2$ and $\Delta(G)=3$
- $c_{V}(G), c_{V, r}(G), c_{E}(G)$ and $c_{\mathrm{E}, \mathrm{r}}(\mathrm{G})$ are unbounded


## Unifying the Surrounding Variants

Upper bounds: Simulation

|  | not <br> not restricted | restricted |
| :---: | :---: | :---: |
| vertices | $\mathbf{c}_{\mathbf{V}}(\mathbf{G})$ | $\mathbf{c}_{\mathbf{V}, \mathbf{r}}(\mathbf{G})$ |
| en <br> edges | $\mathbf{c}_{\mathbf{E}}(\mathbf{G})$ |  |

## Unifying the Surrounding Variants

Upper bounds: Simulation

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A restricted robber is a weaker robber.

## Unifying the Surrounding Variants

Upper bounds: Simulation

|  | not restricted | restricted |
| :---: | :---: | :---: |
| vertices | $c_{V}(\mathbf{G})$ | $\pm \mathrm{c}_{\mathrm{V}, \mathrm{r}}(\mathrm{G})$ |
| edges | $\mathrm{c}_{\mathrm{E}}(\mathrm{G})$ | $\pm c_{E, r}(\mathbf{G})$ |

Replace each cop by a group of $\Delta(\mathrm{G})$ cops.


## Unifying the Surrounding Variants

Upper bounds: Simulation


Replace edge cop by two vertex cops.


## Unifying the Surrounding Variants

Upper bounds: Simulation


Replace vertex cop by a group of $\Delta(\mathrm{G})$ edge cops.

follow original cop

## Unifying the Surrounding Variants

Upper bounds: Simulation


## Lower bounds:

Constructions

Tight examples for all claimed inequalities:

- complete (bipartite) graphs
- regular graphs (with "leaves")
- based on "MOLS"
(mutually orthogonal
Latin squares)
- line graphs of complete graphs


## Separating "Capturing" and "Surrounding"

Observation:
surrounding $\Longrightarrow$ capturing
$\leadsto \mathrm{c}_{\mathrm{X}(\mathrm{r})}(\mathrm{G}) \geqslant \mathrm{c}(\mathrm{G})$
Question: other direction?

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