

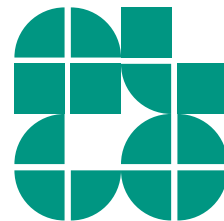
Algorithmen zur Visualisierung von Graphen

Rückblick und gitterbasiertes Graphenzeichnen mit ILP

INSTITUT FÜR THEORETISCHE INFORMATIK · FAKULTÄT FÜR INFORMATIK

Tamara Mchedlidze · **Martin Nöllenburg**

11.02.2015



1. Einführung
2. Baumlayouts
3. Serien-parallele Graphen
4. Geradlinige planare Zeichnungen:
Shift- und Realizer-Methode
5. Datenstrukturen für planare Einbettungen
6. Kontaktrepräsentationen planarer Graphen
7. Lagenlayouts (Sugiyama-Framework)
8. Metro Map Layout
9. Knickminimierung und Kompaktierung in orthogonalen
Layouts (TSM-Framework)
10. Aufwärtsplanarität
11. Kräftebasierte Algorithmen
12. Inkrementelle orthogonale Zeichnungen

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Welche Aspekte haben wir jeweils behandelt? Welche Querbezüge gibt es?

Prüfungsmodalitäten (Master)

Mündliche Prüfung (20 Minuten)

- Termin 1: 25.2.2015 ab 16:15 Uhr
- Termin 2: 25.3.2015 ab 16:15 Uhr
- Termin 3: 15.4.2015 ab 16:15 Uhr
- Anmeldung im Sekretariat (lilian.beckert@kit.edu)

Inhalt: Stoff aus Vorlesung und Übung (Skript, Literatur & Folien/Mitschrieb nutzen – nur was besprochen wurde)

Ziele:

- vorgestellte Algorithmen und Layoutprobleme kennen, erklären und analysieren
- Problemmodellierung/Ästhetikkriterien, Einordnung, Vor-/Nachteile, Beweisideen, ...

Prüfer: beide Dozenten gemeinsam, Tamara: Englisch

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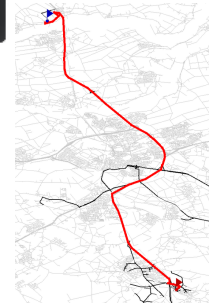
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außer Blatt 6

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Werbung für Sommersemester 2015

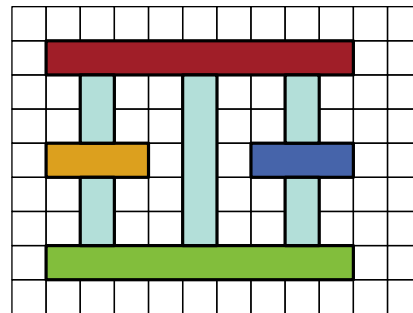
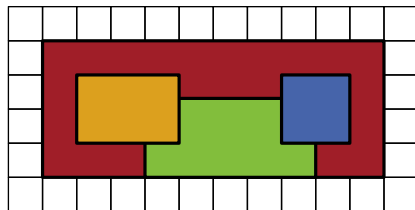
- **Praktikum Graphenvisualisierung (Master)**
aufbauend auf diese VL
Teilnahme am Graph Drawing Contest 2015
- **Algorithmische Kartografie (Master)**
vermutlich dienstags 9:45 Uhr
- **Algorithmen für Routenplanung (Master)**
Montag 14:00 – 15:30
Mittwoch 11:30 – 13:00
- **Masterarbeiten und HiWi-Jobs**
Themen in allen Forschungsfeldern des Lehrstuhls
einfach Mitarbeiter ansprechen



`i11www.iti.kit.edu`

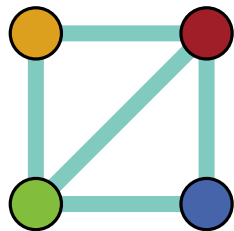
Gitterbasiertes Graphenzeichnen mit ILP

basiert auf [Biedl et al. '14]

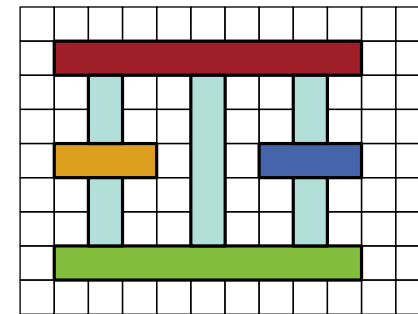
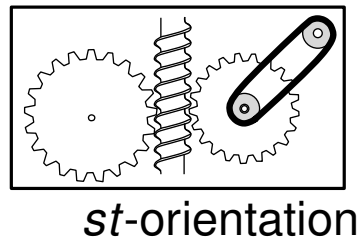
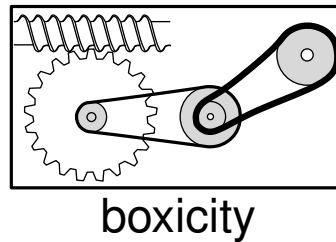
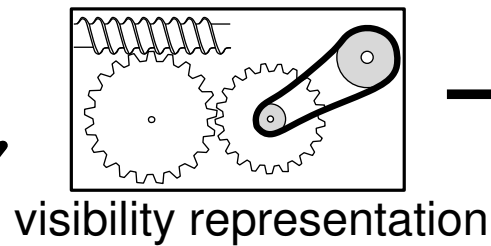


Motivation

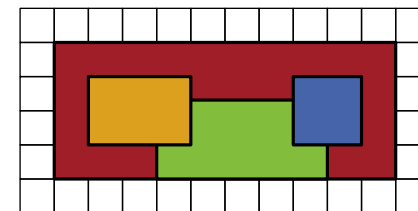
Solve grid-based graph drawing problems in a unified framework:



pathwidth,
bandwidth,
...



minimize
width
NP-hard

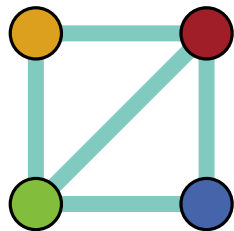


recognition
NP-hard

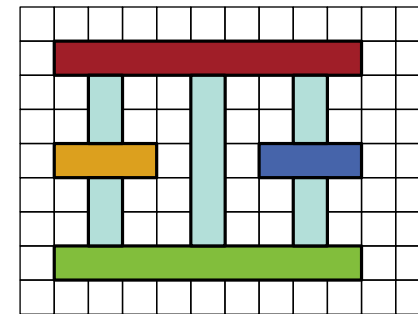
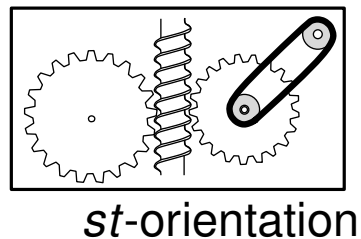
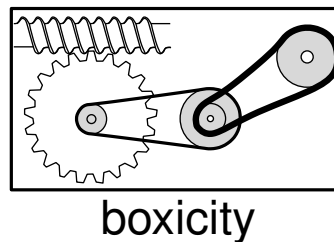
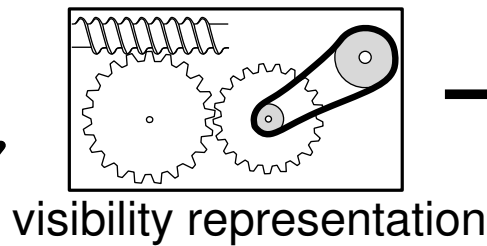


st-orientation
minimize longest *st*-path
NP-hard

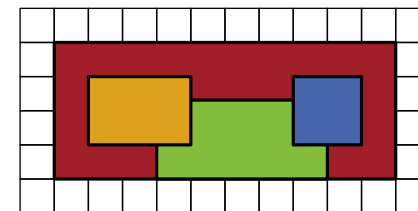
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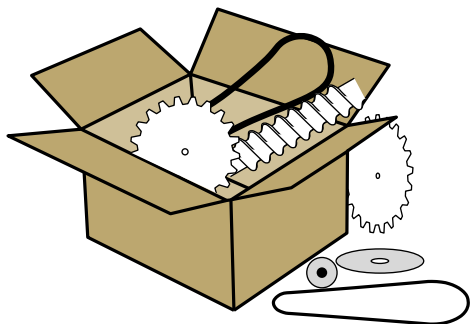
minimize
width
NP-hard



recognition
NP-hard



st-orientation
minimize longest *st*-path
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Apply Integer Linear Programming (ILP):

- Collection of general constraints.
- Solving problem = assemble constraints.

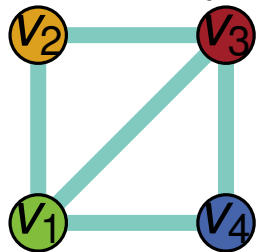
Approach

Given: Grid R , objects, constraints

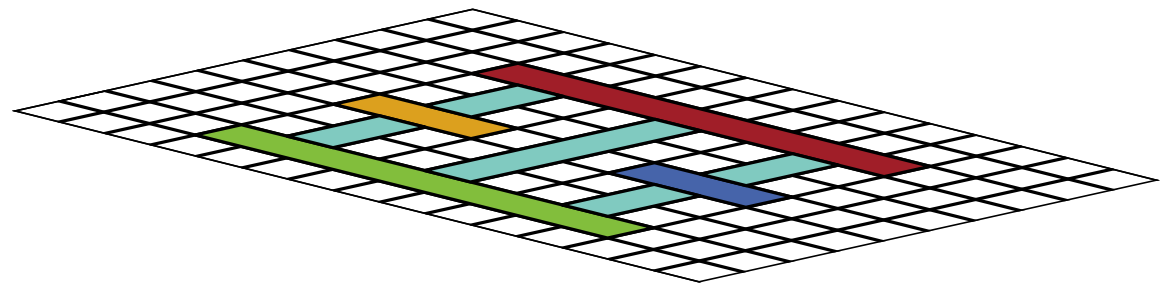
Find: Box on R for each object such that constraints are satisfied.

Example:

Visibility Representation



Vertices = horizontal boxes.
Edges = vertical boxes.



Given: Grid R , objects, constraints

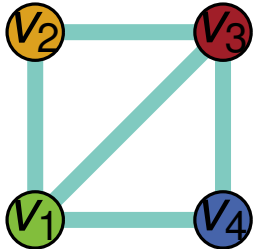
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Representation of single box B :

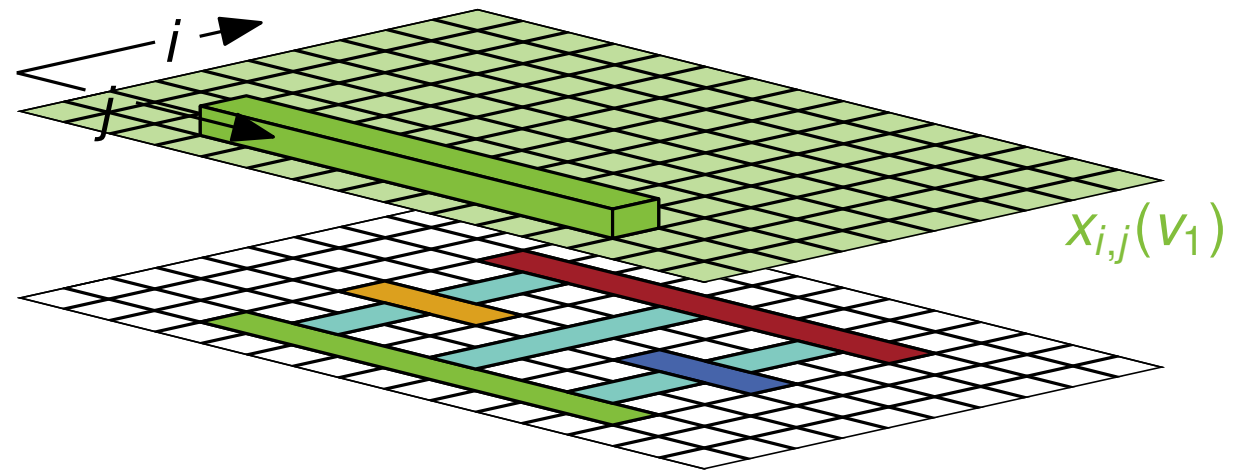
- Grid of binary variables $x(B)_{i,j}$
- Meaning: $x(B)_{i,j} = 1$ iff grid point (i,j) belongs to B .

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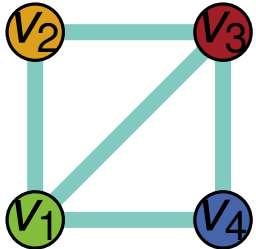
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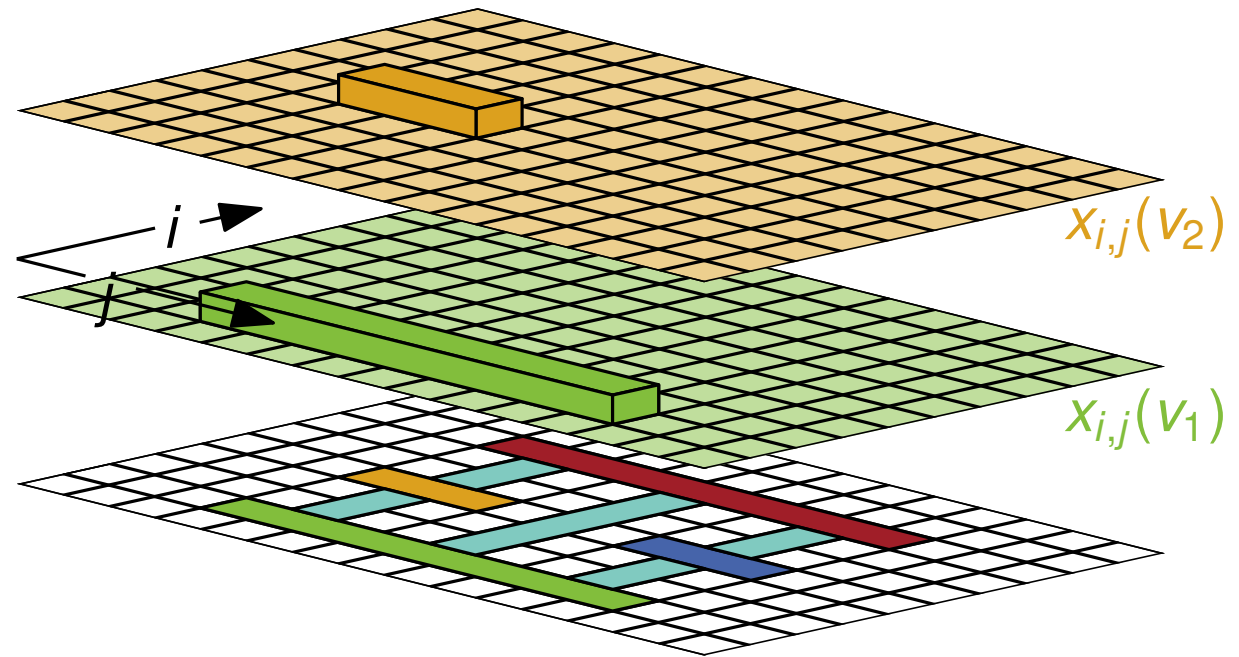
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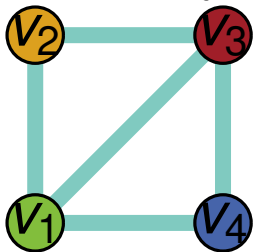
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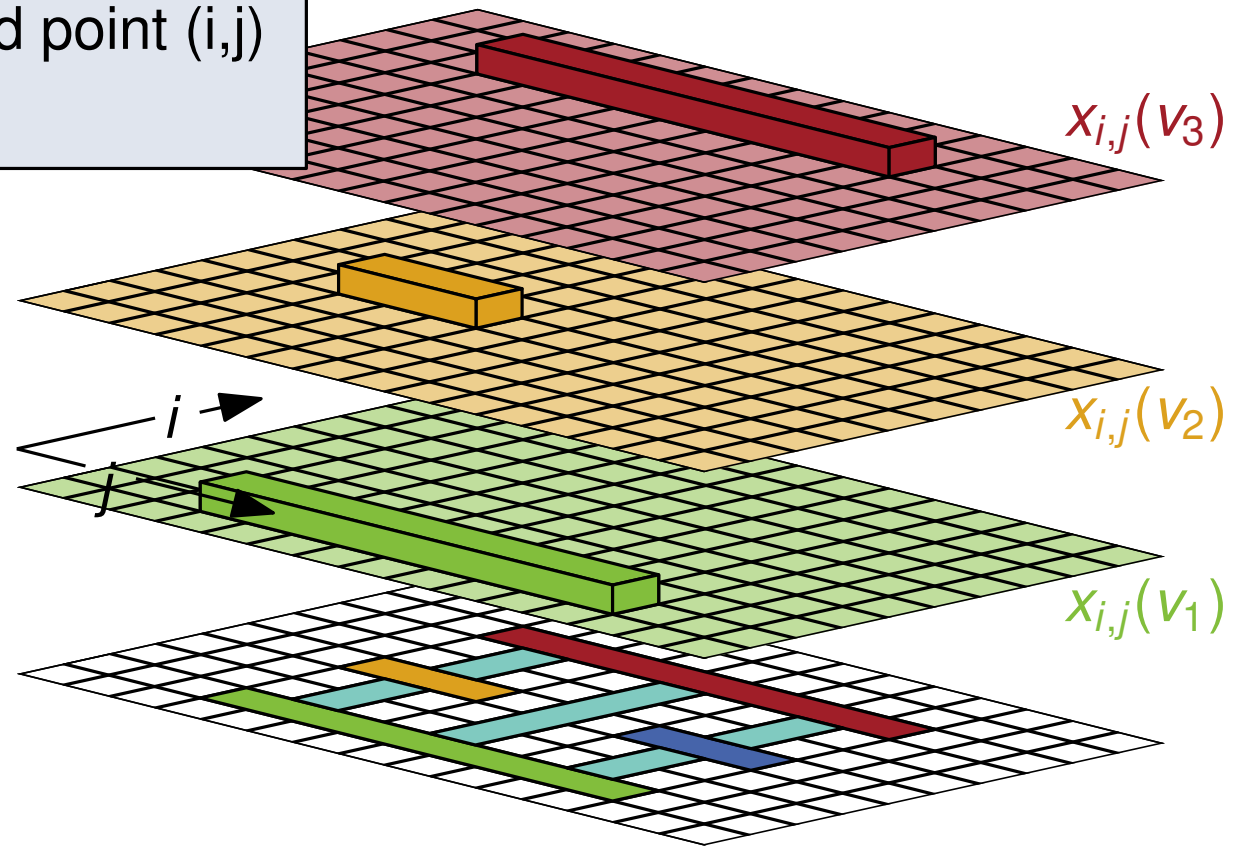
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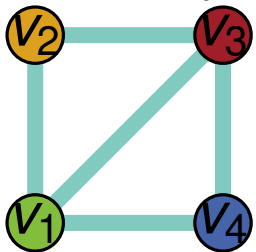
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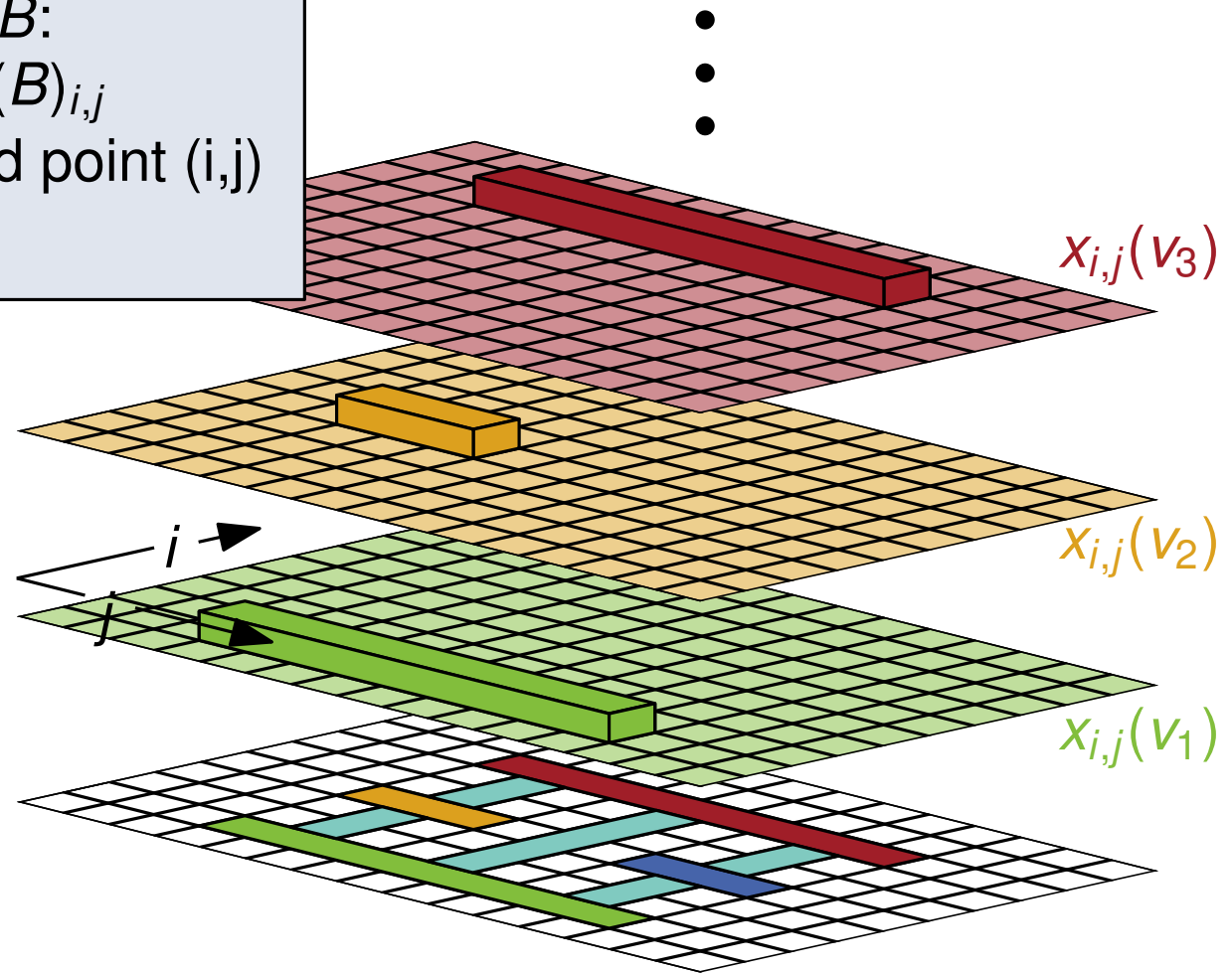
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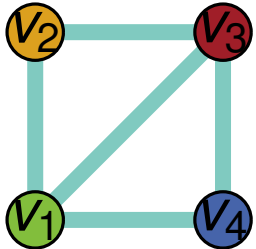
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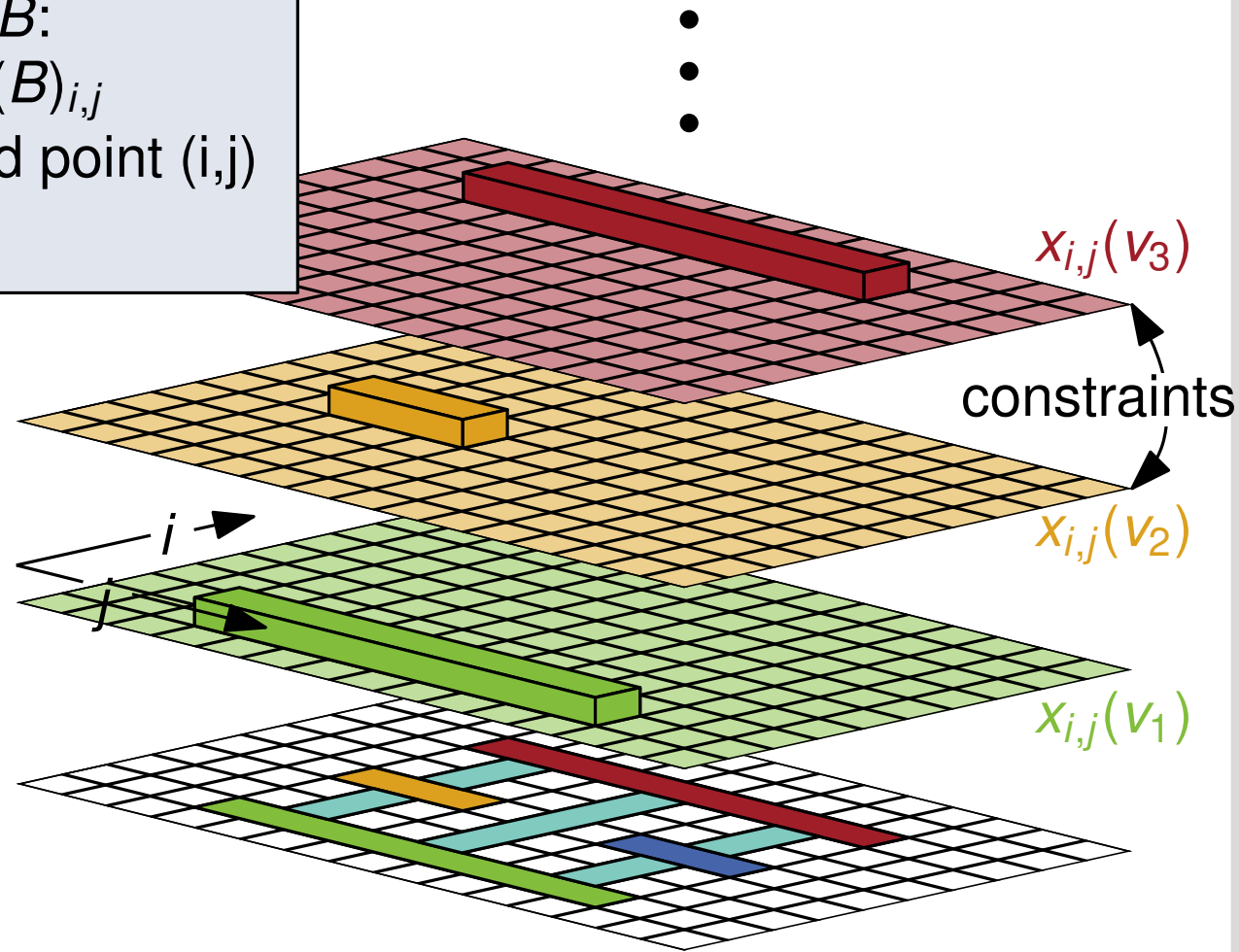
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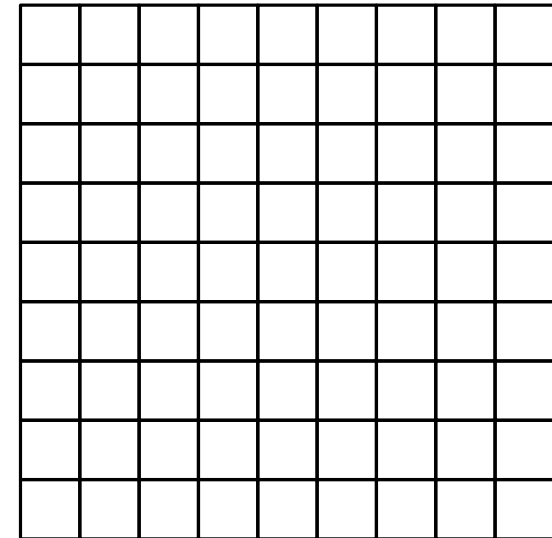
Representing Boxes

Represent box B by ILP constraints (here in 2D).

Grid of binary variables $x_{i,j}$

— j —→

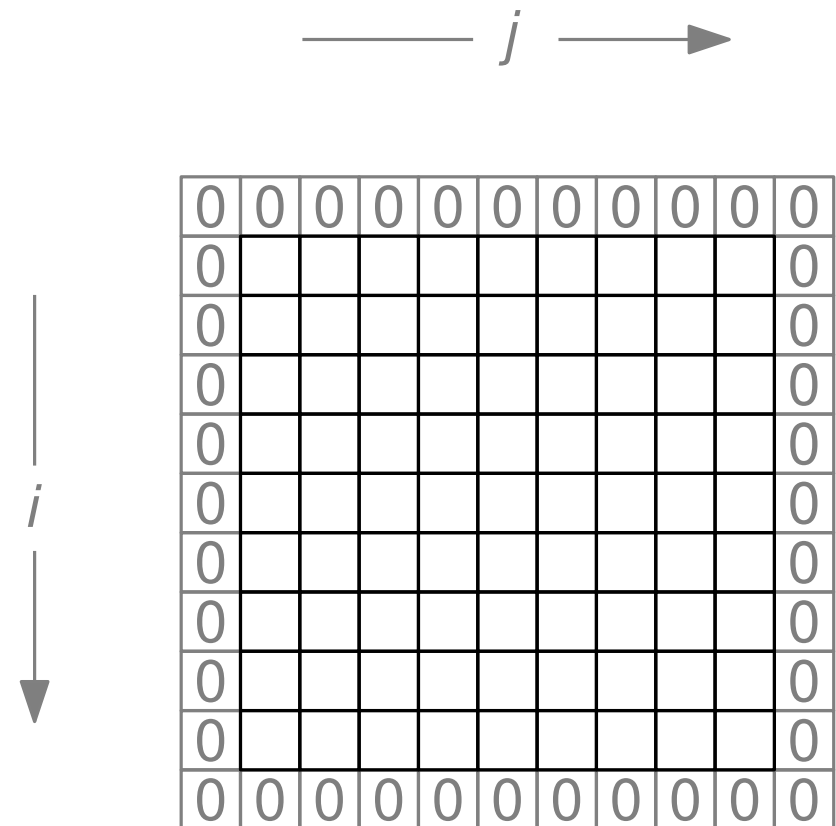
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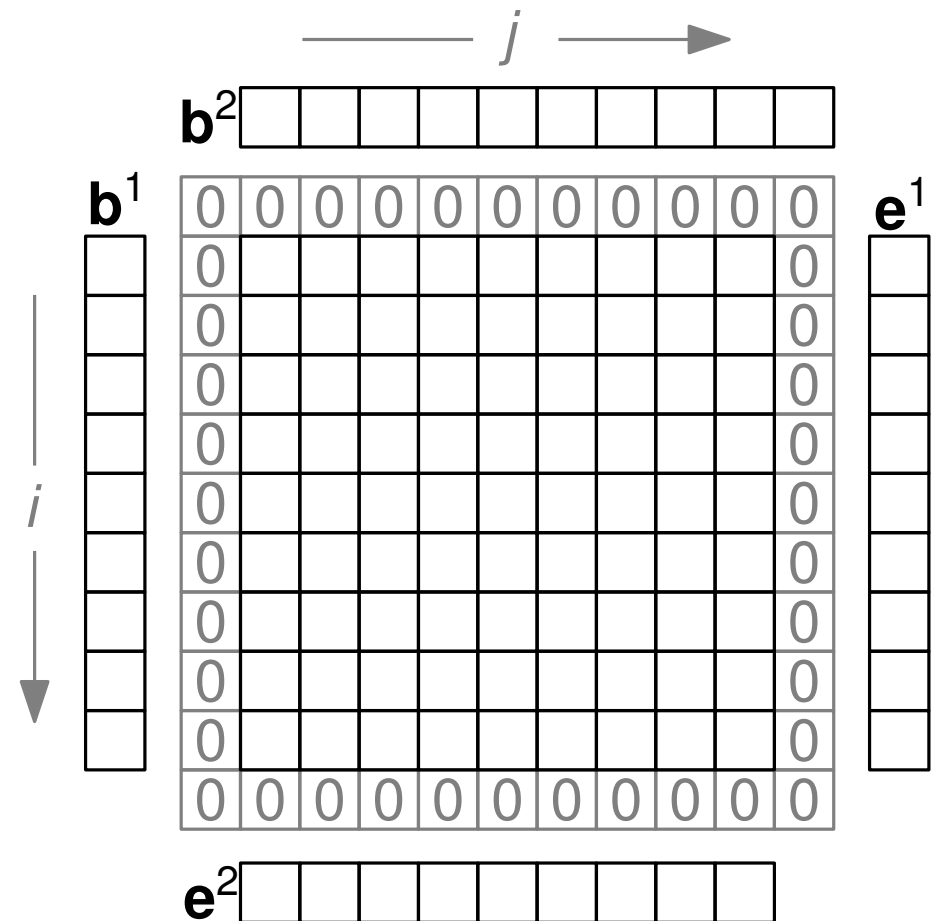
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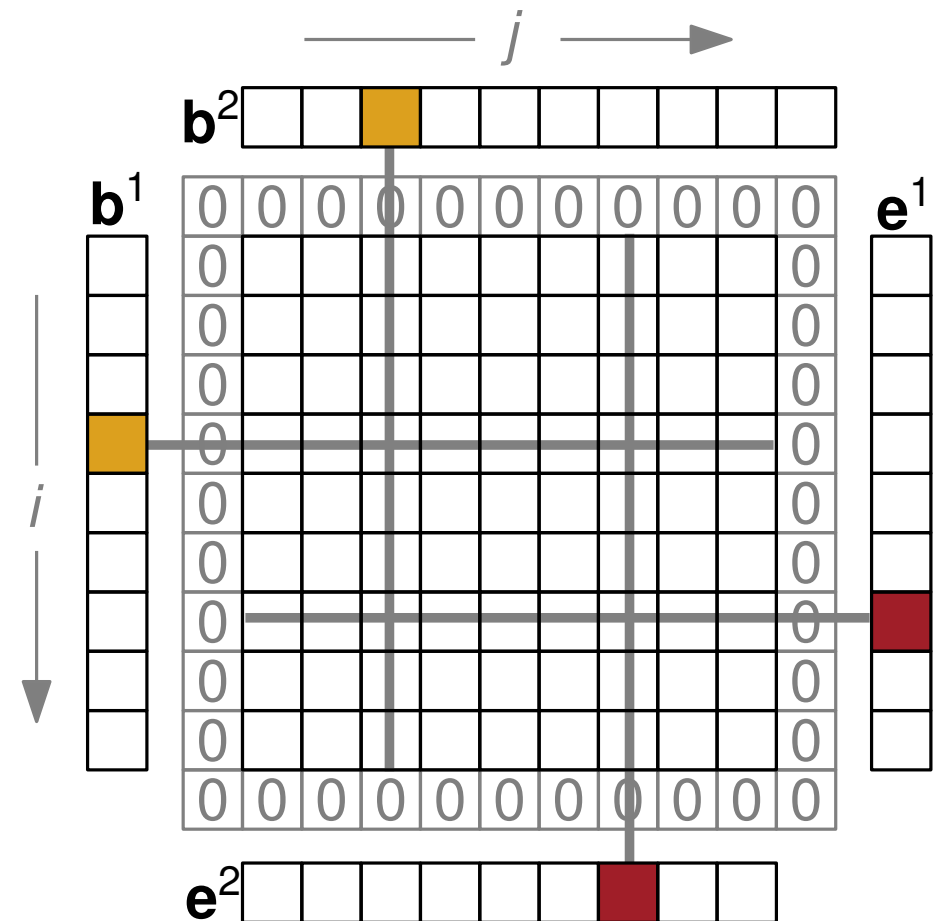
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1.) There is a first/last row/column.

$$\sum_i b_i^d = 1 \quad \sum_i e_i^d = 1 \quad \text{for } d \in \{1, 2\}$$



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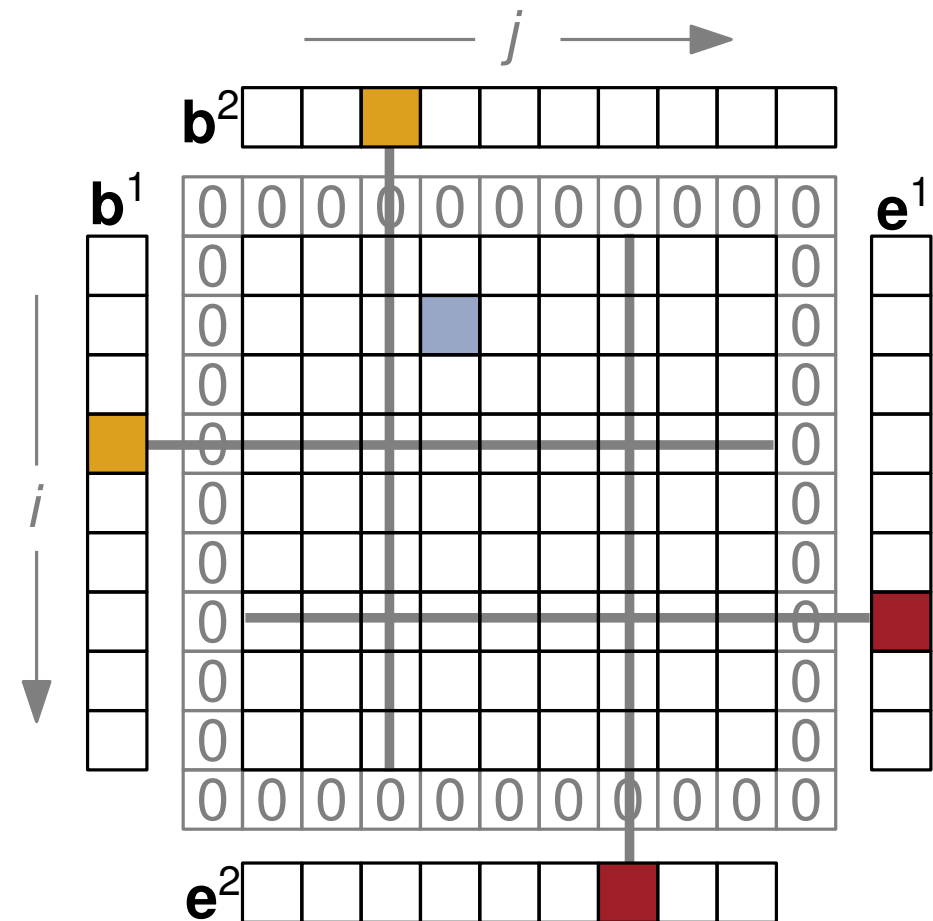
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2.) Boxes are not empty.

$$\sum_{i,j} x_{i,j} \geq 1$$



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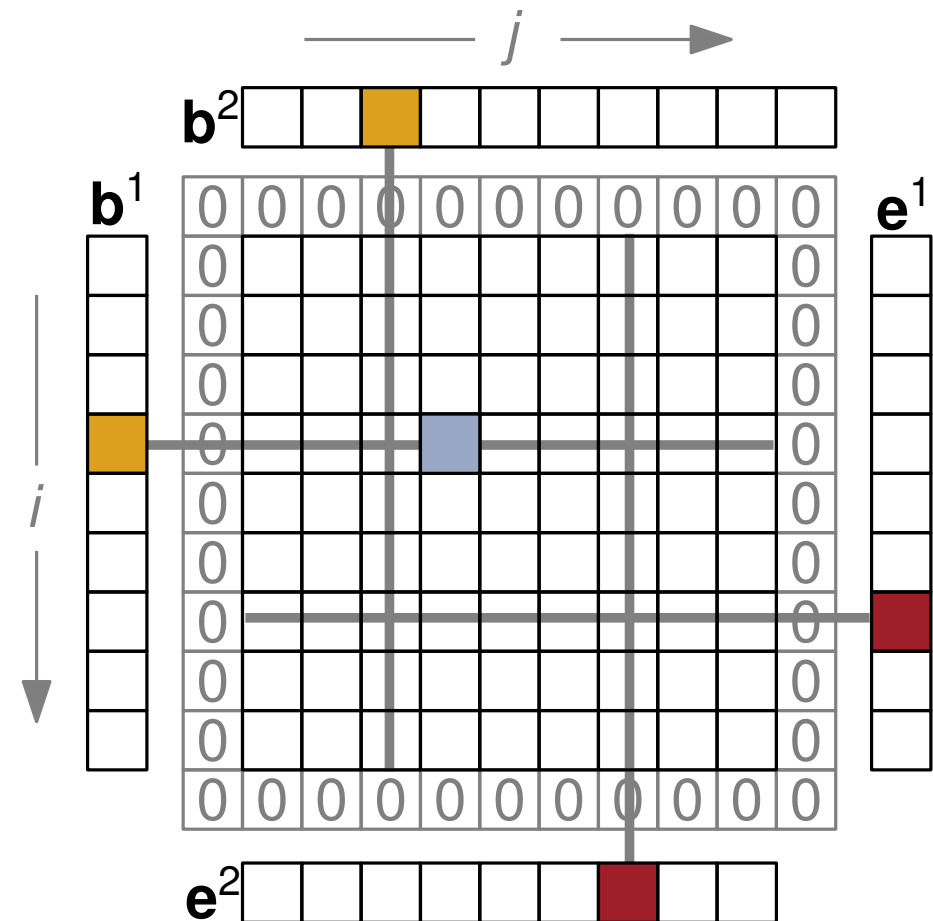
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3.) The first row of B is indicated by b_i^1 .

$$x_{i,j} \leq x_{i-1,j} + b_i^1$$



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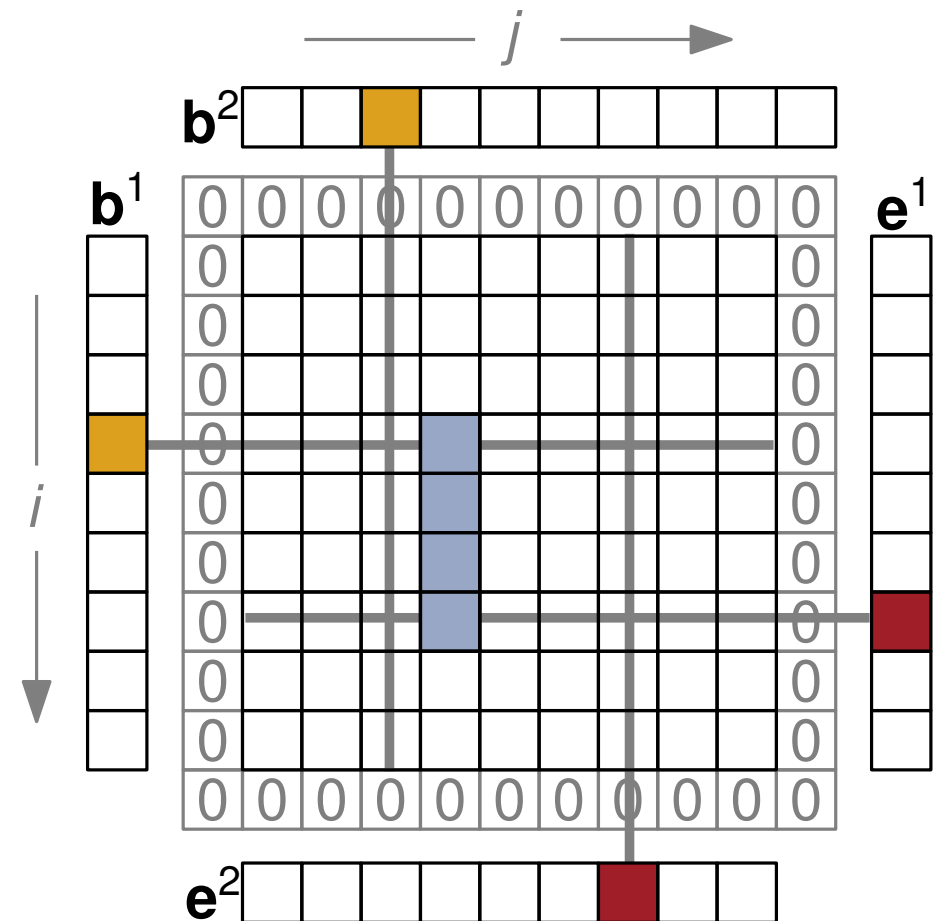
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$$x_{i,j} \leq x_{i-1,j} + b_i^1$$

4.) The last row of B is indicated by e_j^1 .

$$x_{i,j} \leq x_{i+1,j} + e_j^1$$



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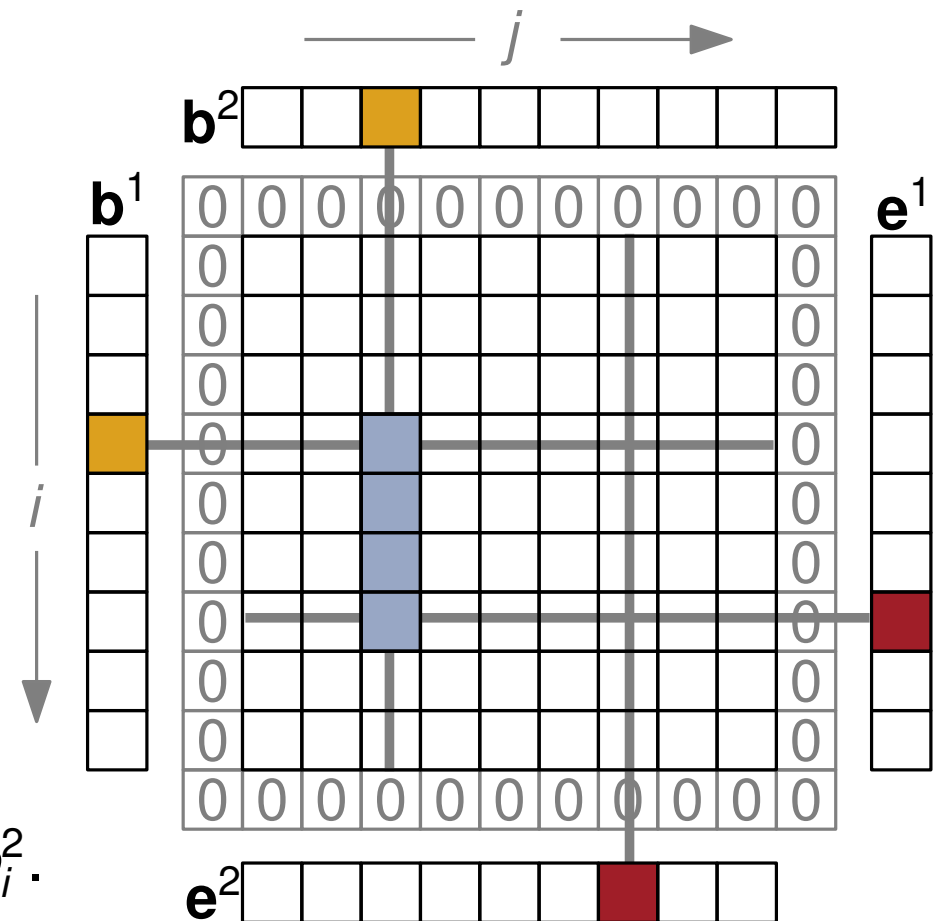
$$x_{i,j} \leq x_{i-1,j} + b_i^1$$

4.) The last row of B is indicated by e_i^1 .

$$x_{i,j} \leq x_{i+1,j} + e_i^1$$

5.) The first column of B is indicated by b_j^2 .

$$x_{i,j} \leq x_{i,j-1} + b_j^2$$



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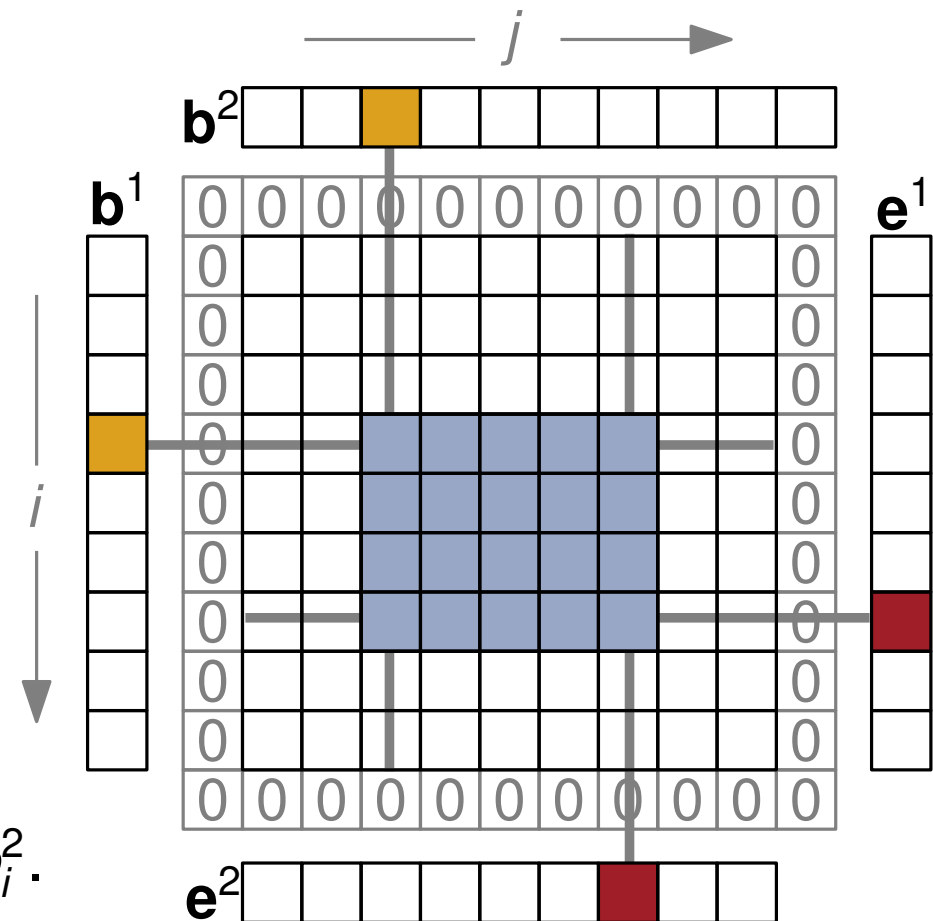
$$x_{i,j} \leq x_{i+1,j} + e_i^1$$

5.) The first column of B is indicated by b_j^2 .

$$x_{i,j} \leq x_{i,j-1} + b_j^2$$

6.) The last column of B is indicated by e_j^2 .

$$x_{i,j} \leq x_{i,j+1} + e_j^2$$



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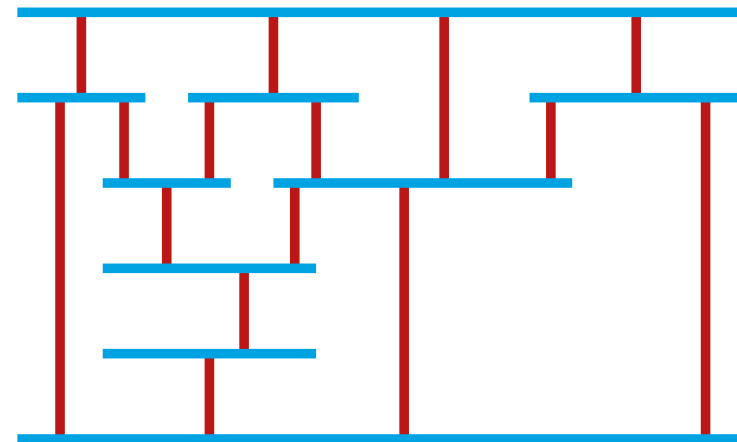
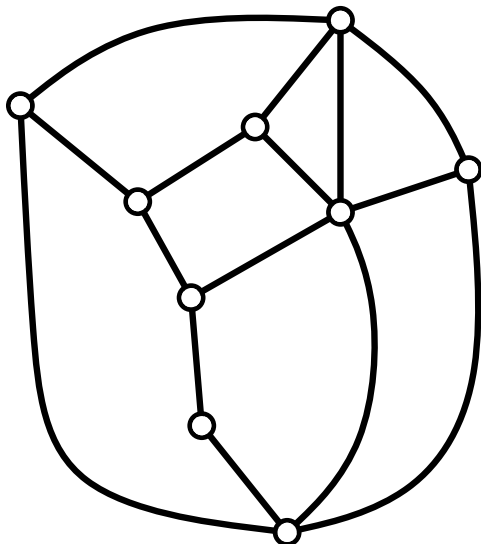
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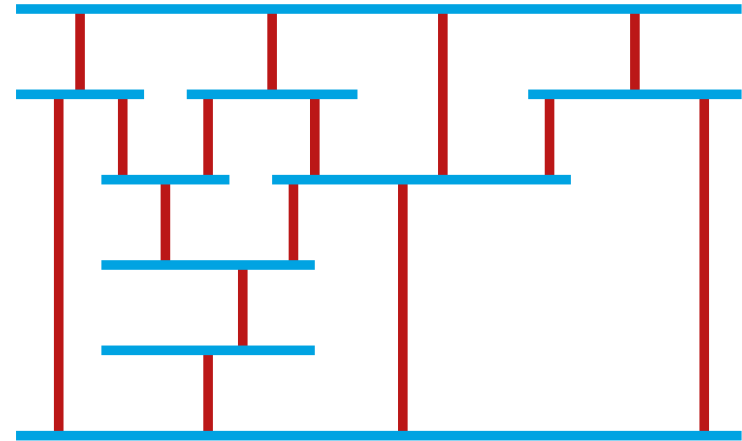
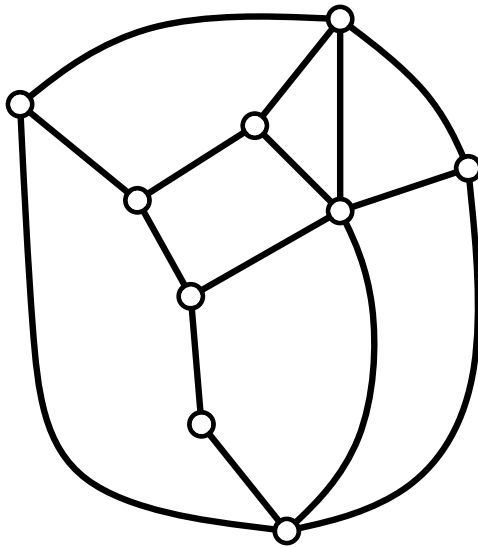
Aufgaben:

- modelliere $\text{Höhe}(B) = 1$
- modelliere $\text{Höhe}(B) = k$
- modelliere $B_1 \cap B_2 = \emptyset$
- modelliere $B_1 \subset B_2$

Def: A **visibility representation** of a graph $G = (V, E)$ draws

- each vertex v as a horizontal line segment $\Gamma(v)$
 - each edge $e = (u, v)$ as a vertical line segment $\Gamma(e)$
- such that
- no two vertex segments intersect
 - no two edge segments intersect
 - each edge segment $\Gamma(u, v)$ has its endpoints on $\Gamma(u)$ and $\Gamma(v)$ and intersects no other vertex segment

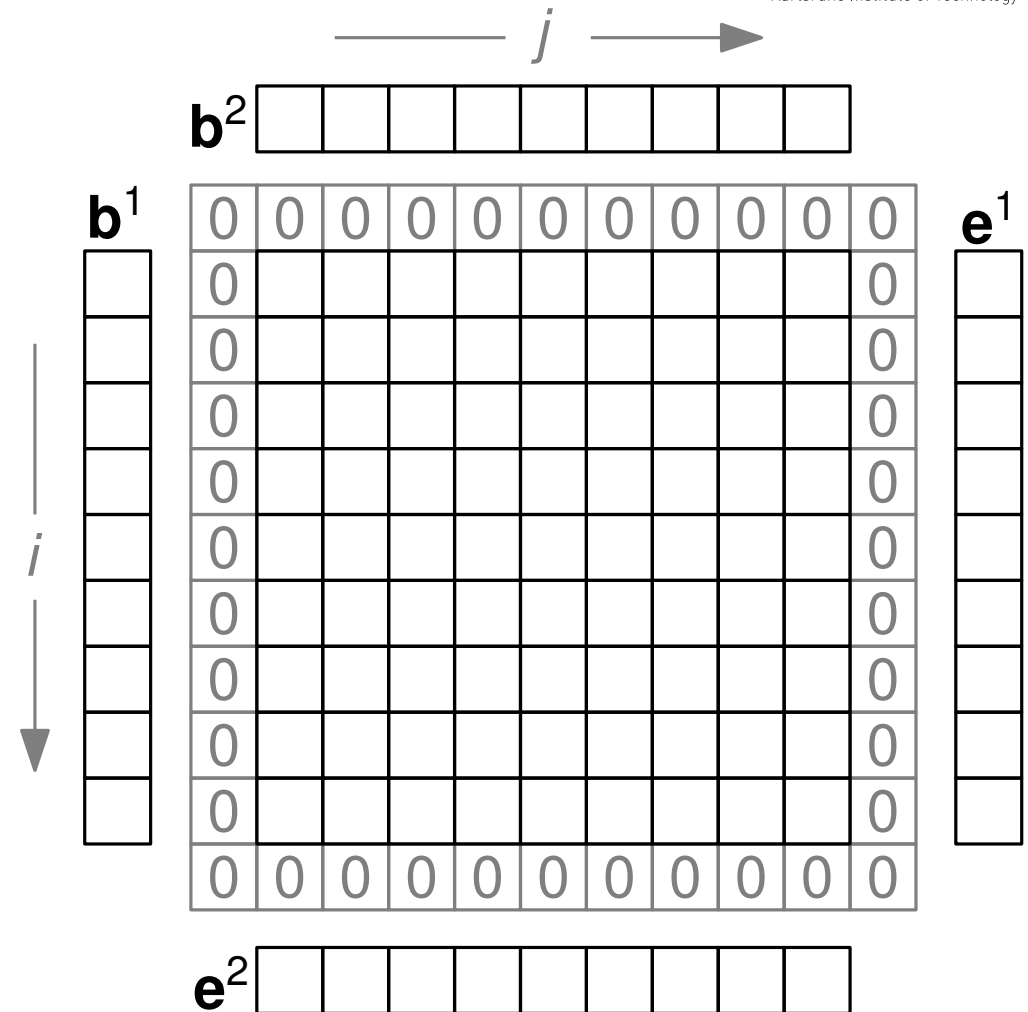
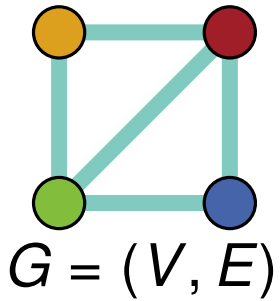




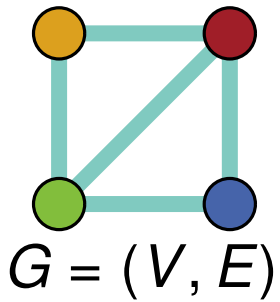
- graph G must be planar (obviously)
- every planar graph has a visibility representation
[Wismath '85], [Tamassia, Tollis '86]
- minimizing the area of a visibility representation is NP-hard

[Lin, Eades '03]

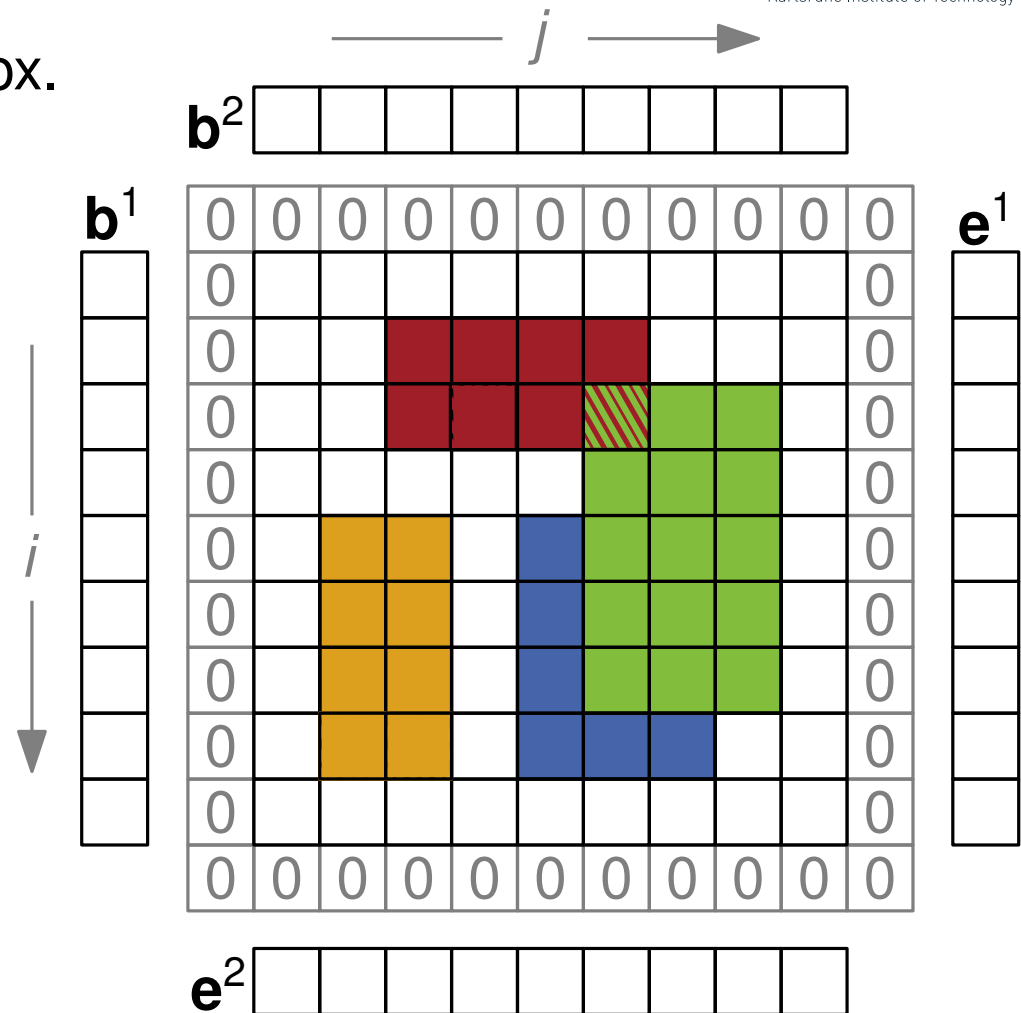
Modeling Visibility Representation



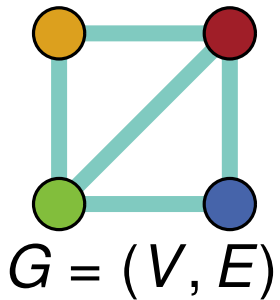
Modeling Visibility Representation



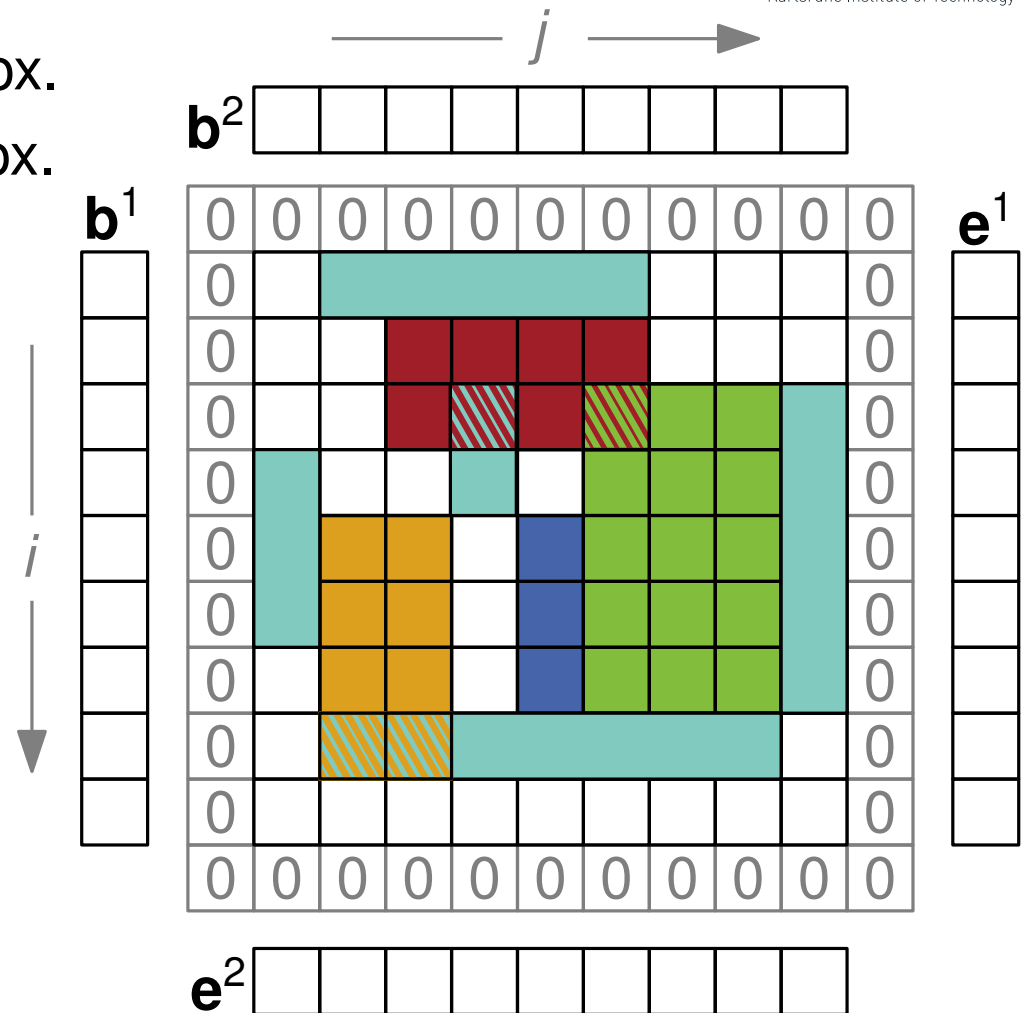
1.) $\forall v \in V$ introduce box.



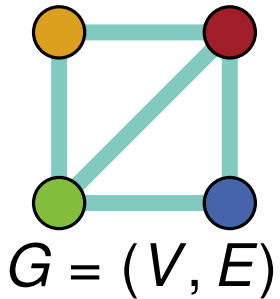
Modeling Visibility Representation



- 1.) $\forall v \in V$ introduce box.
- 2.) $\forall e \in E$ introduce box.

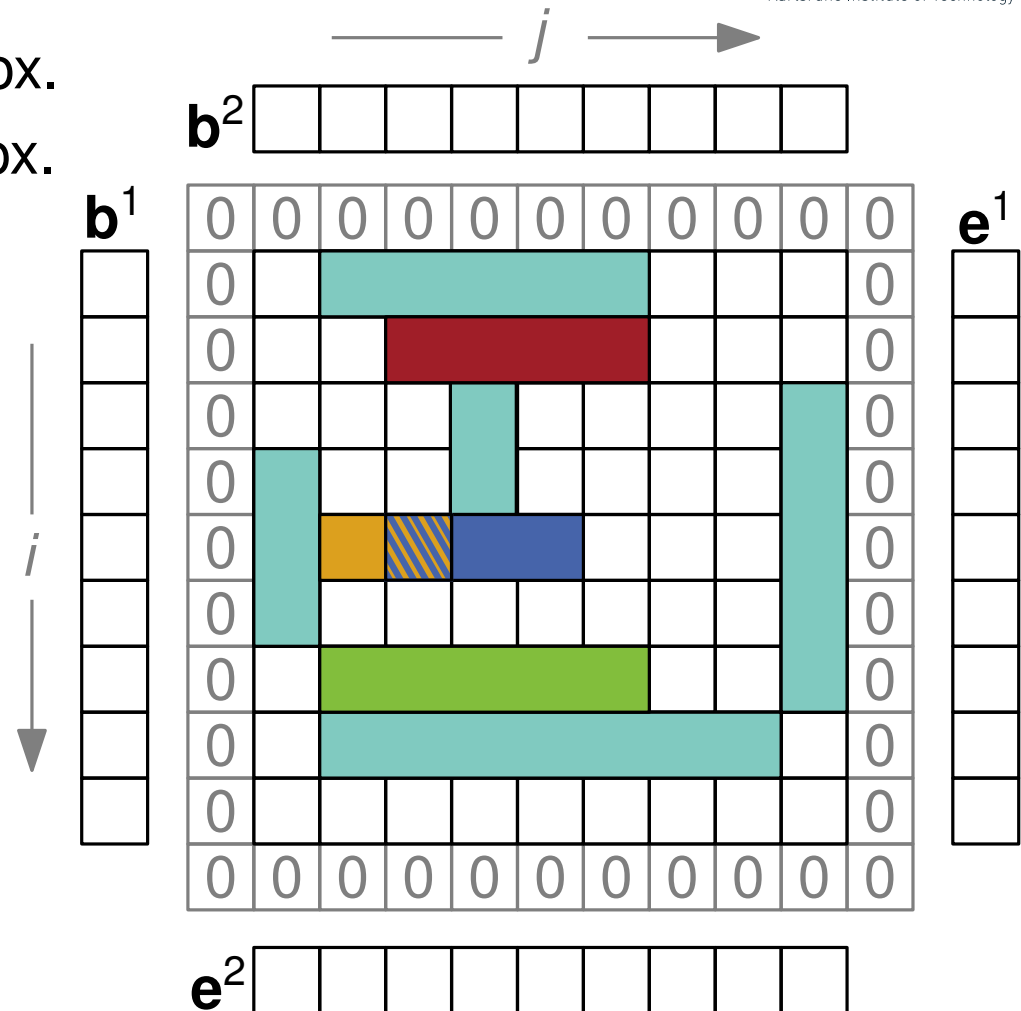


Modeling Visibility Representation

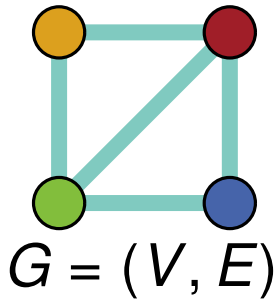


- 1.) $\forall v \in V$ introduce box.
- 2.) $\forall e \in E$ introduce box.

3.) Vertices are horizontal segments:

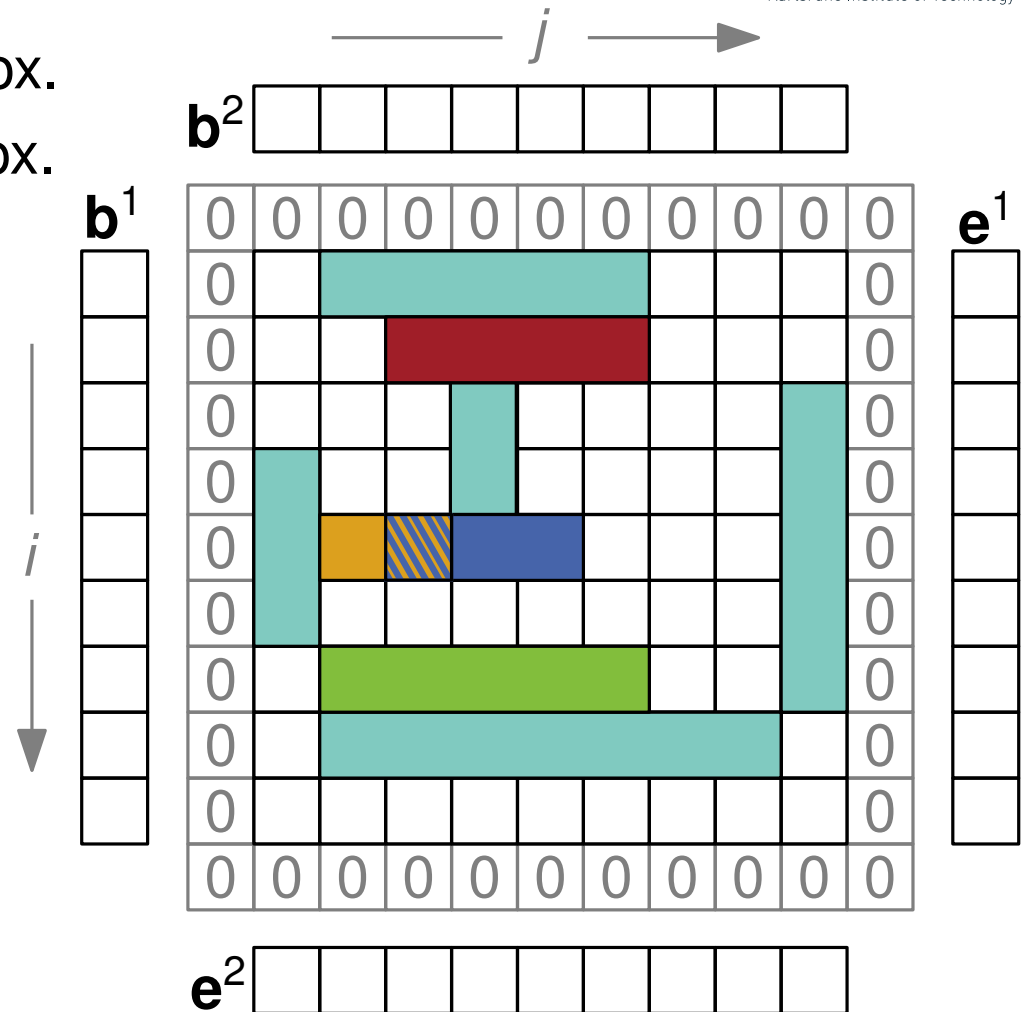


Modeling Visibility Representation

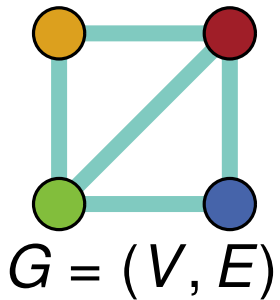


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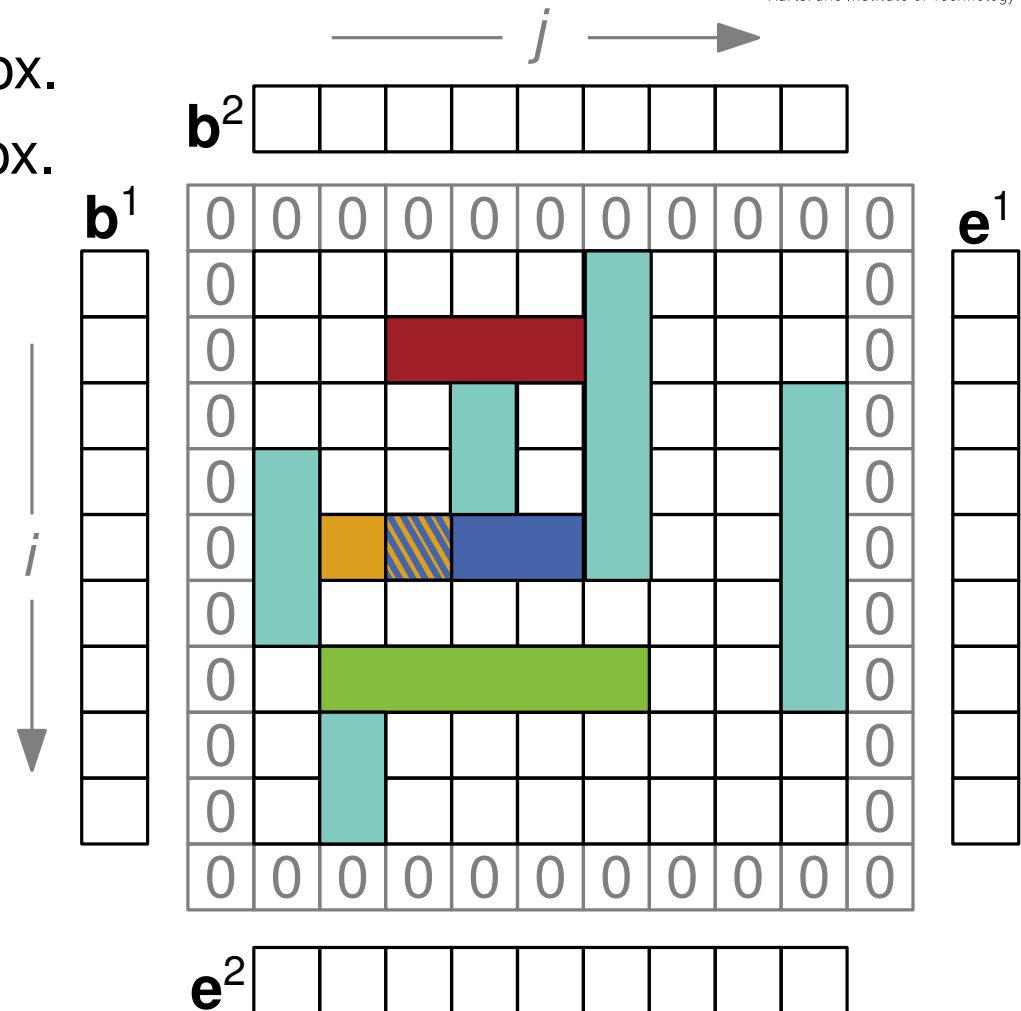


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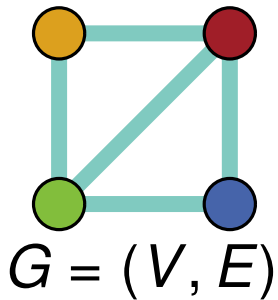
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Modeling Visibility Representation



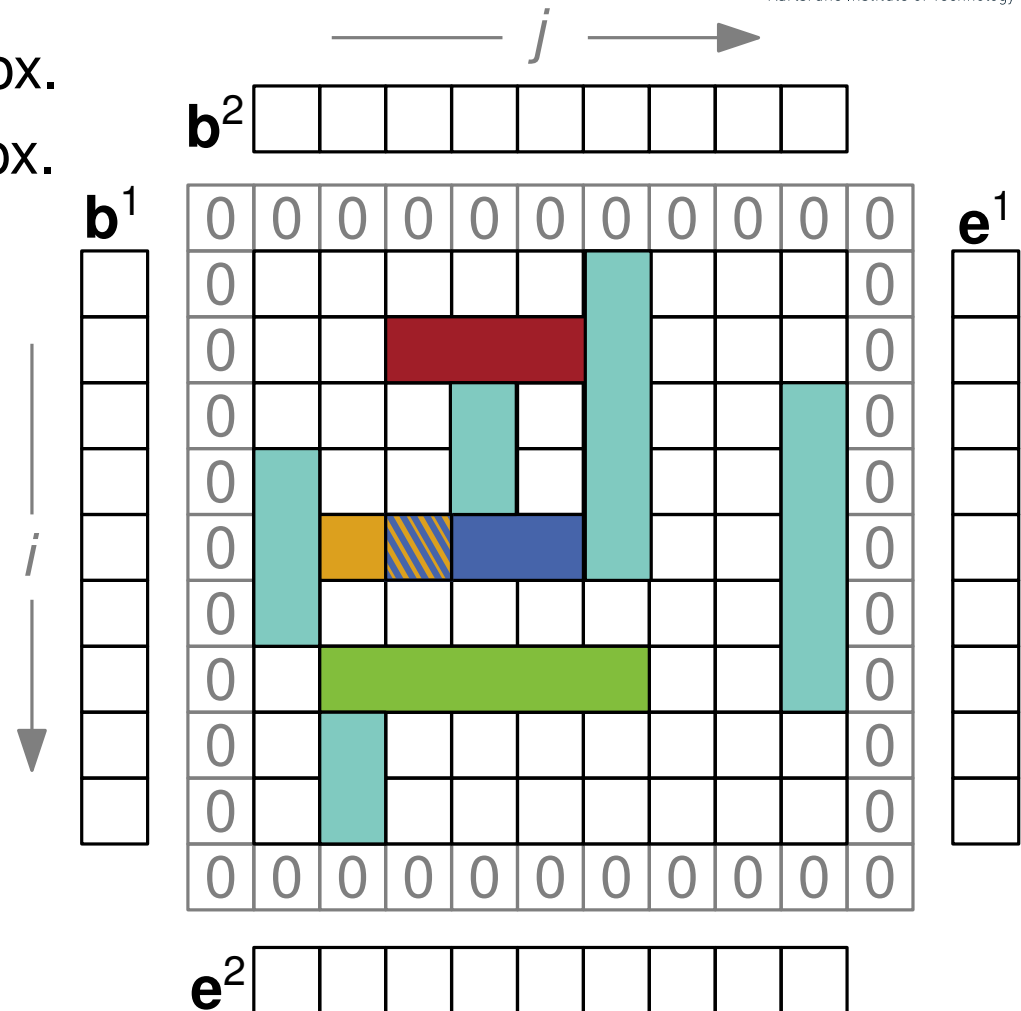
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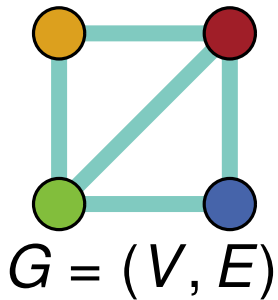
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Modeling Visibility Representation



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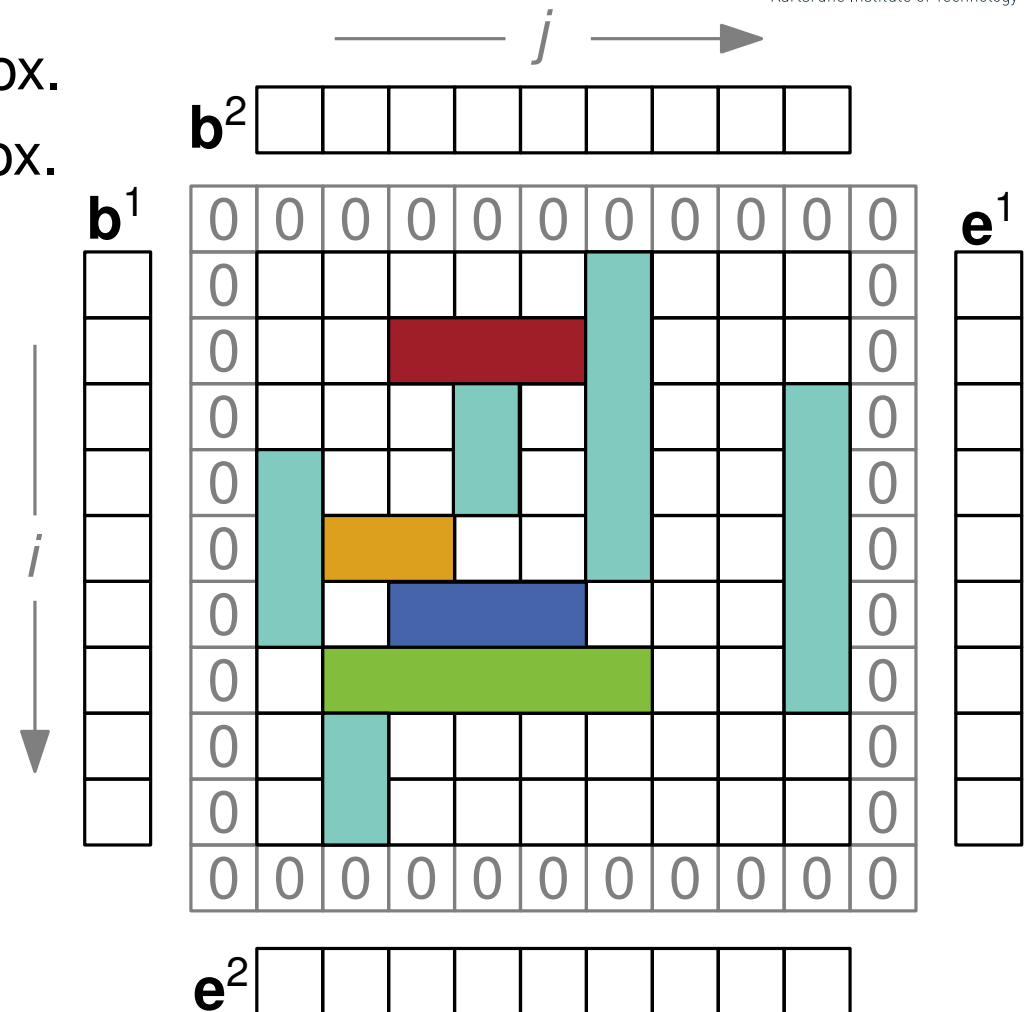
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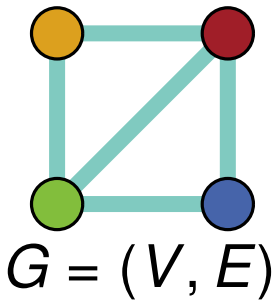
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Modeling Visibility Representation



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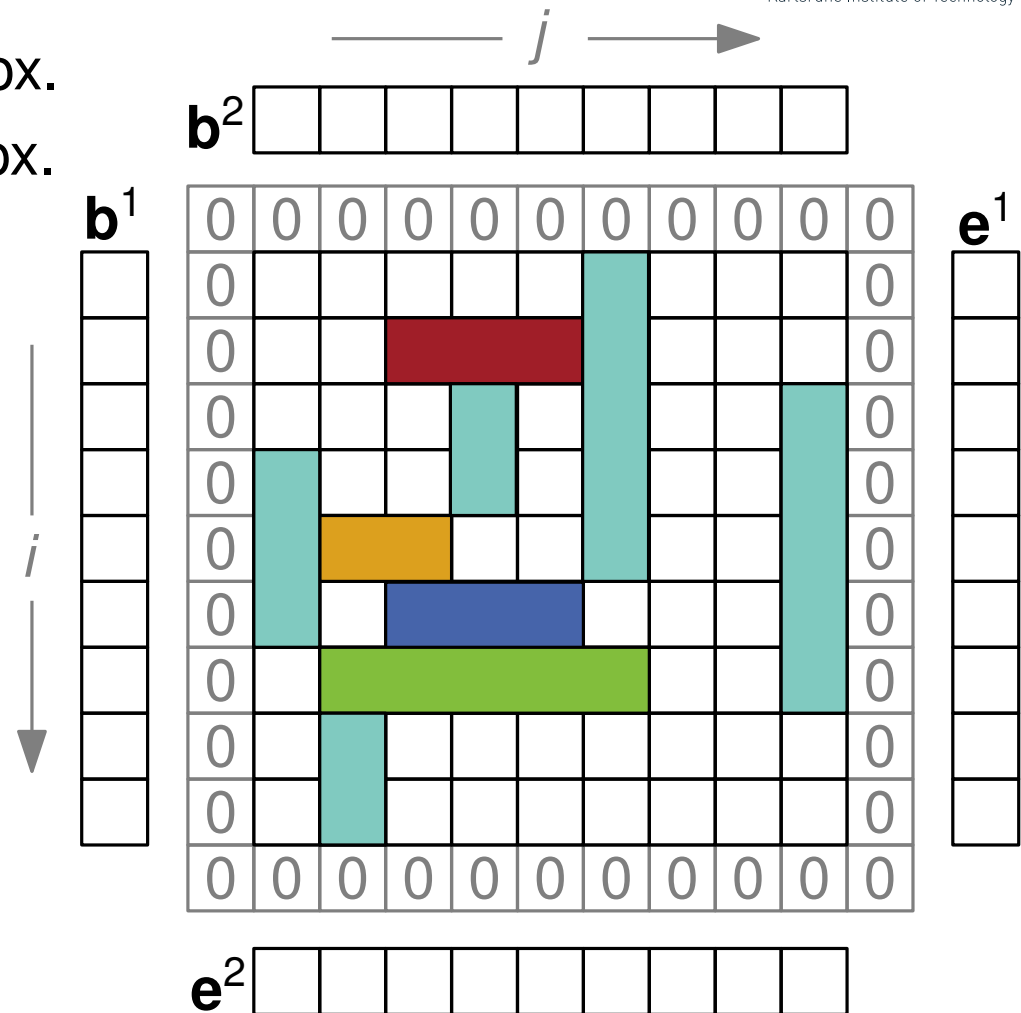
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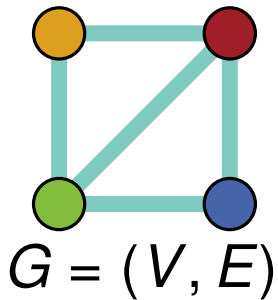
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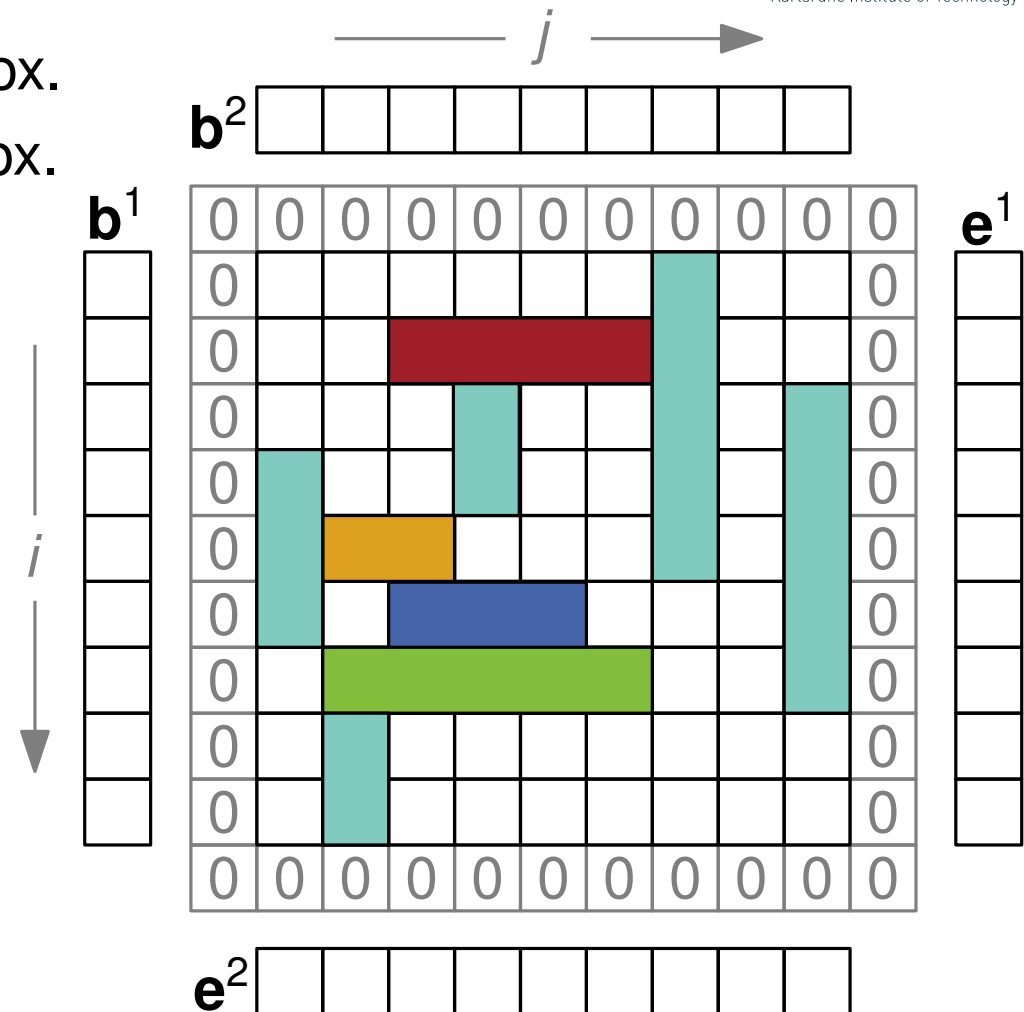
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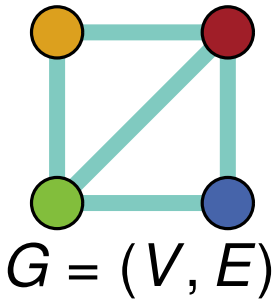
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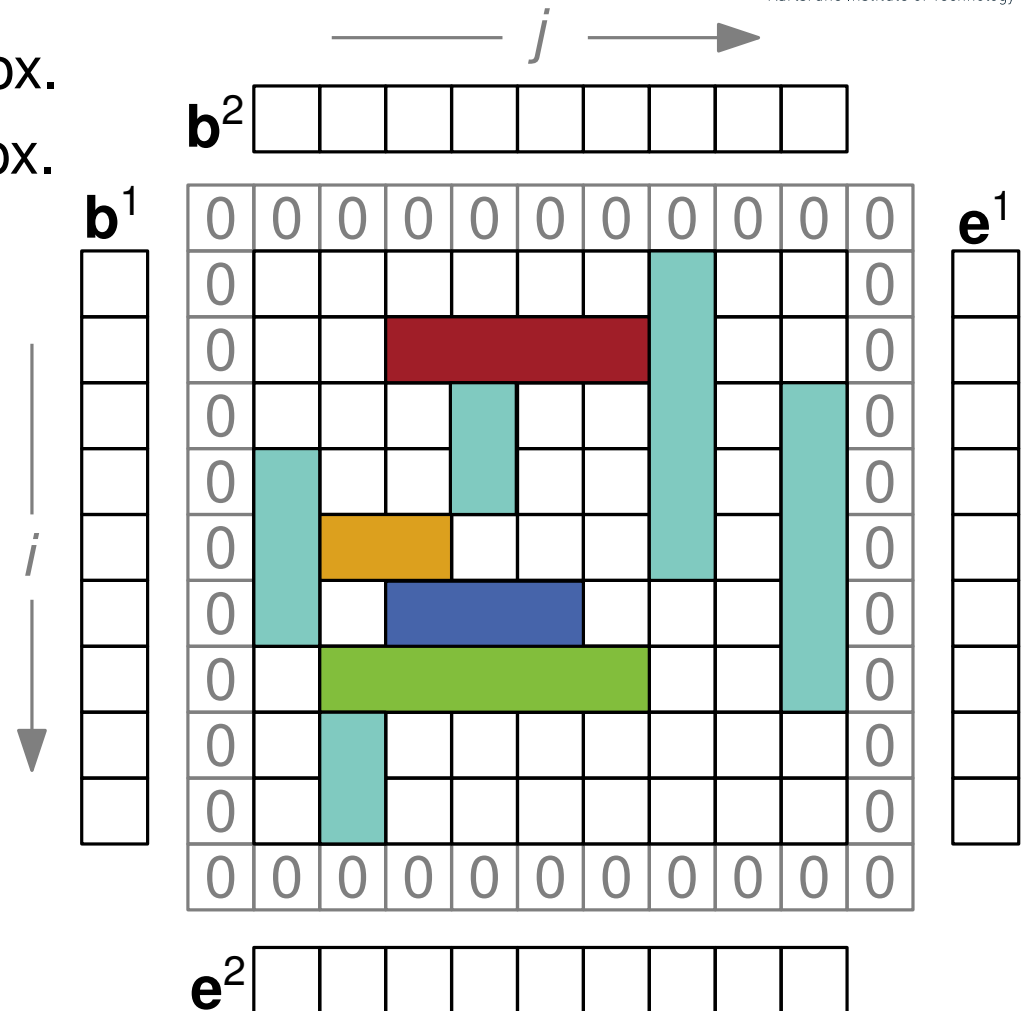
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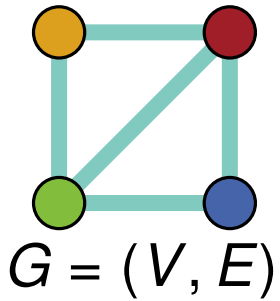
$$\sum_{v \in V} x_{i,j}(v) \leq 1$$

6.) Edges do not overlap non-incident vertices:

$$\sum_{v \in V \setminus e} x_{i,j}(v) \leq (1 - x_{i,j}(e)) \text{ for all } e \in E$$

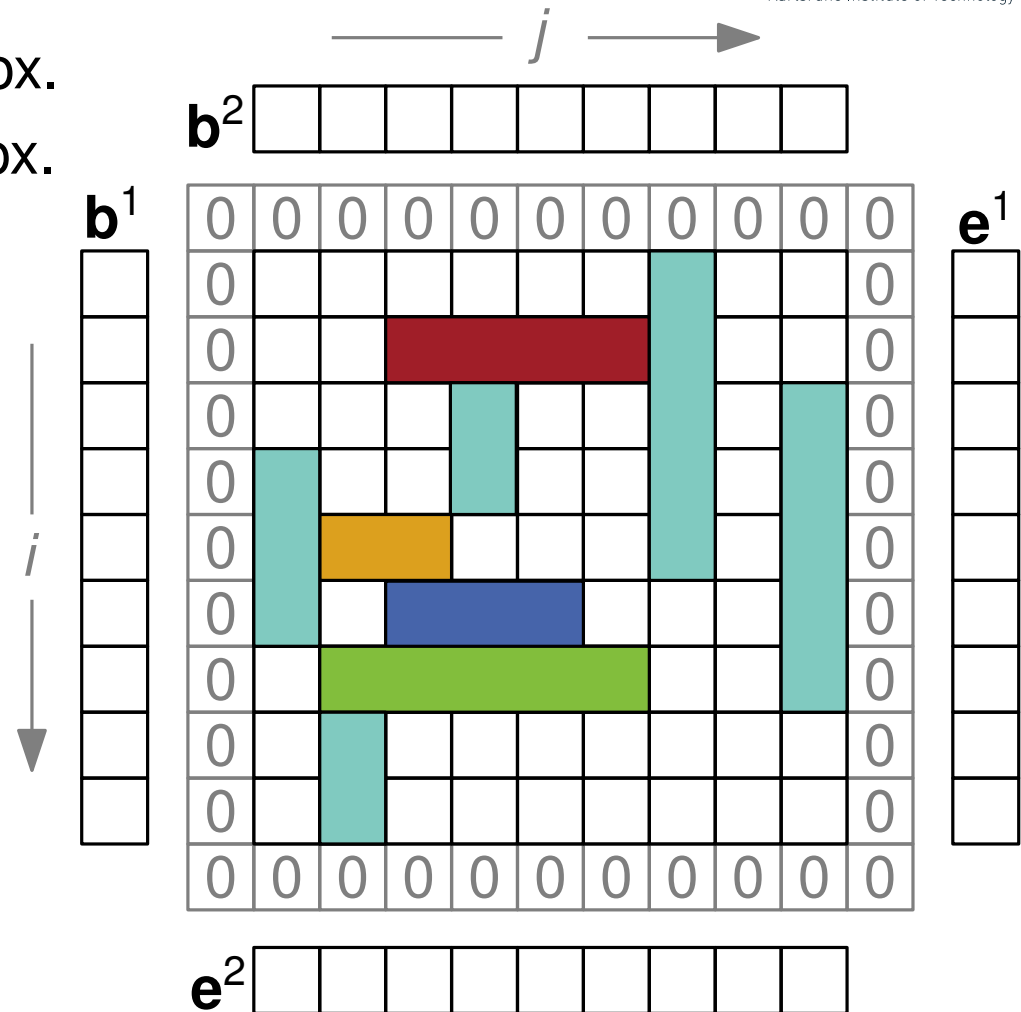


Modeling Visibility Representation

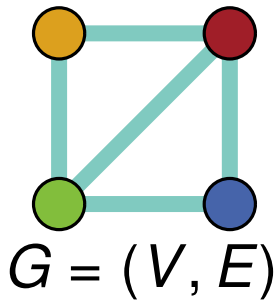


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7.) Edges intersect incident vertices:



Modeling Visibility Representation



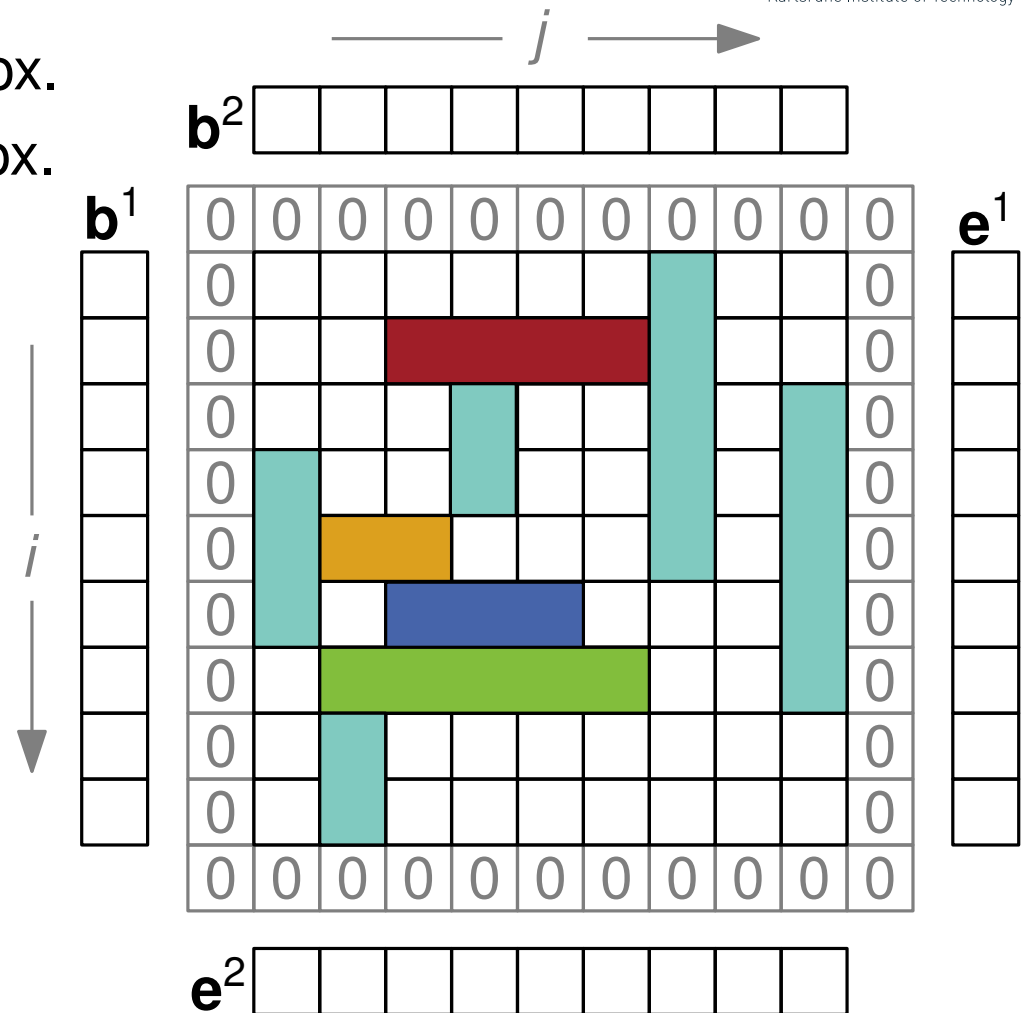
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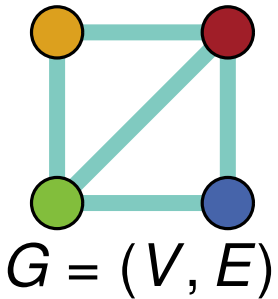
For all $e = \{u, v\} \in E$:

a.) Grid of binary variables

$$x_{i,j}(e, u) \quad x_{i,j}(e, v)$$



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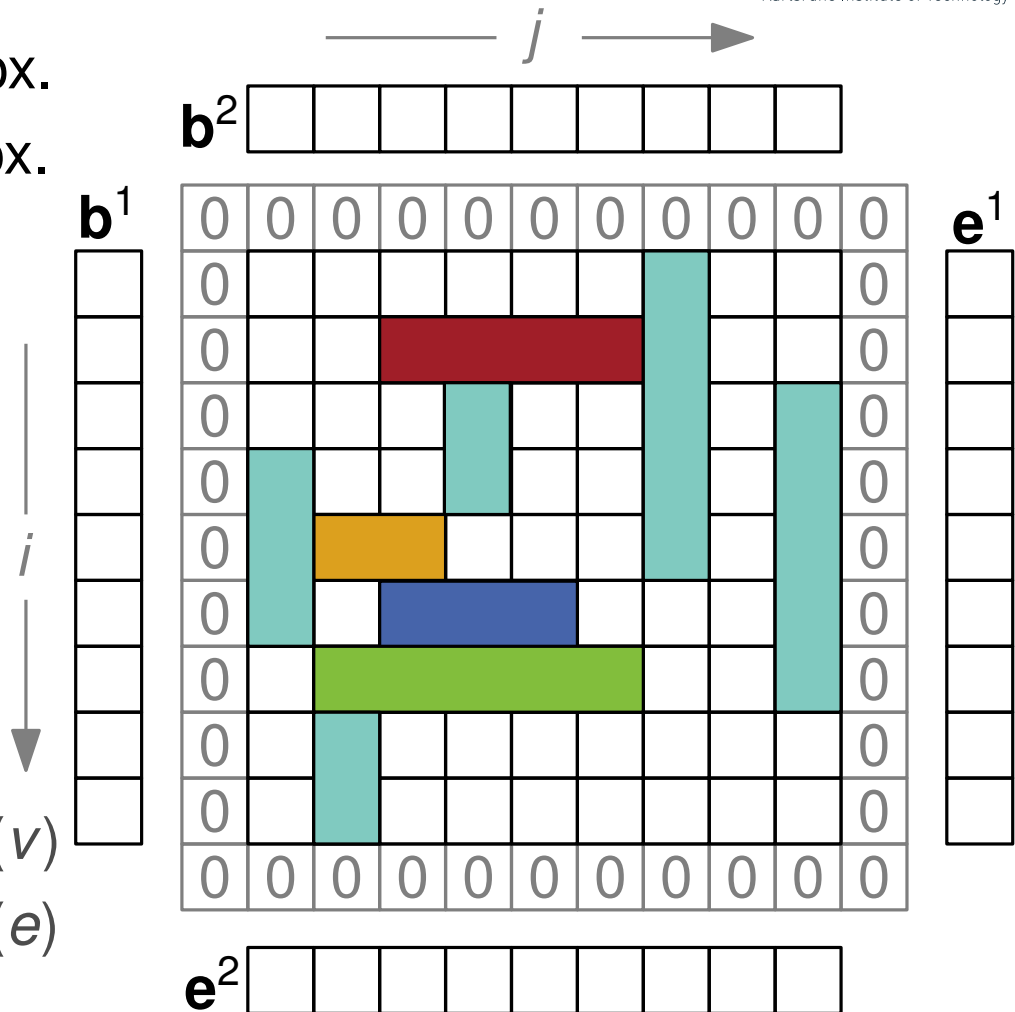
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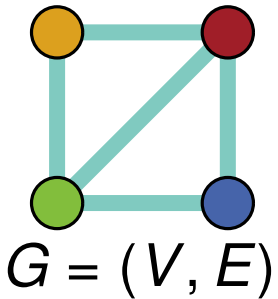
b.) Common point

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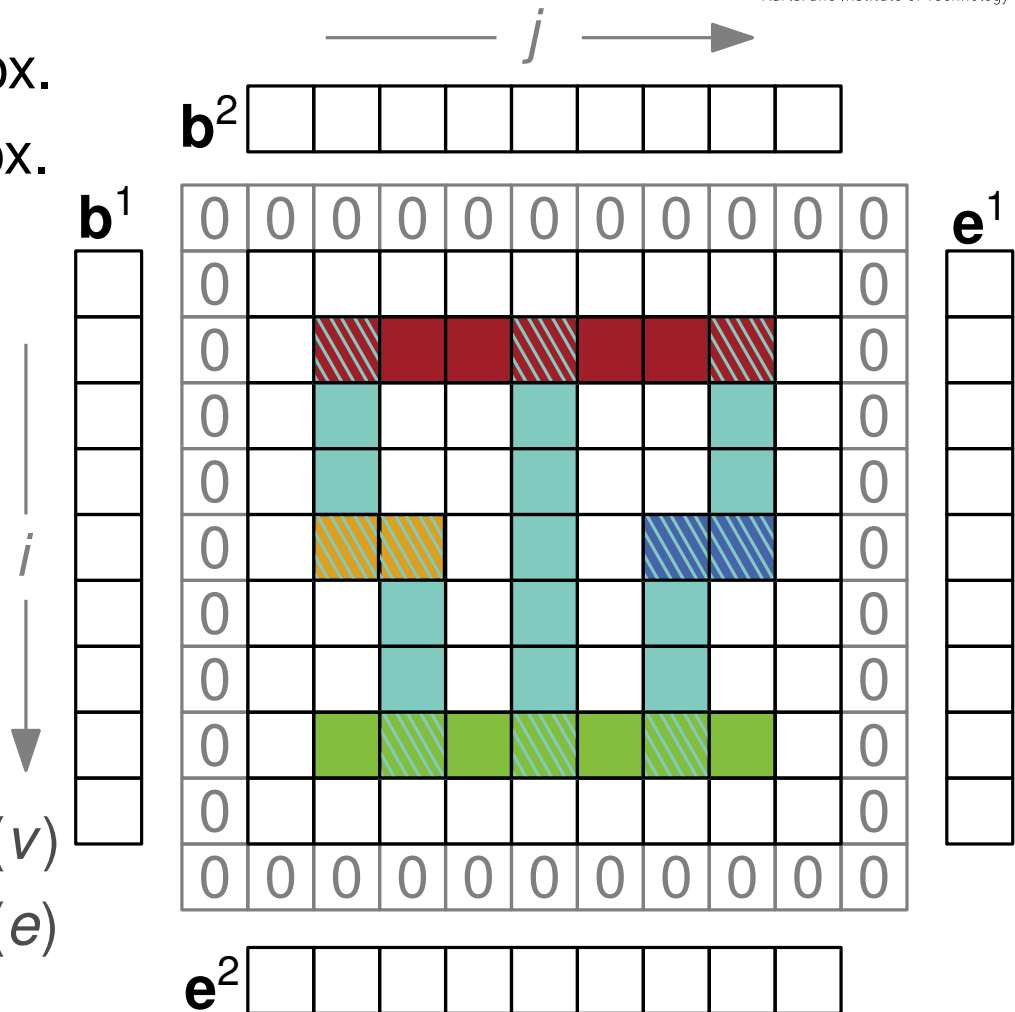
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c.) Existence of common point.

$$\sum_{i,j} x_{i,j}(e, u) \geq 1 \quad \sum_{i,j} x_{i,j}(e, v) \geq 1$$



Optimum st -orientation

Recall: An st -orientation of an undirected graph $G = (V, E)$ and two vertices $s, t \in V$ with $(s, t) \in E$ is an orientation of the edges in E s.t.

- s is the unique source
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Properties:

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- can be computed in linear time
- ingredient for orthogonal layout algorithm (see lecture 14)
- NP-complete to find st -orientation that minimizes the length of the longest st -path, even if G is planar

[Sadasivam, Zhang '10]

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[Sadasivam, Zhang '10]

Find minimum-length *st*-orientation using 1D grid-based model.

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But is it easy to use in practice?

PIGRA Demo (Fabian Klute)